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Mobility Support in IPv6

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> Presented By: Ajay Sharma.







About The Author

Charles E. Perkins: Research Fellow at Nokia Research Center investigating mobile wireless networking and dynamic configuration protocols. He is the editor for several ACM and IEEE journals for areas relating to wireless networking. Charles has served on the Internet Architecture Board (IAB) and on various committees for the National Research Council. He has published a number of papers and award-winning articles in the areas of mobile networking, resource discovery, and automatic configuration for mobile computers.

David B. Johnson: Associate Professor of Computer Science and Electrical and Computer Engineering at Rice University . He was a principal designers of the IETF Mobile IP protocol for IPv4 and primary designer of Mobile IP for IPv6. Currently an Executive Committee member and the Treasurer for SIGMOBILE, also a member of the Editorial Board for IEEE/ACM Transactions on Wireless Networks. 2





Outline

- 1. Why Mobile IPv6
- 2. Benefits of Mobile IPv6
- 3. What is IPv6?
- 4. Address Architecture of IPv6.
- 5. Mobile IPv6 Terminology.
- 6. Mobile IPv6 Mechanism.
- 7. Errors Handling
- 8. Security Handling
- 9. Summery
- 10. Q & A



Why Mobile IPv6? -- Propellant factors.

DSL

Cable

IPv6

all-IP

Packet Core

3G-

SGSN **3G RAN**

3G-

3G-

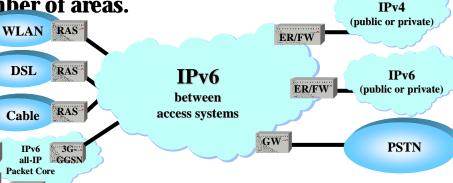
SGSN

3G--

SGSN

3G RAN

- Huge growth of mobile Internet terminals will exhaust IPv4 address space
 - All wireless terminals will have WAP and GPRS
 - IPv6 brings enough IP addresses
- Ease of scalability ٠
 - Supporting billions of new devices and huge amounts of new bandwidth
 - Simplified, cost-efficient architecture without NATs, Proxies, ALGs,...
- Always-on connection establishes a variety of new services.
 - Push, location-based, etc.
- **Integrated Security**
- Efficiency: IPv6 improves efficiency in a number of areas.
 - Routing, Broadcast handling
- **Quality of Service improvements** – Fragmentation, Flows
- Mobility Across Access Technologie ۲







Requirements for Mobility in Internet

Mobility

Transparency

Easy to use

Routing

Security

- Increasing number of users asks for Mobility Support in Internet
- Mobility shall be transparent to all Protocol Layers above IP
- Mobility shall be as easy to handle as with Mobile Phones in GSM
- Mobility shall be compatible to all Routing Protocols and shall optimize routes
- Mobility shall not decrease security in Internet



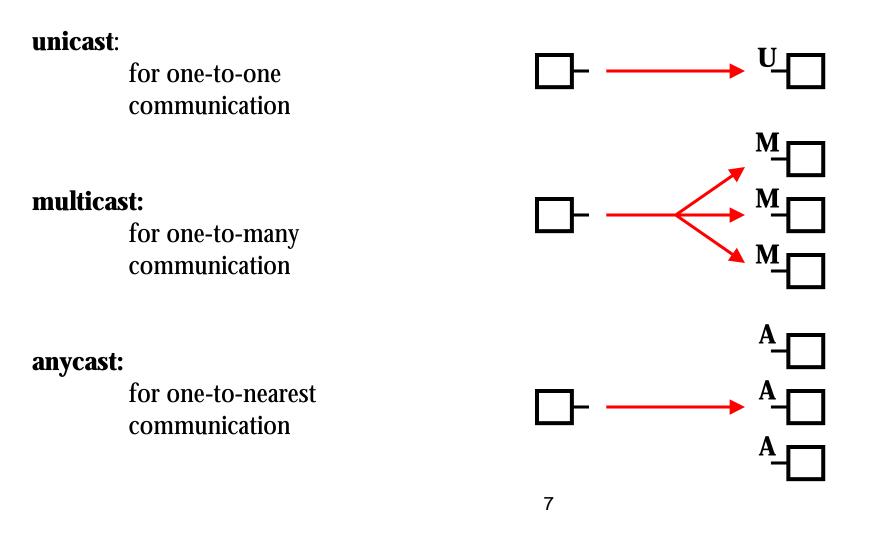
IPv6 features relevant to Mobile IP

- Larger address space => Unique Global address for each device. (6.65 · 10²³ addresses per m² of earth surface)
- Scalable => Run over multiple media i.e. Wireless-LAN, Ethernet, 3G
- Auto configuration capabilities=> Network Plug-and-Play.
- Fixed header format => Fewer fields (8 as compared to 12 in IPv4)
- Router headers => MIP updates are in extension headers. No header length anymore.
- Security extensions => Internet level Security in IPv6 Header.
- Anycast addresses => Special type of address in IPv6.
- Encapsulation =>IP-layer authentication & encryption possible.
- Quality of service and flow labels => efficient routing for real-time applications.
- Elimination of "triangle routing" for mobile IP
- All nodes can handle bindings.
- Small overhead for distributing bindings. Fixed header format
- option extension headers not parsed by intermediate routers anymore





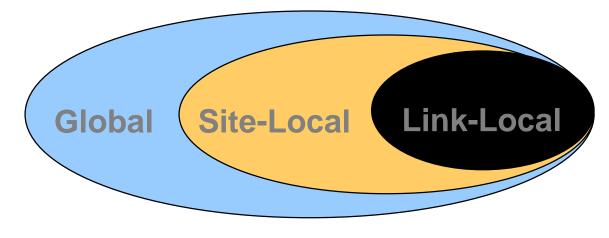
Basic IPv6 Address Types





IPv6 - Addressing Model

- addresses are assigned to interfaces
 - No change from IPv4 Model
- interface 'expected' to have multiple addresses
- addresses have scope
 - Link Local
 - Site Local
 - Global
- addresses have lifetime
 - Valid and Preferred lifetime







Text Representation of IPv6 Address

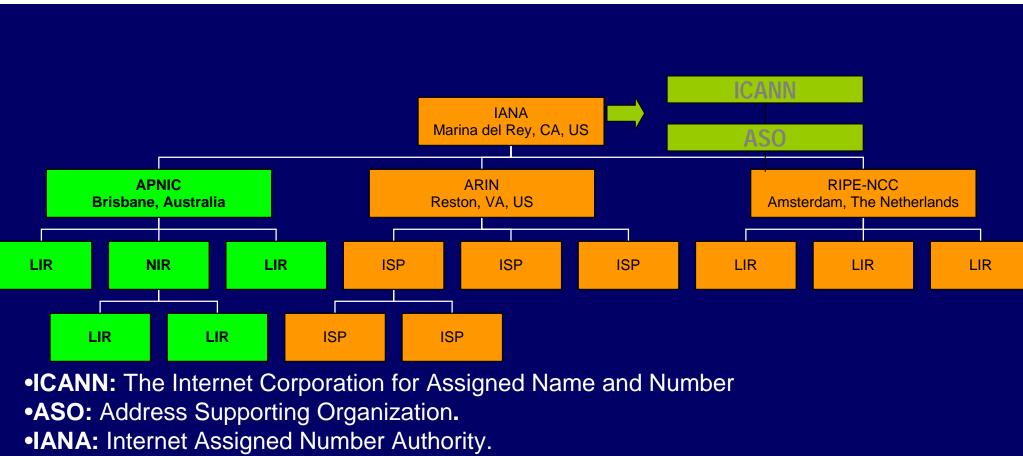
"Preferred" form: 1080:0:FF:0:8:800:200C:417A Compressed form: FF01:0:0:0:0:0:0:0:43 becomes FF01::43 IPv4-compatible: 0:0:0:0:0:0:13.1.68.3 or ::13.1.68.3 There is no broadcast addresses, only multicast. Loopback address is ::1



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Internet Registry Hierarchy



•ARIN: American Registry for Internet Number.

•APNIC: Asia Pacific Network Information Centre.

•RIPE-NCC: Reseaux IP Europeene.







