

A New Method for Intrusion Detection on Hierarchical Wireless Sensor Networks

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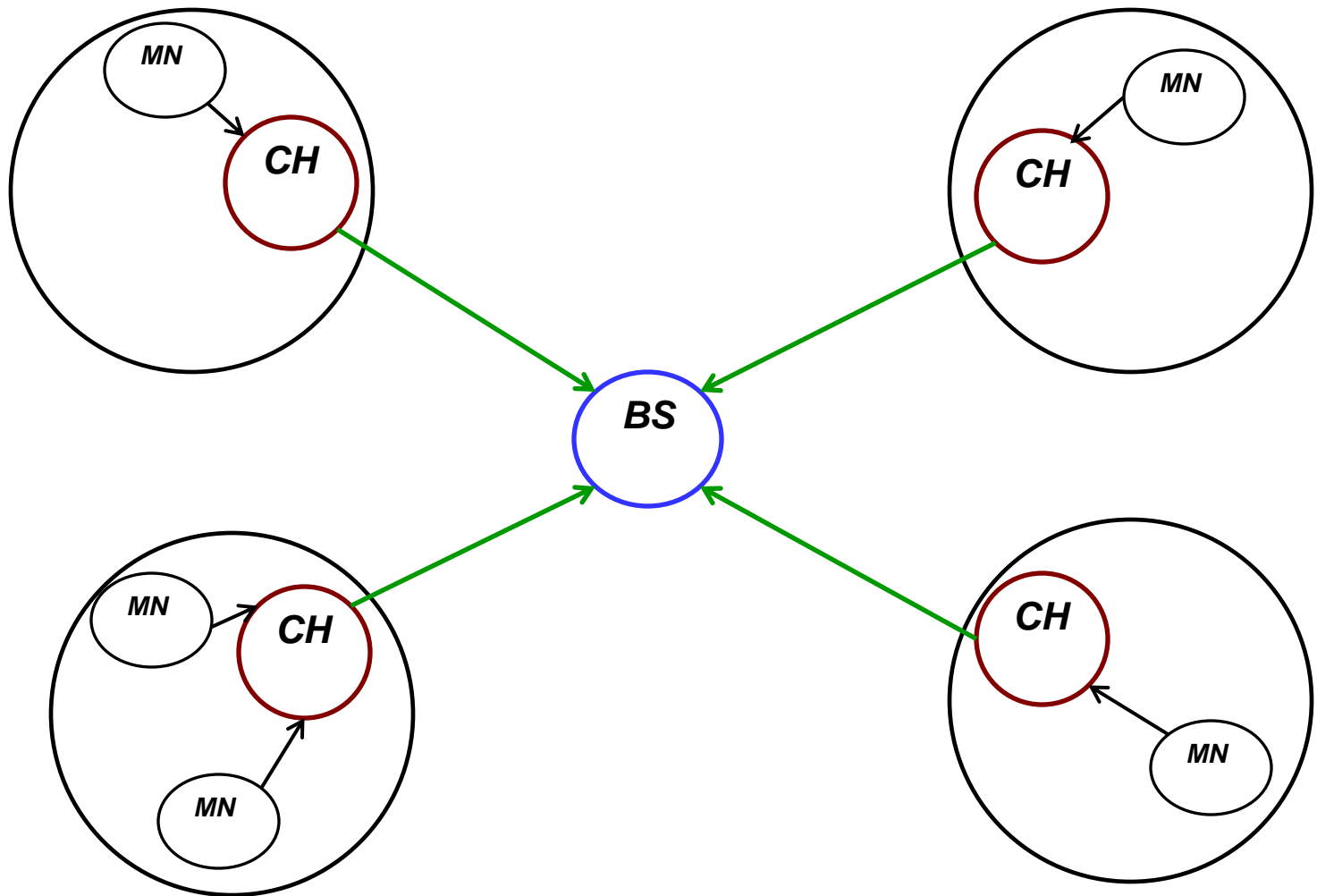
Outline

- *Introduction*
- **Intrusion Detection Systems (IDS) for Wireless Sensor Networks**
- **Collaborated-Based Intrusion Detection**
- **Routing Table Intrusion Detection**
- **Isolated Table Intrusion Detection System**
- **Comparison Experiments**
- **Conclusions/Criticisms**

Introduction

- Energy conservation is the critical performance consideration to extend WSN lifetime.
- This paper adopts **cluster-based WSN (CWSN)** as their choice to extend network coverage and increase lifetime {They reference the Leach paper [Heinzelman]}.

