SPEC CPU2000
Measuring CPU Performance in the New Millennium

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IEEE Computer Magazine
2000

Introduction

- Computers become more powerful
- Human nature to want biggest and baddest
  - But how do you know if it is?
- Even if your computer only crunches numbers (no I/O), it is not just CPU
  - Also cache, memory, compilers
- And different software applications have different requirements
- And whom do you trust to provide reliable performance information?

Standard Performance Evaluation Corporation (SPEC)

- Non-profit consortium
  - Hardware vendors, software vendors, universities…
- High-perf numeric computing, Web servers, graphical sub-systems
- Benchmark suites derived from real-world apps
- Agree to run and report results as specified by benchmark suite

The SPEC consortium's mission is to develop technically credible and objective benchmarks so that both computer designers and purchasers can make decisions on the basis of realistic workloads.

- June 2000, retired CPU95. Replaced with CPU2000.
  - 19 new applications
- How does SPEC do it?
- One-specific day in SPEC release

Outline

- Introduction
- SPEC Benchathon
- Benchmark candidate selection
- Benchmark results
- Summary

SPEC Benchathon

- 6am, a Thursday, Feb 1999
- Compaq employee (author?) comes to work, finds alarm off
  - IBM employees still there from the night before
  - Sub-committee in town for a week-long "benchathon"
- Goes to back room, 85 degrees thanks to workstations Sun, HP Siemens, Intel SGI, Compaq and IBM
- Looks at results of Kit 60 (becomes SPEC CPU2000 10 months later)

Portability Challenge

- Primary goal at this stage is portability
  - 18 platforms from 7 hardware vendors
  - 11 versions of Unix (3 Linux) and 2 Windows NT
  - 34 candidate benchmarks,
  - → but only 19 successful on all platforms
- Challenges can be categorized by source code language
  - Fortran
  - C and C++

Portability Challenges – Fortran (1 of 2)

- Fortran 77 – easiest to port since relatively few machine-dependent features
- But still issues.
  - Ex- 47,134 lines of code, 123 files and hard-to-debug wrong answer when optimization enabled for one compiler
  - Later, determine compiler is to blame and benchmark ships (`200.sixtrack`)
- Several F77 compilers allocate 200 MB memory
  - When static, takes too much disk space
  - When dynamic, another vendor has stack limits exceeded
  - SPEC later decides dynamic but vendor can choose static if needed

Portability Challenges – Fortran (2 of 2)

- Fortran-99 more difficult to port since F90 compilers less common
  - "Language Lawyer" wants to use
  - One platform with F90 has only 3 applications working
  - Later, works on all but does reveal bugs in current compilers
  - And causes change in comparable work category

(sidebar next)

Comparable Work (Sidebar 1 of 3)

- Want comparable work across platforms. But difficult. Consider `187.facerec`:

```
If ((NewSim – OldSim) > SimThresh) Then
  CoordX (IX, IY) = NewX
  CoordY (IX, IY) = NewY
  Hops = Hops + 1
  Improved = .TRUE.
EndIf
Sweeps = Sweeps + 1
  If ((.NOT. Improved) .OR.
    (Sweeps >= Params%Match%MaxSweeps))
Exit
```

- Algorithm to look through images
- Attempt to recognize a face

Comparable Work (Sidebar 2 of 3)

- The loop exit depends on floating-point comparison. That depends upon accuracy of flops, as implemented by vendors
- If two systems recognize a face but take different iterations, is that the same work?
  - Could argue same work, different path
  - But SPEC wants mostly the same path
  - And don't want to change spirit of algorithm with fixed number of iterations

Comparable Work (Sidebar 3 of 3)

- Solution? File-by-file validation tolerances
- Modify `187.facerec` to get number of iterations, and summary of total iterations

	Detail	Summary
reltol	0.2	0.001
abstol	5	2.e – 7
skiptol	4	0

- Valid if iterations within 20% (reltol=0.2) or no more than 5 different (abstol=5).
  - Allowed to fail 4 times (skiptol=4)
- For overall run, iterations within .1% (reltol = 0.001) and all iterations checked (?) (skiptol = 0)
- So, two platforms may do different amounts of work on 1 face, but similar on many faces

Portability Challenges – C and C++

- C has more hardware-specific issues
  - How big is a `long`? A pointer? Does a platform have `calloc()`? Little endian or big endian byte order?
- Do not want configure scripts because wants to minimize source code differences
  - Instead, prefers `#ifdef` directives to manually control
- C++ harder (standard was new)
  - Only 2 C++ candidates, and 1 too hard to make ANSI
  - Ultimately, only 1 ships (`252.eon`)

February 1999 Benchathon Results

Table 1. February 1999 benchathon results.