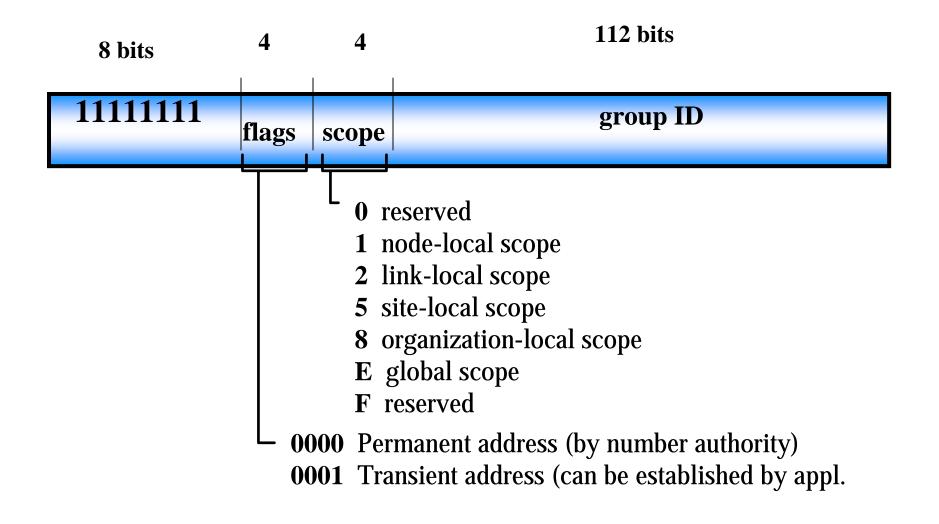
IPv6 Address Formats

Aggregatable Global	3	13	8		24			16			64		
Aggregatable Global Unicast Address	001	TLA ID	RES NLA ID			SLA ID Int			Interface I	D			
Address Format with	3	13	13	n	m	19-n-m	0	16-0			64		
Address Format with Field Substructures	001	TLA ID	Sub-TLA	NLA1	NLA2	Site ID	SLA1	Subnet			Interface I	D	
Link Local		10 54							64				
Link-Local Unicast Address	1111	1111010		0							Interface I	D	
Cite Level		10	38				16	64					
Site-Local Unicast Address	111111011			0			Sub	onet ID	Interface ID			D	
Reserved Subnet				6	64				6	1	50		7
Anycast Address	Subnet Prefix								111111	ul	1		Anycast ID
	8 4 4 80							32					
Multicast Address	11111111 Flags Scope					0				Group I	D		





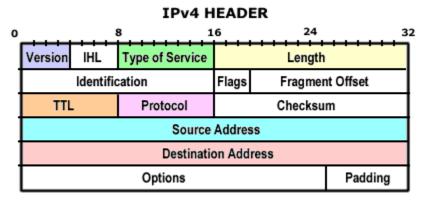
Multicast address







IPv4 vs. IPv6 Header



- 14 fields, at least 20 octets
- 32 bit addresses
- fragmented packet processing at every hop
- header checksum recalculation at every hop
- variable Options field for extra processing information

	IPV6 HEADER								
0		8	1	6	24 3				
	Version	Traffic Class	Flow Label						
		Payload Length		Next Header	Hop Limit				
4	Source Address								
8	Destination Address								

IPv6 HEADER

- 8 fields, fixed 40 octet size
- 128 bit addresses
- fragmentation only in src and dst endpoint, or lower layer
- no checksums
- new 20 bit flow label field
- options in Extension Headers



Changes in IPv4 Header

- 20 bytes
- 13 fields
- removed
- moved to extension headers

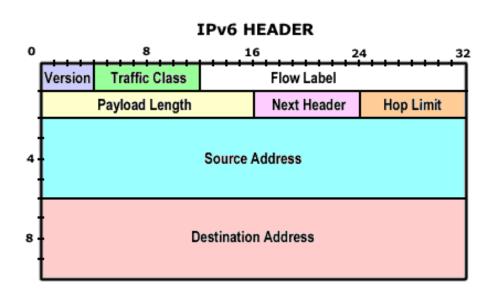
- renamed
 - –precedence \rightarrow class
 - -total length
 - \rightarrow payload length
 - –time to live \rightarrow hop limit
 - -protocol \rightarrow next header

Version Hdr Len	Prece- dence ToS	Total Length					
Identif	ication	Flags Fragment Offset					
Time To Live	Protocol	Header Checksum					
Source Address							
Destination Address							





IPv6 Header Simplifications



• Base header is fixed size - 40 octets

<u>Simplifications</u> Fixed format headers

no options -> no need for header length options expressed as Extension headers No header checksum

> reduce cost of header processing, no checksum updates at each router minimal risk as encapsulation of media access protocols (e.g..., Ethernet, PPP) have checksum

No segmentation

hosts should use path MTU discovery otherwise use the minimum MTU (536 bytes)

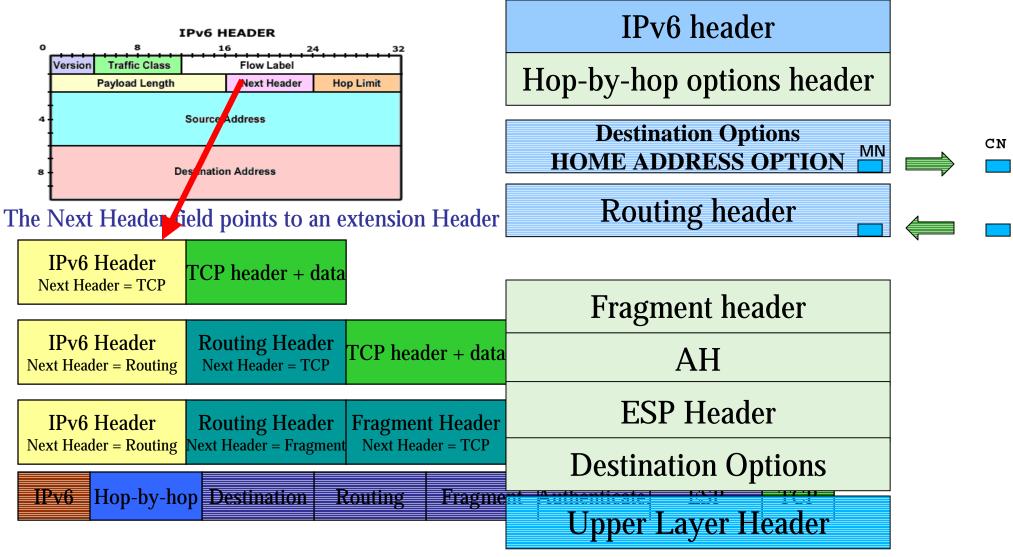
- NEXT HEADER field in base header defines type of header
- Appears at end of fixed-size base header
- Some extensions headers are variable sized
 - NEXT HEADER field in extension header defines type
 - HEADER LEN field gives size of extension header





Extension Header

Store optional internet-layer information [Placed between IPv6 header and upper-layer header]



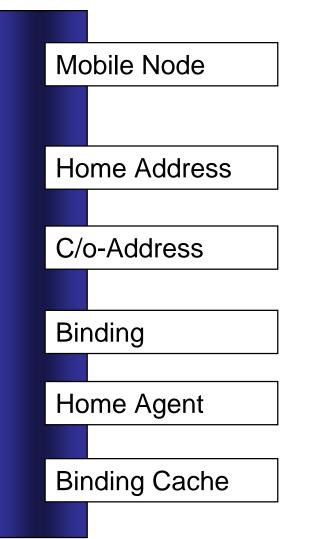


Extension Header

Hop-by-hop **Destination** TCP IPv6 Routing Fragment Authenticate ESP ins **Enge**man done by Carries Bittidinal Alts **よPath BAGBOOGRA** must bethisardingeneyeeten so yet the sed to along the sticked with the determine smallest allowed packet size. **Does not give** Supports data authentication for IP confidentiality. Charlie can't read header fields that change Alice or Bob's value along route. Message (ESP) Alice is Alice, Bob is Bob **(AH)** 17 W DI



Terms used in Mobile IPv6



Node, which can change its access point to the Internet while still being reachable under its Home Address.

Static IP Address of the Mobile Node valid at its home network.

Temporary IP Address of the Mobile Node valid at the actually visited network of the Mobile Node (c/o = care-of).

Association of the Home Address with the c/o-Address.

Router located at the Mobile Node's home network used by the Mobile Node for registering its c/o-Address.

Cache for received Bindings.



Binding Update Option Header Format

4 bytes (32 bits)										
1	Vex	t Header	Hdr Ext Len	Option Type = 16	Option Length					
A	Н	L F	Reserved	Lifeti	me					
	Identification (8 bytes)									
					-					
	Mobile node's Home Address (16 bytes)									
	-									
	Care-of Address (16 bytes)									
	-									
	-									
	Mobile node's Link-Local Home Address (16 bytes) (only present if 'L' bit equals 1)									
	(only present in Libit equals 1)									

- **A Bit** : Indicates whether receiver should reply or not with Binding Acknowledgement.
- **H Bit**: Use when mobile node wants the receiving node to act a Home Agent.
- **L Bit:** Set if the mobile node want to receive packet destined to its link-local address.
- Lifetime: Lease time for the address.
- **Identification Field**: Counter is use to insure Binding Updates are order-wise. Counter increment for each new BU (not for retransmission).
- **Care-of Address**: current address of MN. When care-of address = Home address. Destination Cache entries should be deleted.