Cluster Architecture





Cluster Hierarchy

- **MN** (Member Nodes) deliver sensed data to the **BS** (Base Station) through their **CH** (Cluster Head).
- **CHs** are chosen as center of cluster and data from **MN** is aggregated before being sent to **BS**.
- {Note - paper ignores dynamically changing CHs.}

WSN Security

- Authors divide WSN security roles into IDS (Intrusions Detection Systems) and IPS (Intrusion Protection System).
- IPS uses **authentication** and **symmetric keys** to defend system from outside attackers (**Not part of this paper**).
- IDS identifies attackers using a rules database via anomaly data (e.g., signatures).



WSN IDS

- This paper uses term '**anomalies**' and relies on detecting them.
- With WSN, IDS database must be smaller due to limited resources.
- Attack behavior is different in WSN and draining sensor energy and breaking network connectivity are possible attack strategies.



Stated Objectives

1. Use IDS to **isolate** intruders in WSN.
2. Reduce IDS energy consumption to extend WSN lifetime.
3. Find balance between energy consumption and WSN security.

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IDSs for Wireless Sensor Networks

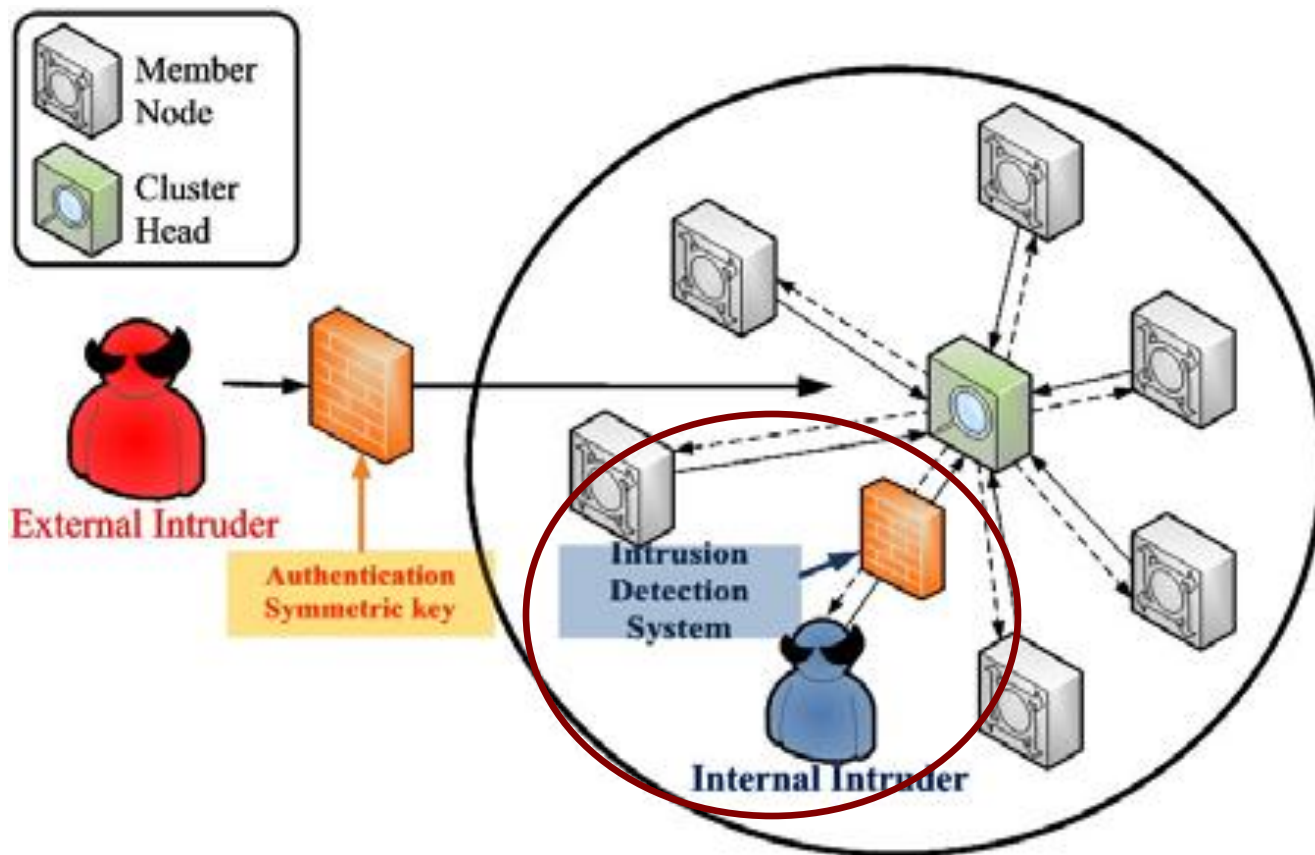


Figure 1. The Secondary Defense of WSNs



Stages of WSN Attack Behaviors

1. Preparation stage

- Probing the communication of the **BS** and gathering info from the WSN.
- Intrusion often happens in the connection between the **BS** and **CHs**.

2. Attack and Occupy stage

- Uses victim components (spooft, alter or replay routing info and selective forwarding).

Stages of WSN Attack Behaviors

3. Doom Stage

- This paper focuses on these forms of attack to crash WSN:
 - including Hello Flooding, Denial of Service, Denial of Sleep, Sinkholes and Wormholes.
- Goal is to attack **CH** and need to occupy a sensor (an **MN**) to intrude **CH**. Essentially, try to consume **CH** resources until **CH** energy exhausted.

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Collaboration-based Intrusion Detection (CID)

- **CH** runs whole network: controls and monitors **MNs** and communicates with **BS**.
- Uses several security levels. Level is determined by **thresholds**.
- **MNs** divide into Monitor Groups (**MGs**) to reduce energy and to monitor **CH** using authentication methods.
- Structure and alarm thresholds set by an **administrator**.



Collaboration-based Intrusion Detection (CID)

- CH detects anomaly MNs and isolates them.
- When MNs raise alarms above threshold, they can depose CH.
- CID weaknesses when CH changes:
 - when CH changes, new CH does not obtain old isolation information.
 - new CH consumes extra energy monitoring MNs again for malicious nodes.

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Routing Table Intrusion Detection (RTID)

- WSN routing table used to detect anomaly behaviors.
- Example in Figure 3, static sensors topology uses a **spanning tree** with hops metric in the IASN table.
- IASN (Information Authentication for Sensor Networks) Table holds information such that arriving sensor data is compared against valid info coming correct path.



Routing Table Intrusion Detection (RTID)

- Both information sent along the wrong path or invalid information deemed **anomalous** by IASN table.
- Actual routing is **DSDV** (Destination-Sequenced Distance Vector).
- Problem is each sensor must store DSDV routing table and IASN table.