	19 Feb	26 Feb
Compile errors	22	2
Runtime errors	18	6
Validation errors	60	41
Total	100	49

- Goal of benchathon is to have project leaders in place to resolve technical issues from multiple stakeholders
  - Employees from different companies, helping each other debug

Project Leader Structure

- Project leader shepherds candidate benchmarks
  - "Owns" resolution of portability problems
  - One has 10, but later lightens load
  - One has only 3, but difficult challenges
- Example 1: simulator gets different answers on different platforms
  - Later dropped
- Example 2: another requires 64 bit integers. Compilers for 32-bit platform can specify
- Example 3: app constructs color pixmap. Subtle differences in shades. Since not detectable by eye, deemed ok.

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Benchmark Selection (1 of 3)

- Porting is clearly technical. Answer question "does benchmark work?"
- Selecting benchmarks harder
- Solicit candidates through search process on Web
- Members of SPEC vote. "Yes" if:
  - Many users
  - Exercises significant hardware resources
  - Solves interesting technical problem
  - Published results in journal
  - Or adds variety to suite

Benchmark Selection (2 of 3)

- "No" if:
  - Too hard to port
  - Does too much I/O so not CPU bound
  - Was previously in SPEC CPU suite
  - Code fragment rather than complete application
  - Is redundant
  - Appears to do different work on different platforms

Benchmark Selection (3 of 3)

Benchmark	Language	KLOC	Resident size (Mbytes)	Virtual size (Mbytes)	Description
SPECint2000					
164.gzip	C	7.6	181	200	Compression
175.vpr	C	13.6	50	55.2	FPGA circuit placement and routing
176.gcc	C	193.0	155	158	C programming language compiler
181.mcf	C	1.9	190	192	Combinatorial optimization
186.crafty	C	20.7	2.1	4.2	Game playing: Chess
197.parser	C	10.3	37	62.5	Word processing
252.eon	C++	34.2	0.7	3.3	Computer visualization
253.perlbmk	C	79.2	146	159	Perl programming language
254.gap	C	62.5	193	196	Group theory, interpreter
255.vortex	C	54.3	72	81	Object-oriented database
256.bzip2	C	3.9	185	200	Compression
300.twolf	C	19.2	1.9	4.1	Place and route simulator
SPECfp2000					
168.wupwise	F77	1.8	176	177	Physics: Quantum chromodynamics
171.swim	F77	0.4	191	192	Shallow water modeling
172.mgrid	F77	0.5	56	56.7	Multigrid solver: 3D potential field
173.applu	F77	7.9	181	191	Partial differential equations
177.mesa	C	81.8	9.5	24.7	3D graphics library
178.galgel	F90	14.1	63	155	Computational fluid dynamics
179.art	C	1.2	3.7	5.9	Image recognition/neural networks
183.equake	C	1.2	49	51.1	Seismic wave propagation simulation
187.facerec	F90	2.4	16	18.5	Image processing: Face recognition
188.ammp	C	12.9	26	30	Computational chemistry
189.lucas	F90	2.8	142	143	Number theory/primality testing
191.fma3d	F90	59.8	103	105	Finite-element crash simulation
200.sixtrack	F77	47.1	26	59.8	Nuclear physics accelerator design
301.apsi	F77	6.4	191	192	Meteorology: Pollutant distribution

Objective Criteria

- Want objective technical reasons for choosing/not choosing benchmarks
- But often at odds since technical reasons may be confidential
- Solution was all members provided some objective data and kept confidential
- Info: I/O, cache and main memory behavior, floating-point op mixes, branches, etc.

Subjective Criteria (1 of 3)

- Confidence in benchmark maintainability
  - Some have errors that are difficult to diagnose
  - Some have error fixed then re-appears
  - Some have easy to fix errors, but take sub-committee time
- All contribute to confidence level
- Needs to be manageable

Subjective Criteria (2 of 3)

- If stable quickly enough then can be analyzed
- Can be complex, but should not be misleading
- Workload should be describable in ordinary English and technical language

Subjective Criteria (3 of 3)

- Vendor interest matters. Temptation to vote accordingly. Two factors reduce influence
  - Generally, do not know numbers on competitors hardware. Hardware may not even be released. So, hard to vote for a benchmark because it is bad. Better to just vote on *merit*
  - Hard to argue the converse. I.e.- "you should vote for `999.favorite` because it helps my company".
- Of course, vendor interest represented. Want to keep level playing field

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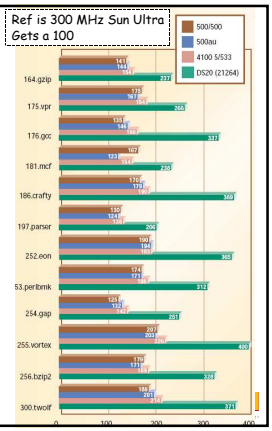
Results

- Want something we cannot learn by looking at clock speed
- 21164 differ by 92.3 to 331 (non-green)
- 500 MHz differ by more than 5% on 18 of 26 (clock diff is 7%)
- 533 MHz wins by 10% 3 times, < 3 % 3 times and loses 3 times
- Difference in memory hierarchy?