Server-less Autoconfiguration ("Plug-n-Play")

Host autoconfiguration: Host autoconfiguration is a mechanism whereby addresses and other parameters can be assigned to network interfaces. This can be done in two different ways, known as stateful and stateless autoconfiguration. Duplicate Address Detection (DAD) is also performed here.

Router autoconfiguration: Neighbor Discovery protocol the mechanisms for automatic router configuration Keeping a router updated means ensuring that it has an exact knowledge of the organization of the subnet to which it is connected, which in turn means assigning the correct prefixes to each link with which the router has an interface.

DNS autoconfiguration: To facilitate man-machine interfacing, applications generally handle domain names rather than numerical addresses. DNS, database contains name-address mappings for each Internet domain. A6 record type has been defined facilitate the adoption of an automatic DNS management mechanism.

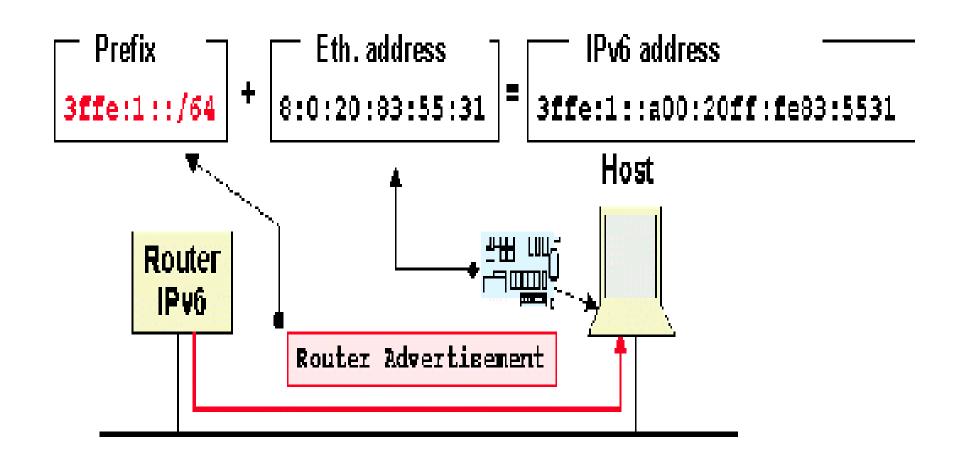
Service autoconfiguration : to make use of the services available on the network, users must know at least the name of the network host on which they are installed. Service Location Protocol (SLP), which provides a flexible and scalable structure whereby hosts can access information concerning the existence, location and configuration of network services.