Outline

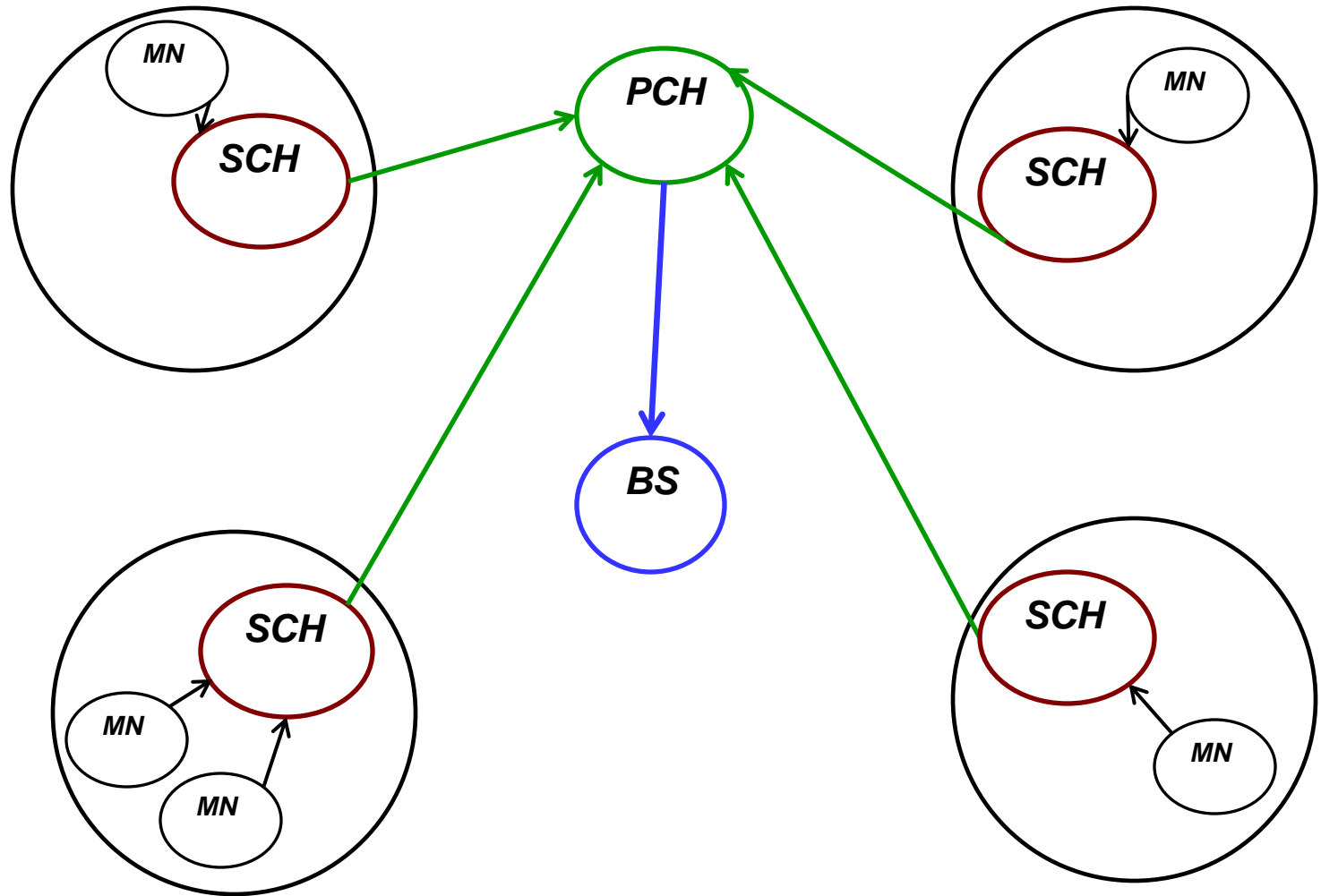
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Isolation Table Intrusion Detection System (ITIDS)

- Authors claim: to save on energy consumption of sensors employ an **isolation table** for IDS.
- Node types: 1 **BS**, 1 Primary Cluster Head (**PCH**), 1 Secondary Cluster Head (**SCH**) per **MG** where **MNs** are assigned to **MGs**.