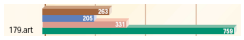


Memory Hierarchy Differences

- Memory differences not clear
  - 500/500 has largest L3 cache
  - 500au has best memory latency
  - 4100 533 has highest main memory bwidth

CPU	Alpha 21164		
System	500/500	500au	4100 5/533
CPU MHz	500	500	**533**
L1 cache on chip	8 Kbytes (instruction) + 8 Kbytes (data)		
L2 cache on chip	96 Kbytes		
Off-chip cache			
Size (Mbytes)	**8**	2	4
Latency (ns)	82	**58**	62
Latency (processor cycles)	41	**29**	33
Main memory			
Latency (ns)	341	**247**	248
Latency (processor cycles)	171	**124**	132
Bandwidth (Mbytes/s)	200	238	**272**

- 533 wins largest on `179.art`.



- Why? Perhaps because the benchmark fits in 4 Mbyte cache, but not 2 Mbyte cache

Effects of Cache

- 500/500 outperforms 533 most on `181.mcf`
- Looking at profile, can see benefits from larger cache
- `252.eon` is only place 500au wins. Very small, maybe in range of validation test. But could be from lower cache latency

Effects of Main Memory

- Cache matters, but many apps depend upon main memory
  - Where 533 is best

Benchmark	2 Mbytes	4 Mbytes	8 Mbytes
179.art	29.1	0.5	0.4
171.swim	23.9	23.7	23.6
183.equake	23.6	22.2	21.0
189.lucas	19.6	19.3	18.9
173.applu	14.2	14.0	13.0

- 2 of top 5 generators do not get better when same system has bigger cache
- 533 versus 500au has 14% better mem bandwidth. Provides 1%, 12%, 7% improvement for `171.swim`, `189.lucas`, `173.applu`

189.Lucas Analysis

- 21164 can have only two outstanding mem requests. Stalls after third
  - So code that spreads out memory requests will work better than if bunched
- 189.Lucas was hand-unrolled before submitting to SPEC and spread memory references
- So, overall, if you want good performance, benchmarks show not just CPU speed

Processor Performance (1 of 3)

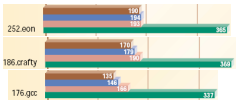
- DS20 has new chip 21264, but still 500 MHz
- Memory system different

CPU	Alpha 21164			Alpha 21264
System	500/500	500au	4100 5/533	DS20
CPU MHz	500	500	**533**	500
L1 cache on chip	8 Kbytes (instruction) + 8 Kbytes (data)			64 Kbytes (instruction) + 64
L2 cache on chip	96 Kbytes			None
Off-chip cache				
Size (Mbytes)	**8**	2	4	4
Latency (ns)	82	**58**	62	32
Latency (processor cycles)	41	**29**	33	16
Main memory				
Latency (ns)	341	**247**	248	184
Latency (processor cycles)	171	**124**	132	92
Bandwidth (Mbytes/s)	200	238	**272**	1,232

- Cache latency by factor 1.8x, memory latency by 1.3x,
- Bwdith latency by 4.5x

Processor Performance (2 of 3)

- Three integer benchmarks biggest



- `176.gcc` greatest because of workload
  - Was only 1.61x better for CPU95
- CPU95 gcc ran for 79 seconds, 47 MB vm
- CPU2000 gcc runs 327 seconds and uses 156 MB vm

Processor Performance (3 of 3)

- `171.swim`, `189.lucas`, `173.applu` (earlier table) should benefit from 4.5x mem bandwidth
  - 4.9x, 2.2x, and 2.8x respectively
- `183.equake` improves by only 1.7x despite high miss rate. Analysis shows because program is bound by latency not memory bwidth.

Compiler Effects

- All results in article use single compiler and "base" tuning
  - No more than 4 switches and same switches for all benchmarks in a suite
- Different tuning would have different results
- Highlights:
  - 400,000 lines of new float code with '-fast' flag make it tougher to be robust
  - Unrolling can really help. Ex: `178.galgel` unrolled had 70% improvement versus base tuning
- Note, recommend continued compiler improvements but should improve general applications and not just SPEC benchmarks

Summary

- SPEC encourages industry and academia to study more
- Now is the time for CPU 200x (CPU 2004)
- (Have ordered CPU2000 for those that want a go)

SPEC/Academia

SPEC encourages research and academic usage of the new CPU2000 suite, as described in section 4.5 of the run rules (http://www.spec.org/cpu2000/docs/runrules.txt). During the year 2000, SPEC is offering two incentives to the academic community:

- SPEC has reduced the cost of membership for associate members, as described at http://www.spec.org/news/y2kspecial.html.
- Universities can obtain a free copy of SPEC CPU2000 (or certain other SPEC products) through 31 December 2000, by following the instructions at the above URL.