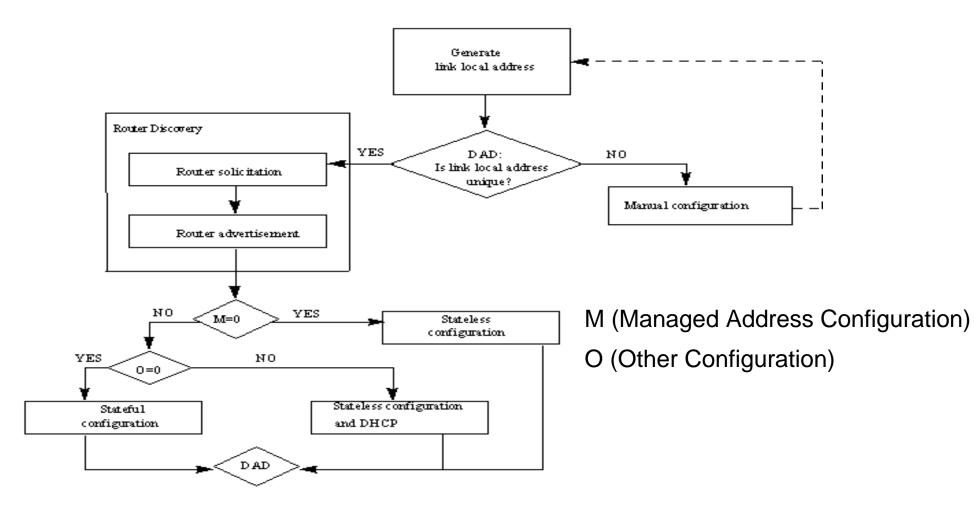


Configuring Network Prefix





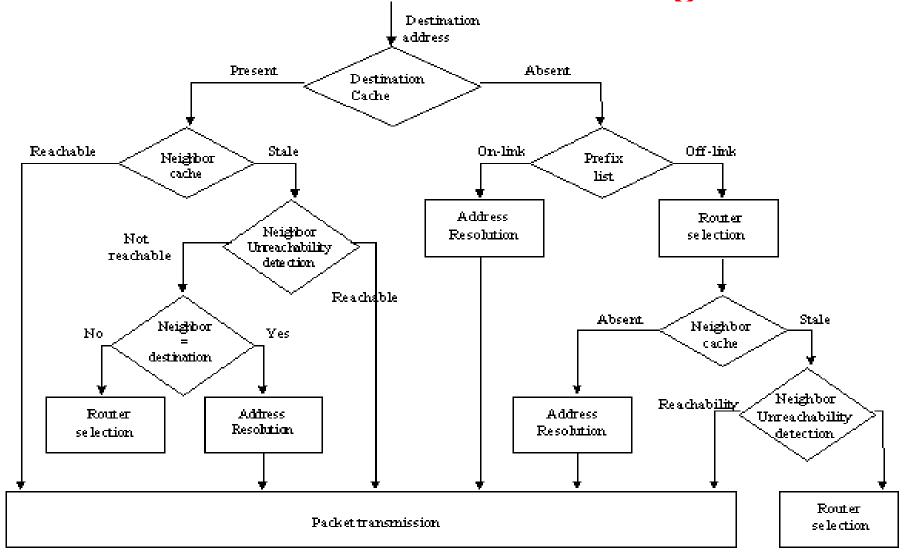
Autoconfiguration Algorithm







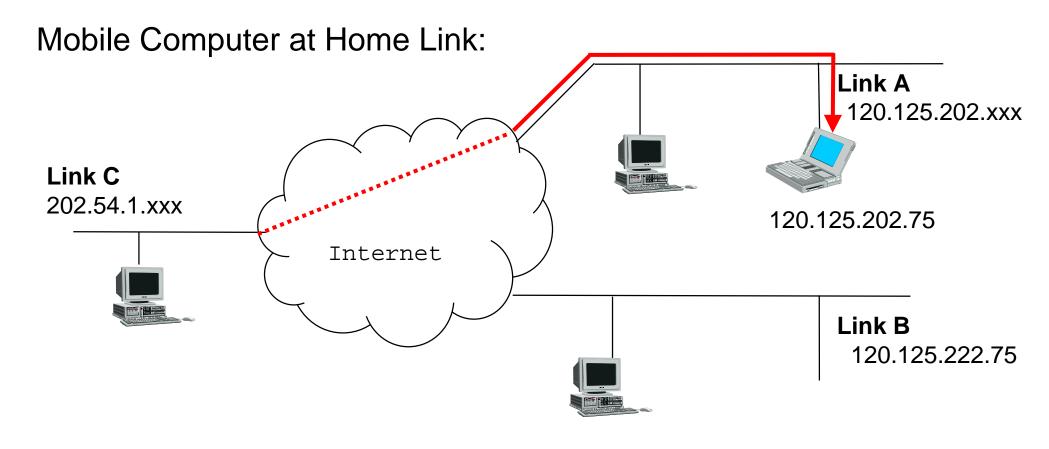
Packet Transmission Algorithm







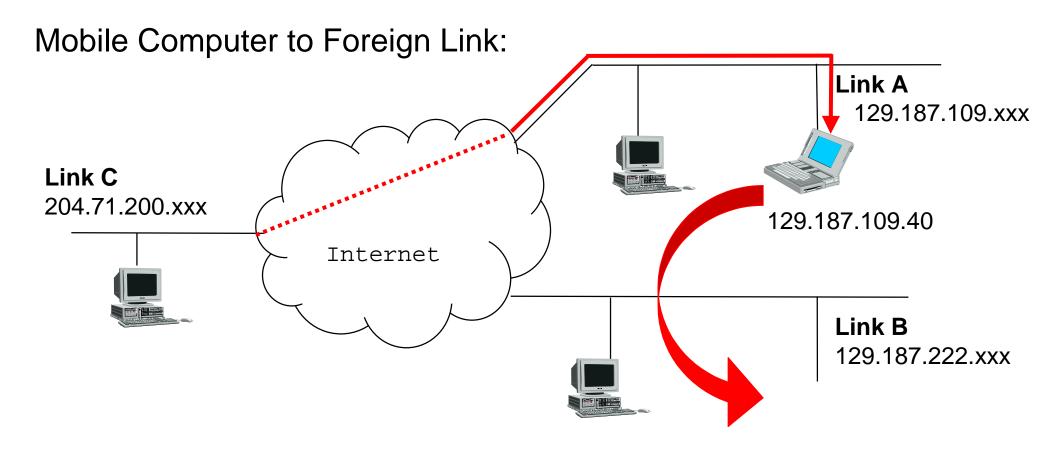
Mobility Problem with IPv4







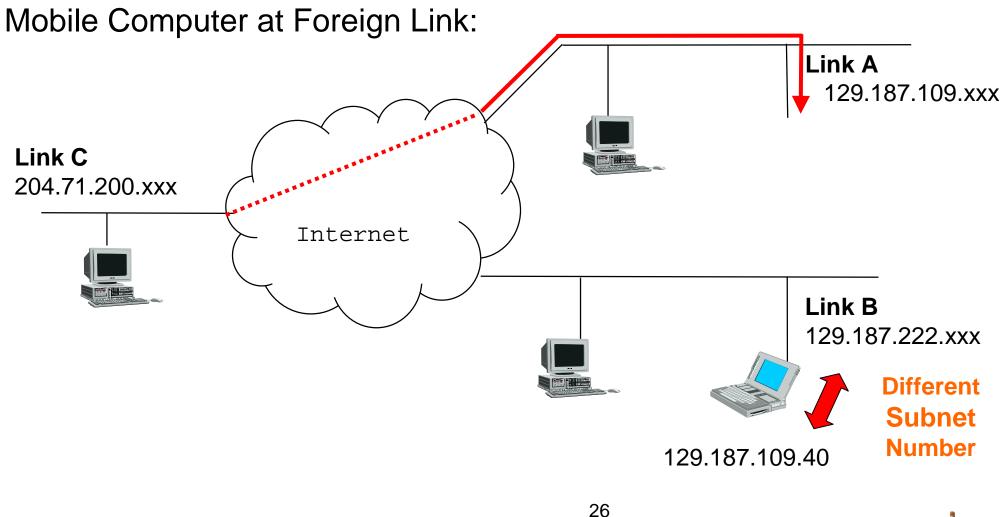
IP Mobility Problem with IPv4







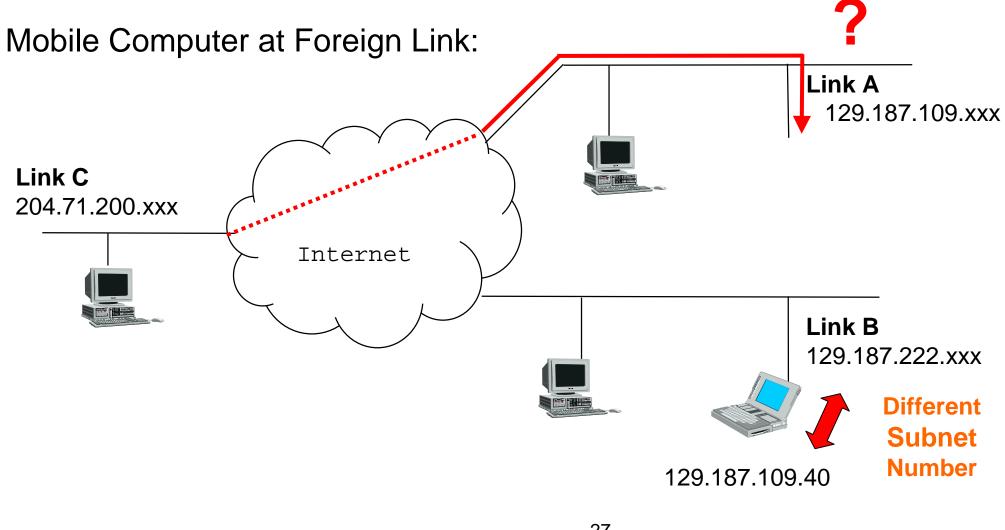
IP Mobility Problem on Movement







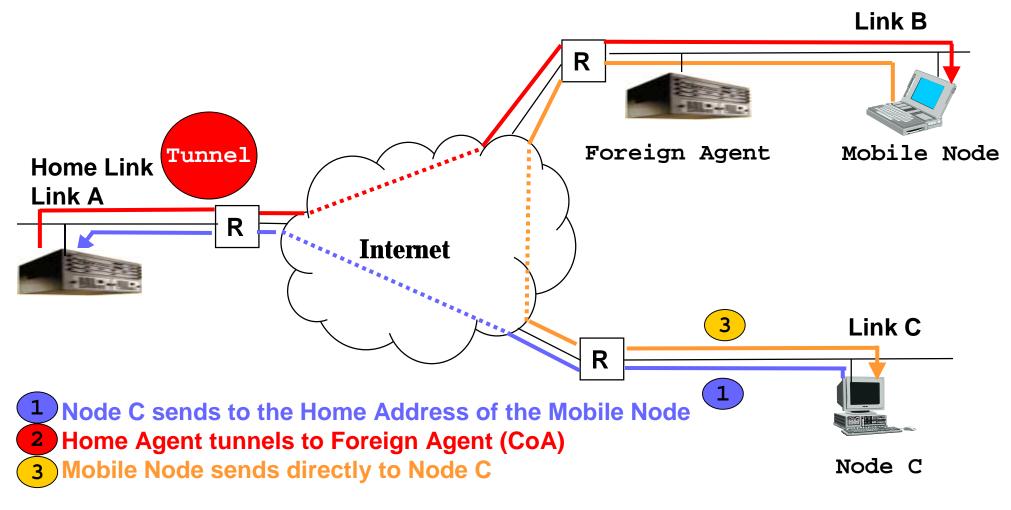
IP Mobility Problem with IPv4



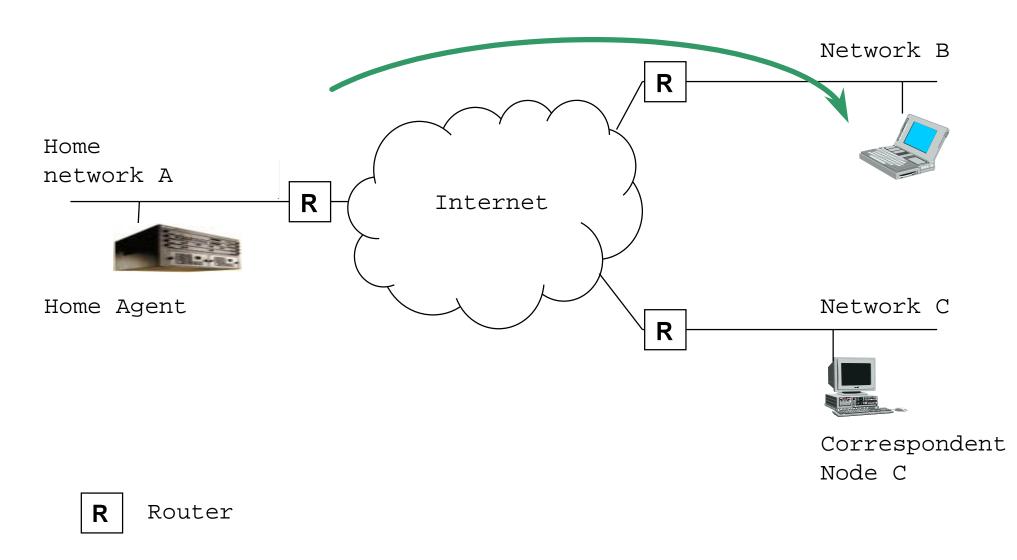




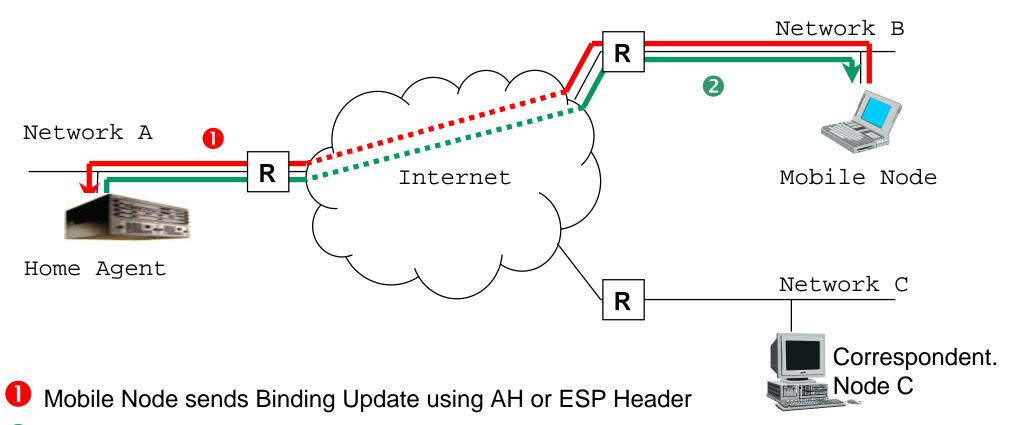