Modified Cluster Architecture



ITIDS Architecture

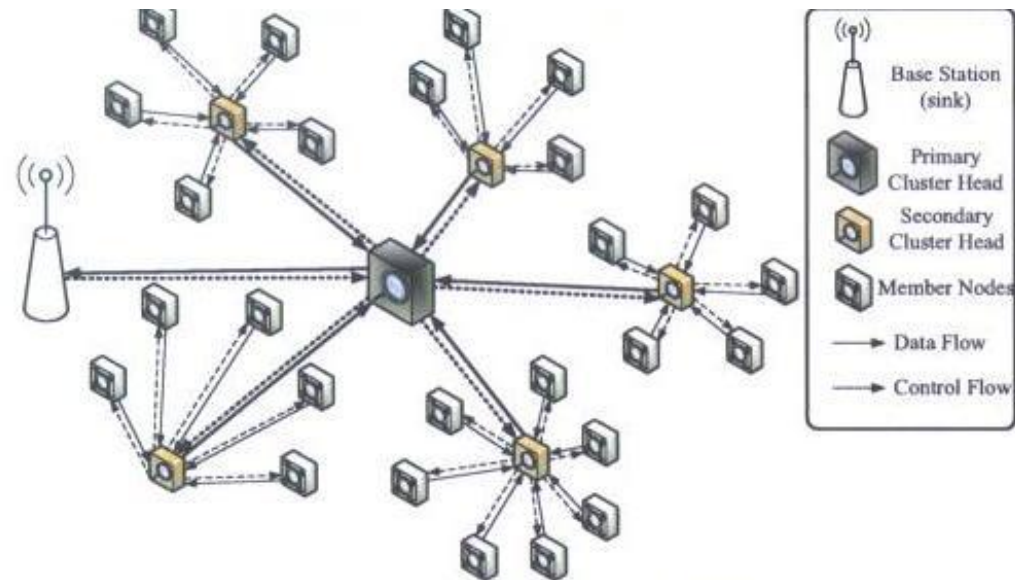


Figure 5. The Architecture of ITIDS.

Table 4. The Routing Table of SCH

N_{id}	G_{id}	N_{INFO}	$E_{fl}(\mu J)$
S406 1	1	INFO ₁	$7.2 \cdot 10^6$
S128 2	2	INFO ₁	$7 \cdot 10^6$
S564 9	3	INFO ₃	$7.1 \cdot 10^6$
S710 7	2	INFO ₂	$6.5 \cdot 10^6$
S392 5	5	INFO ₅	$6.8 \cdot 10^6$

ITIDS Sensor Roles

- **BS** - used by administrator to control WSN; receives sensing data and isolation tables.
- **PCH** - gathers sensing data and isolation table from **SCH** to **BS**.
- **SCHs** - calculate trust value to find malicious **MNs** and monitors **PCH** with **MNs** in its **MG**.
- **MNs** - send sensed data to **SCH** and rotate monitoring **PCH**.



ITIDS Four Stages

1. System Predefinition of IDS
2. SCHs monitors MNs.
3. SCH and MNs monitor PCH.
4. IDS backups the isolation table in BS.

Predefinition Stage

- Set up and define all the roles.
 - Set sensing types and number of **MGs**.
 - **PCH** randomly selects a sensor node in each **MG** to be **SCH**.
 - * Authors discuss duty-cycle between **PCH** and **SCH** {unclear!}.
- Anomaly thresholds set per **MG** {set to $2/3$ of N_k }.



Predefinition Stage

- **MNs** report info (including remaining energy) to **SCHs**.
- **SCHs** authenticate info from their **Mns** and isolate anomalies.
- Anomalies recorded in the **SCH** isolation table.
- **SCH** isolation tables integrated by **PCH**.



Stage 2 -SCHs Monitors MNs

- SCHs authenticate info from MNs.
- Anomalies recorded in the SCH isolation table (transmitted immediately to PCH).
- Possible MN anomalies:
 - Routing info changed by intruder.
 - Record attack behaviors (spoofed, altered, replayed routing info or selective forwarding) in isolation table.

Stage 2 -SCHs Monitors MNs

- Possible **MN** anomalies (cont.):
 - Remaining energy has increased instead of decreasing (implies sinkhole attack).
- Defense against doom attacks:
 - Measure frequency of **MN** traffic to **SCH** during time slot. If **too frequent**, **MN** is isolated by **SCH**.
 - Routing tables record energy information, if energy has increased **SCH** isolates the MN.

Stage 3 - SCH and MNs monitor PCH

- When enough **MNs** raise alarm total above threshold, **SCH** deposes **PCH**.
- New **PCH** integrates isolation table from each **SCH** and sends **BS** latest isolation table.
- **PCH** selects new **SCHs** - one from each **MG** randomly.
- Duty cycle of **PCH** divided into **SCH** duty cycles equally.

Table 5 - Isolation Table

- **Sample attack behaviors:**
 - **Fault information::** deliver info is a fault type from routing table.
 - **Detection error::** **MN** raises detection alarm when no attack occurs.
 - **Redundancy::** the same **MN** repeatedly sends data to **SCH**.
 - **Wrong source::** the data source of the transmission is wrong.

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Comparison Experiments

- Use NS-2 to compare ITIDS against CID and RTID.
- Only comparisons are Number of Alive Nodes for ITIDS and CID and Transmission Accuracy between ITIDS and RTID.
- Not enough detail to understand details of wireless energy usage.
- No discussion of MAC protocol simulated (assume 802.11).



Comparison Experiments

- **No indication of radio (byte or packet)**
 - 100 or 200 sensors in 10,000 sq. meters.
 - **PCH** in center, **BS** ??
 - 50 meter transmit radius
 - Energy formula: **ONLY** transmission + detection energy (calculation).
- **No mention of routing or multi-hop.**
- **ONLY** simulated doom attacks.



Comparison Experiments

- Remained resources formula used to determine number of alive nodes.
- Again, **CHs NOT** rotated to balance energy among **MNs**.

Figure 6 Number of Alive Nodes

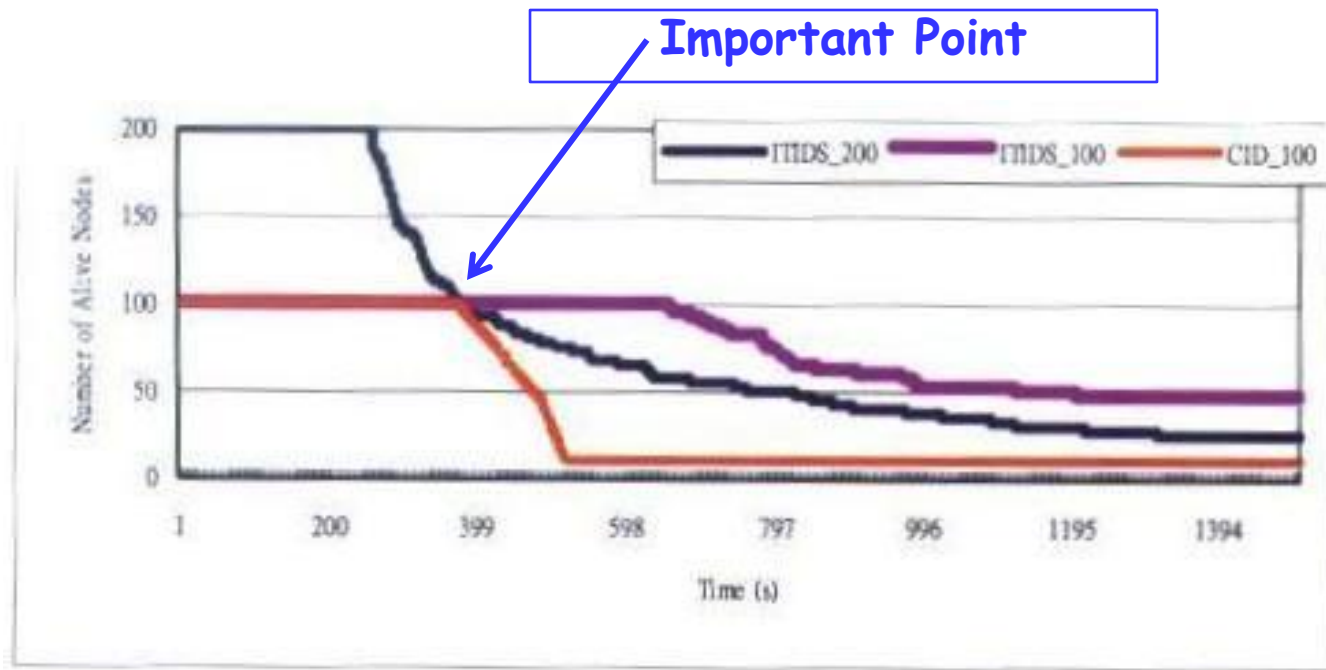


Figure 6. The Comparison of the Number of Alive Nodes Between ITIDS and CID.