Packet Delivery with IPv4



Mobile Node Moves: IPv6 consideration



Mobile Node registers at its Home Agent

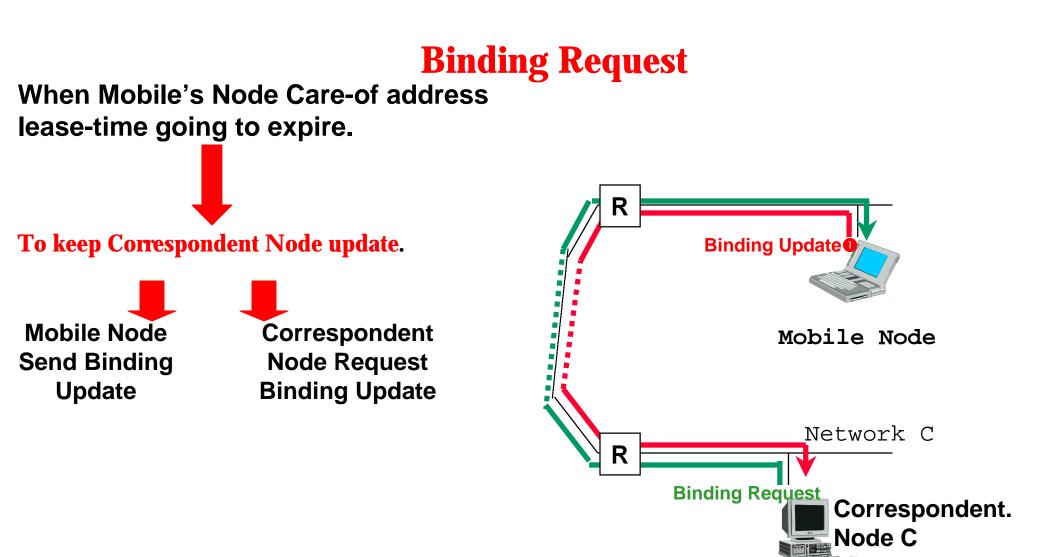


2 Home Agent replies with Binding Acknowledgement using AH or ESP Header







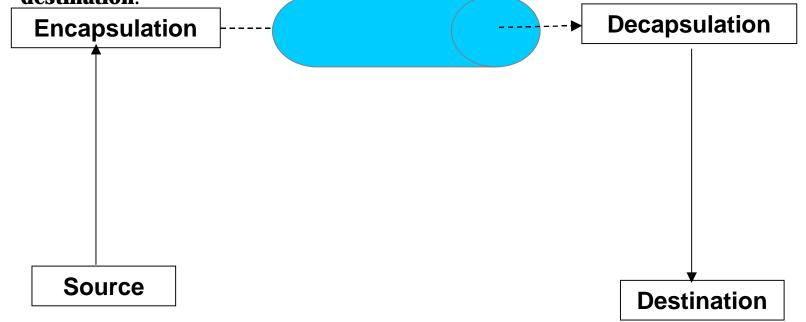






Tunneling

Tunnel: The path followed by a datagram while it is encapsulated. While encapsulated, a datagram is routed to a knowledgeable agent, which decapsulates the datagram and then forwards it to its ultimate destination.

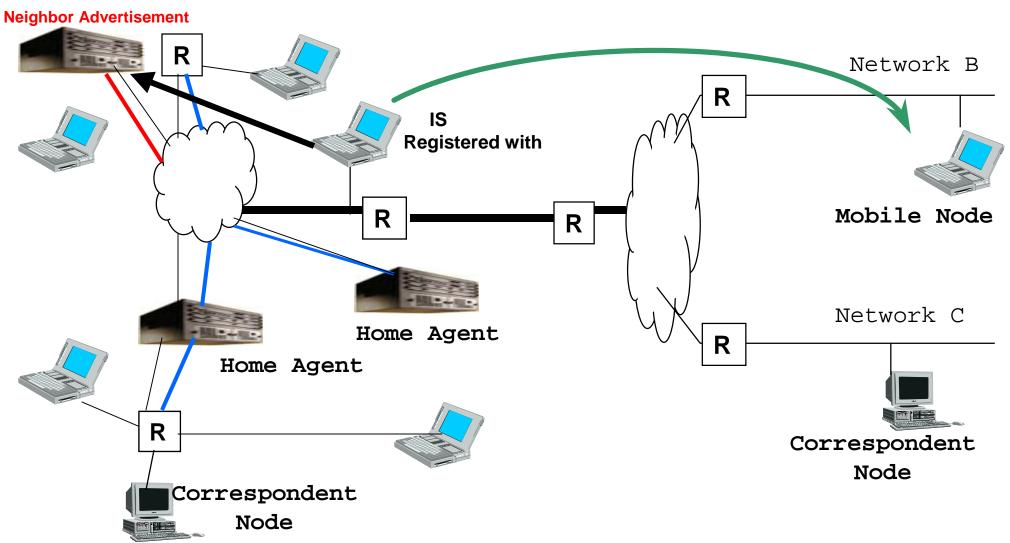






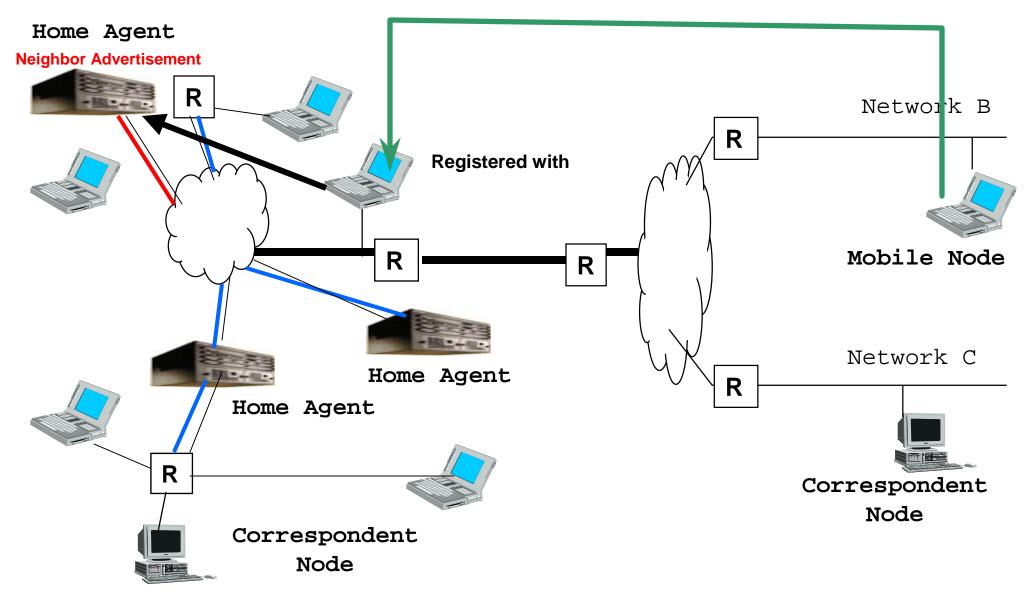
On Mobile Node Movement: HA Takes Action