Figure 7 Transmission Accuracy

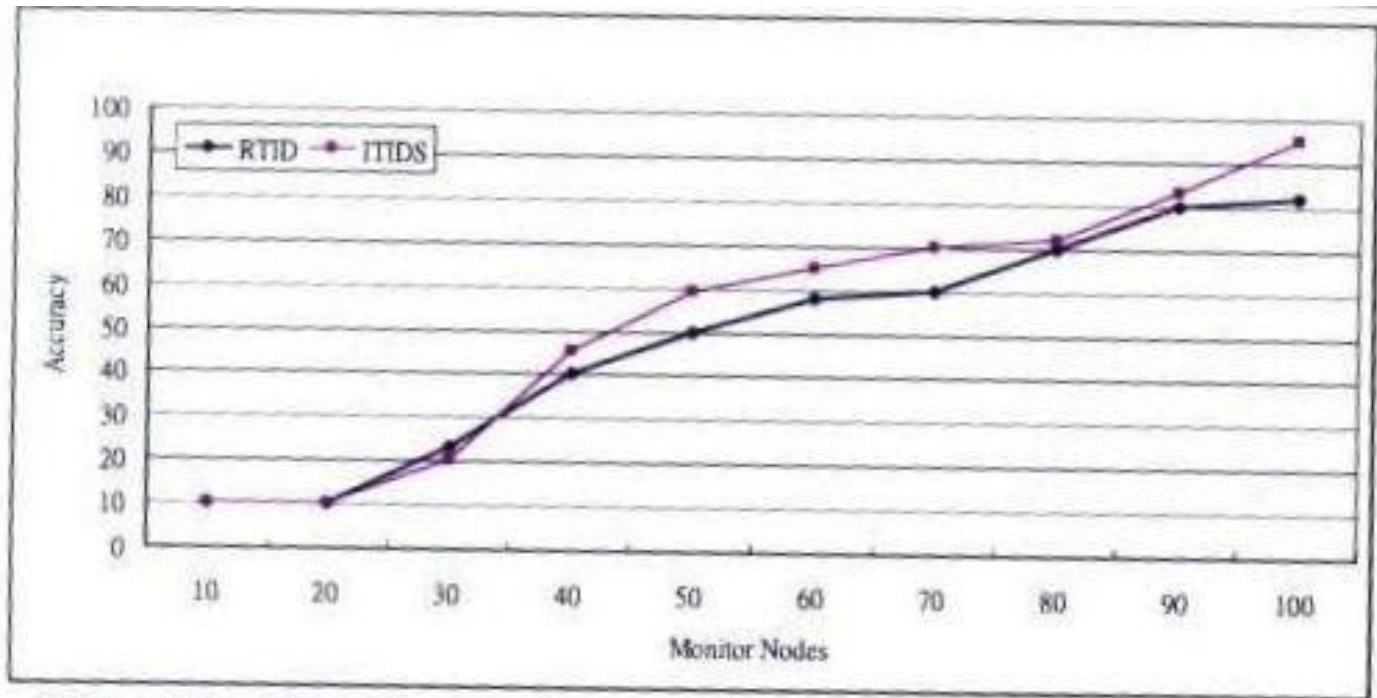



Figure 7. The Comparison of the Transmission Accuracy Between ITIDS and RTID.



Conclusions/Criticisms

- Authors conclude that evidence shows ITIDS can prevent attacks effectively.
 - No info on number of attacks attempted and number prevented.
- Generally, experiments are few weak with too few details to assess or to reproduce.
- Difficult to see cause-and-effect from two graphs presented.



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Questions ??

Thank You!