Home Agent



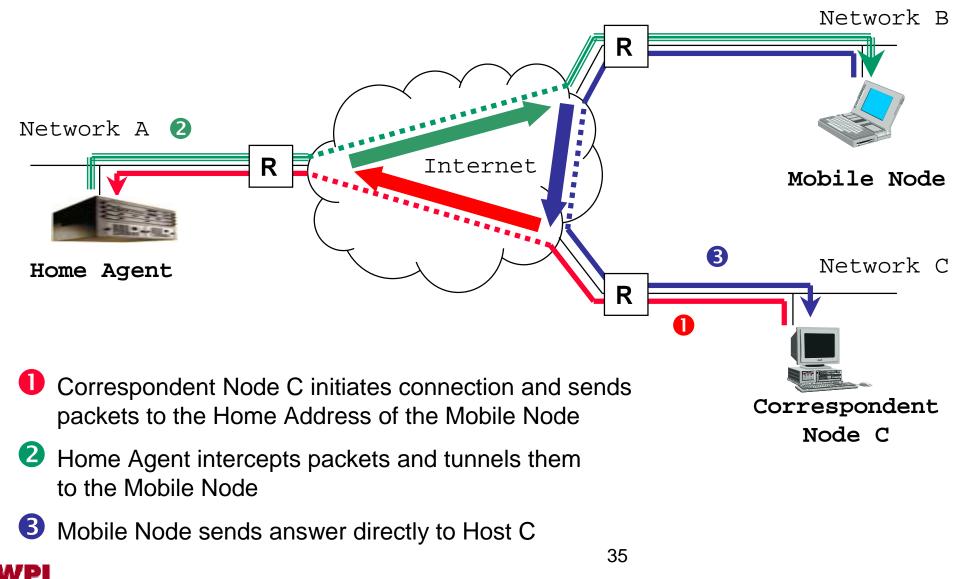


HA Takes Action: When MN Return its Home Subnet



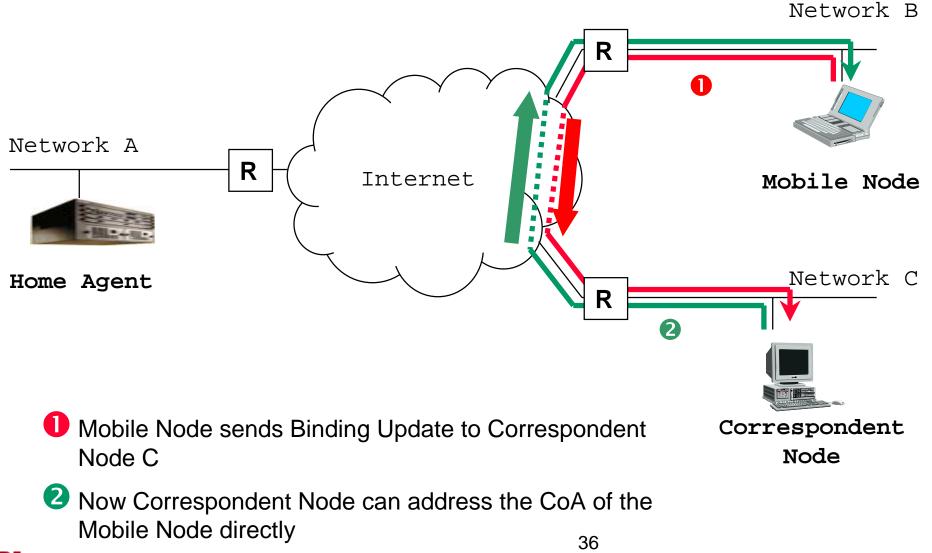


Triangular Routing during Initial Phase



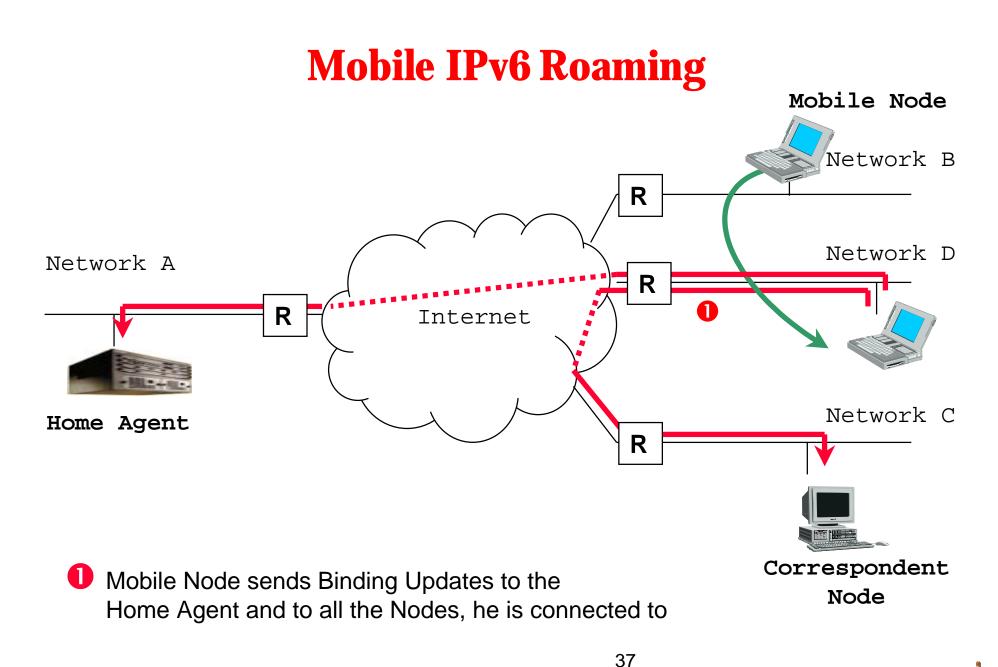


Normal Operation by Route Optimization



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MPI



Movement Detection

Scenario-I : Mobile node to know quickly when the when the Default router will be unavailable

Indicator Neighbor Advertisement unreachable detection by using upper-layer TCP timeout mechanism.

When Mobile node don't receive Neighbor Advertisement Message from default router in response to Neighbor Solicitation message.

Scenario II : When Mobile node become unreachable to default Router

Some sort of time setting its network interface so This that it can receive all the packets through that receive area.

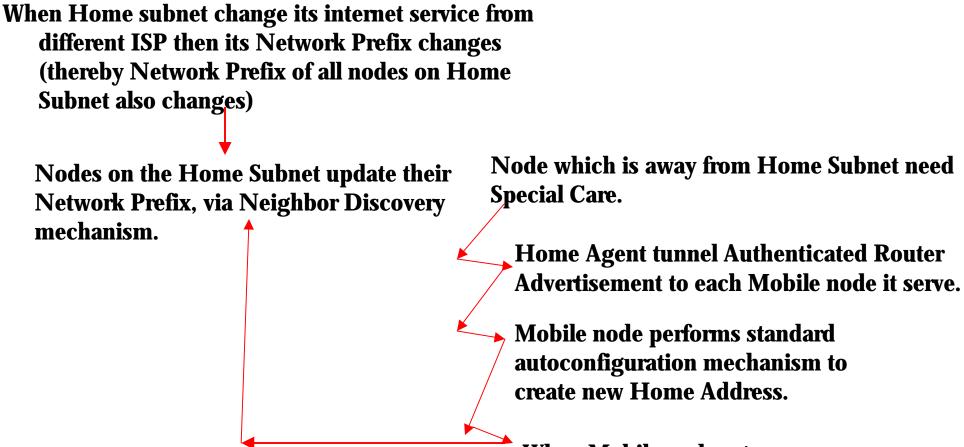
Through Router Advertisement messages. receipt of packets from default router indicate reachable.







Renumbering Home Subnet

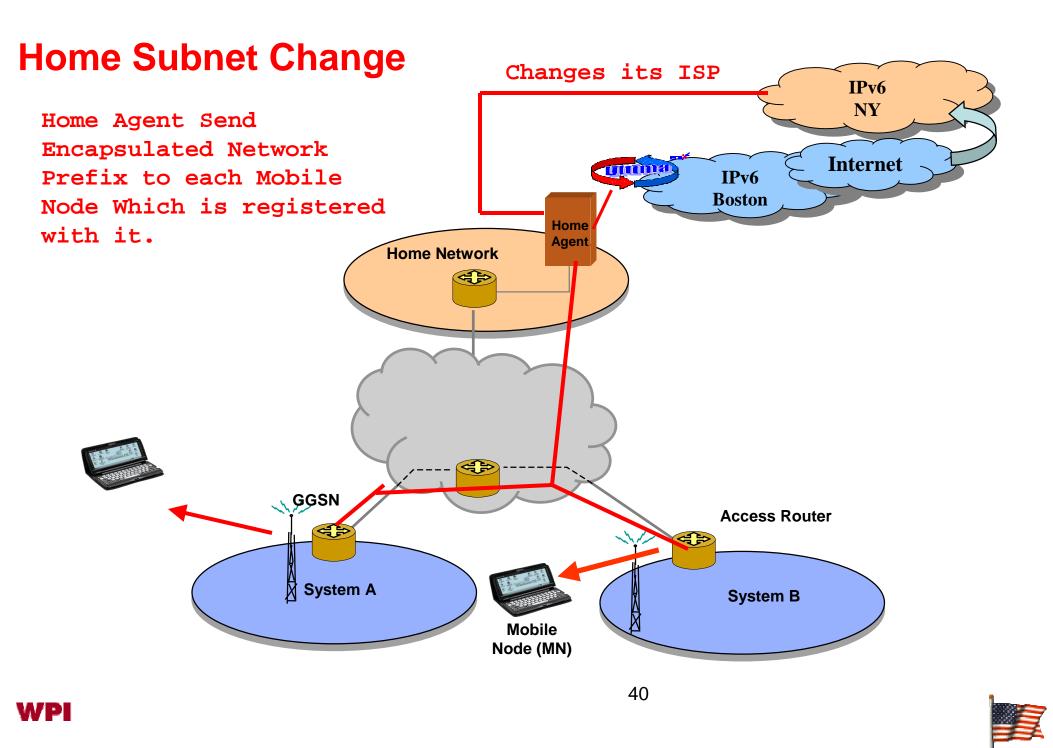


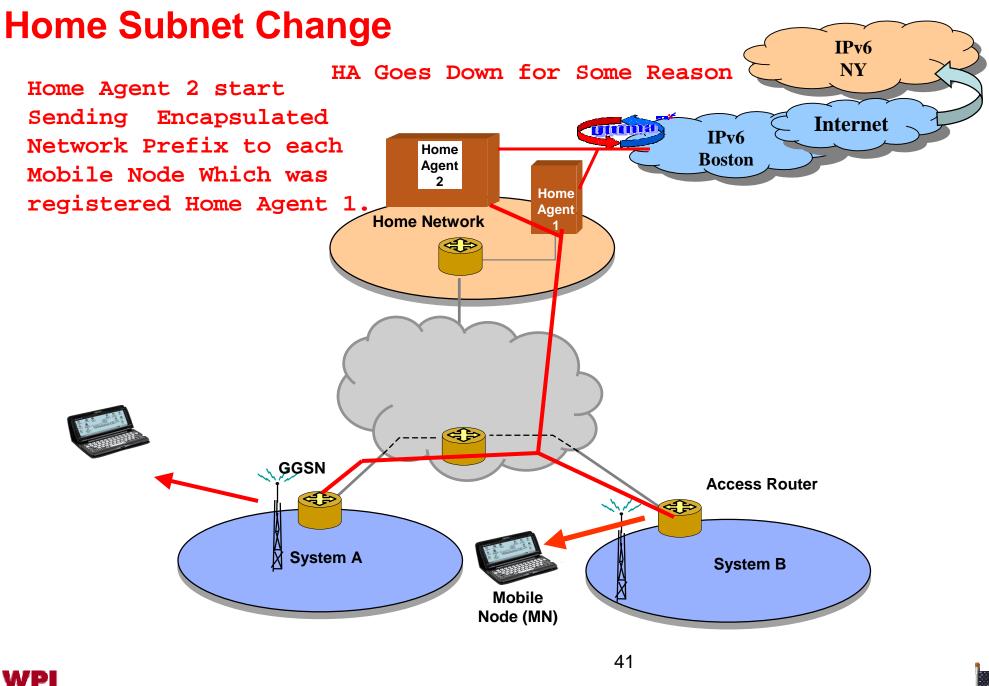
When Mobile node return home, it first performs duplicate address detection.



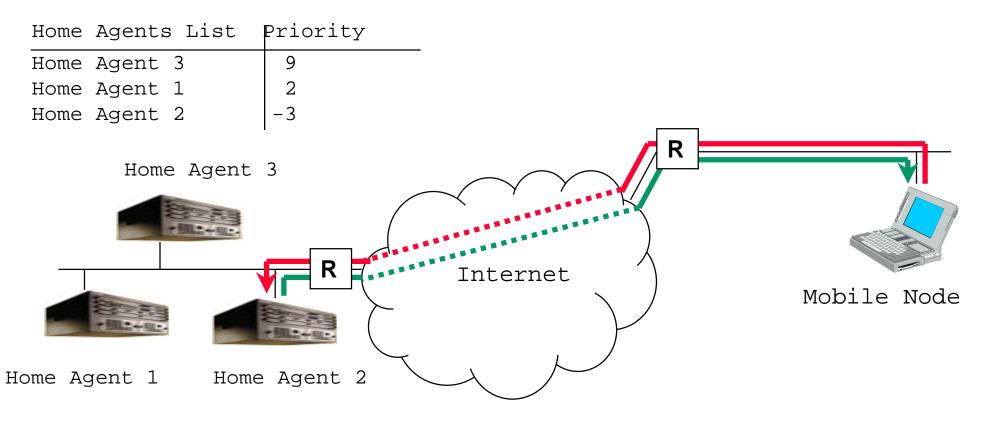








Dynamic Home Agent Address Discovery (Renumbering)

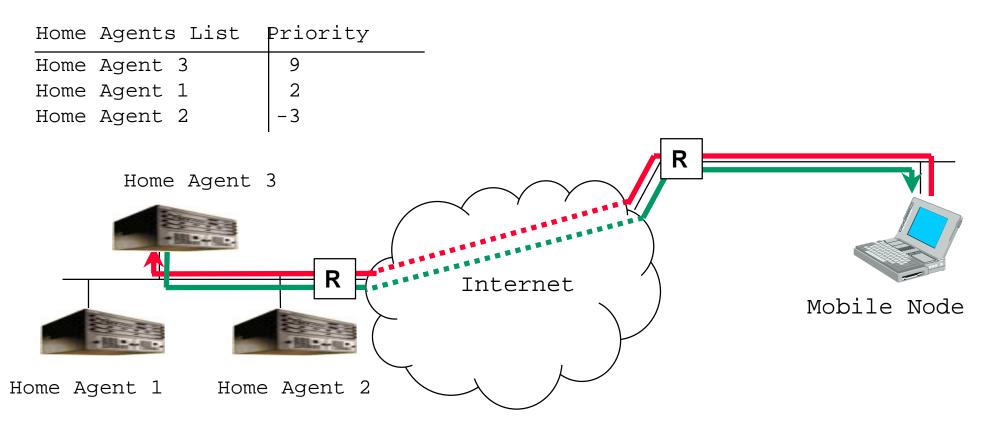


- Mobile Node sends Binding Update to the Home Agents Anycast Address of its home network
- One Home Agent answers with Binding Acknowledgement containing a list of available Home Agents





Registration at selected Home Agent

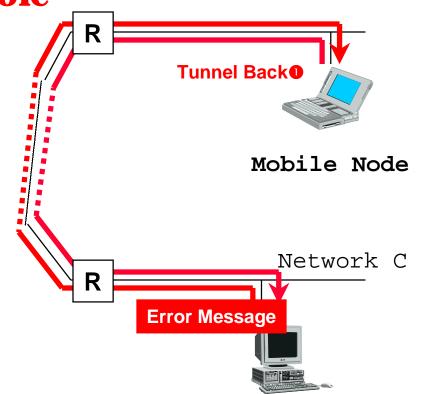


- Mobile Node sends Binding Update to the first Home Agent contained in the Home Agents List
- **2** Binding Acknowledgement completes Registration process



ICMP Role

- When an IPv6 node discards a packet, it sends an error message to the source. There are four types of message:
- 1. Destination unreachable (type=1). Sent by a router to the source when a packet cannot be forwarded to its destination.
- 2. Packet too big (type =2). Used when the link MUT on the forwarding link is smaller than the packet.
- 3. Time exceeded (type=3). Indicates that the packet's hop limit field is zero.
- 4. Parameter problem (type=4). Indicates that a field of the datagram is not recognized as valid and the packet can thus not be processed.



ICMP: Includes the so-called Neighbor Discovery mechanisms, the terminal autoconfiguration mechanisms and address resolution mechanisms.

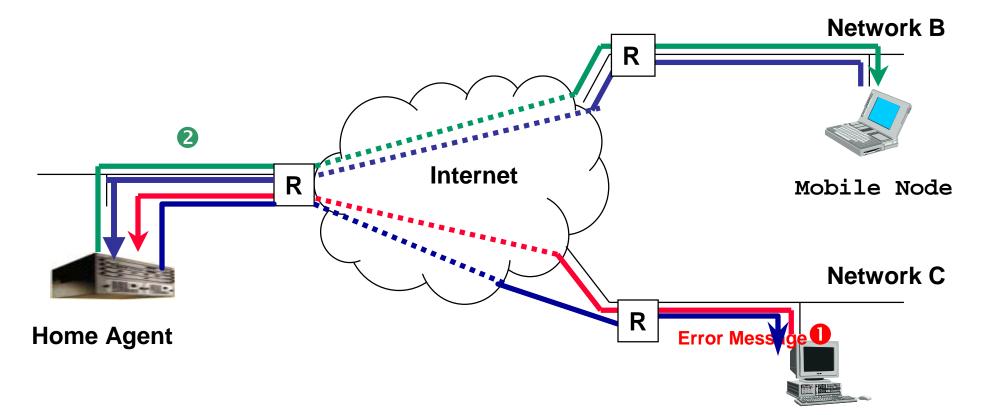






Handling ICMP Scenario 2

When CoN send error message through Home Agent.







Smooth/Fast/Seamless Handover

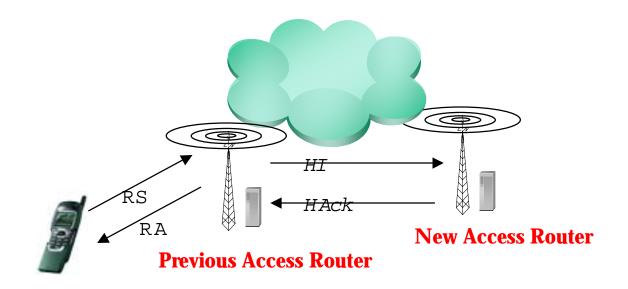
- Smooth handover == low loss
- Fast handover == low delay
 - **30 ms**?
 - Duplicate Address Detection?? (can router pre-empt this?)
- Seamless handover == *smooth* and *fast*







Mobile-controlled seamless handover



One scenario: mobile sends special Router Solicitation (RS)

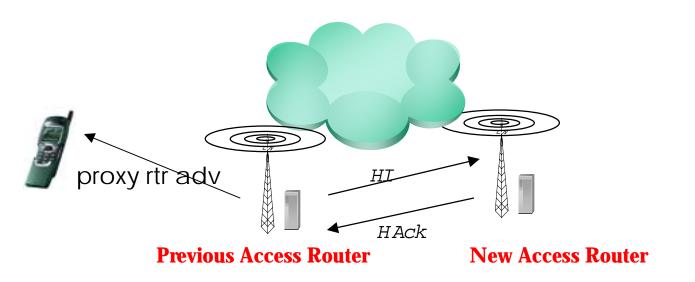
- Previous Access Router replies with *Proxy* Router Advert. (RA)
- Previous Access Router sends Handover Initiate (HI)
- New Access Router sends Handover Acknowledge (HACK)







Network Controlled Handover



- Previous access router sends Proxy Router Advertisement on behalf of the new access router contains prefix and lifetime information, etc.
- Previous access router sends *Handover Initiate* message to new access router
- Mobile node *MAY* finalize context transfer at new access router





Ongoing Work for Open Questions

- Security issues: Firewalls, cause difficulty for Mobile IP because they block all classes of incoming packets that do not meet specified criteria.
- Ingress filtering: Many border routers discard packets if the packets do not contain a source IP address configured for one of the enterprise's internal networks
- Deficiency of Mobile IPv6, is that it does not support fast handoff – (this is the ability to switch to another subnet without significant delay or loss of packets). Excessive signalling in rapidly changing cells.

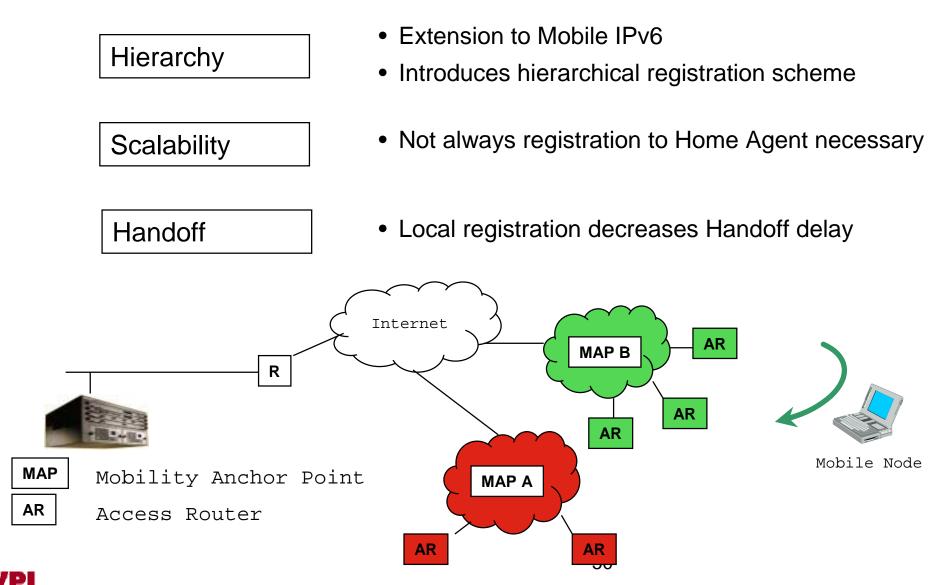
- -Gupta and Glass have proposed a firewall traversal extend Mobile IP operation across firewalls, even when multiple security domains are involved.
- Montenegro has proposed the use of reverse tunnels to the home agent to counter the restriction imposed by ingress filtering.

Extension to Mobile IPv6 called "HIERARCHICAL MOBILE IP v6".



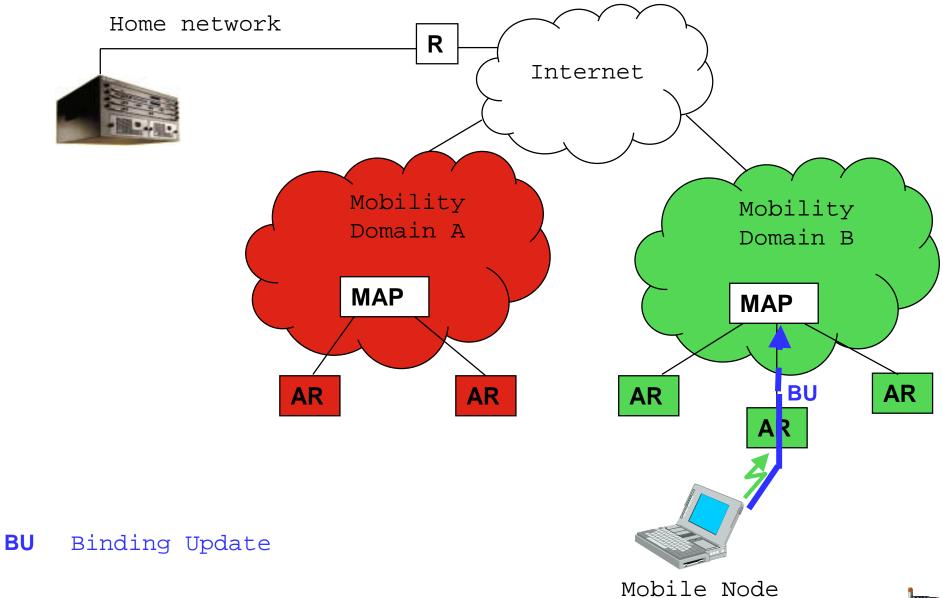


Hierarchical Mobile IPv6



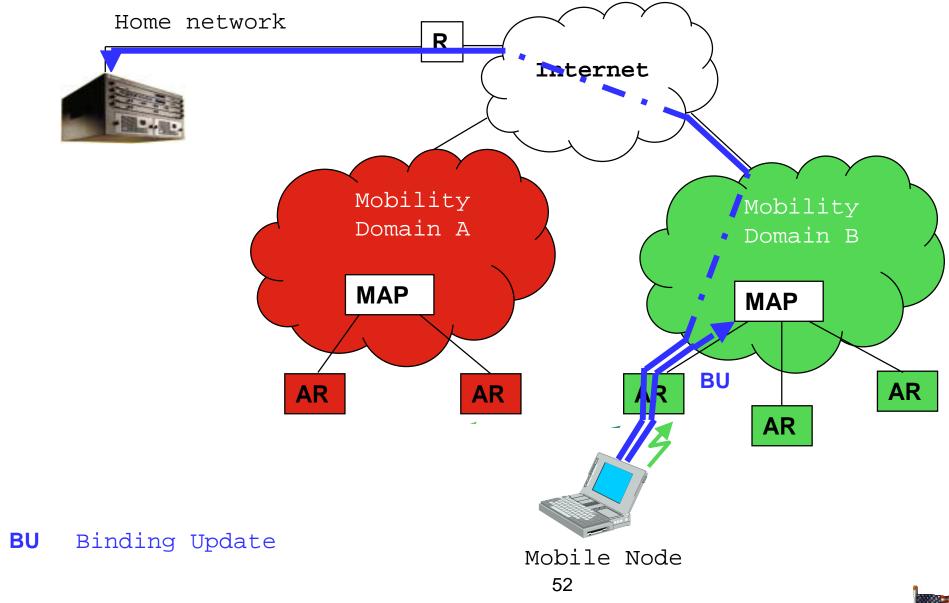


Example 1: Mobility within Domain





Example 2: Mobility between Domains







Summary

Both "sides", Internet and Cellular Communication, have recognized the promising potential of the Mobile Internet market

Co-operation between organizations of the Internet and Cellular Communication side are established

IPv6 and Mobile IPv6 are seen as an efficient and scalable solution for the future Mobile Internet

Numerous research activities take place in the area of IPv6 for mobile users

From the technical side not all problems are solved now - but we are doing a good job here



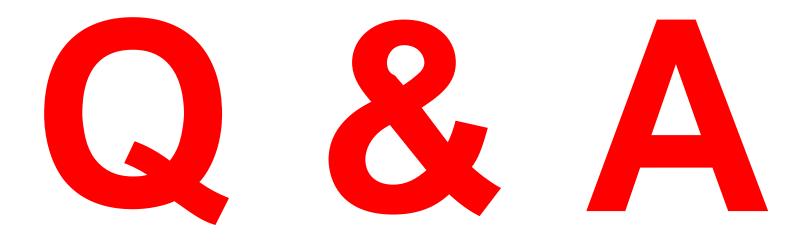


Diversity of today's available mobile devices













Thanks for your attention!



