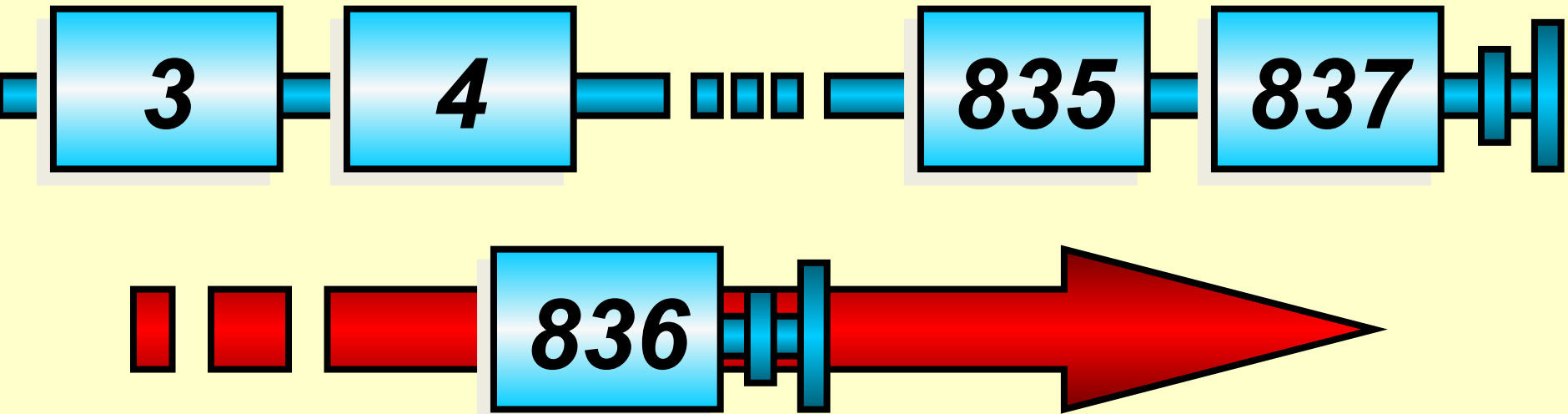


Hashing

Desire

- We want to store objects in some structure and be able to retrieve them extremely quickly.
- The number of items to store might be big.

Hashing--Why?



Steps

15

10

5

0

5

10

15

20

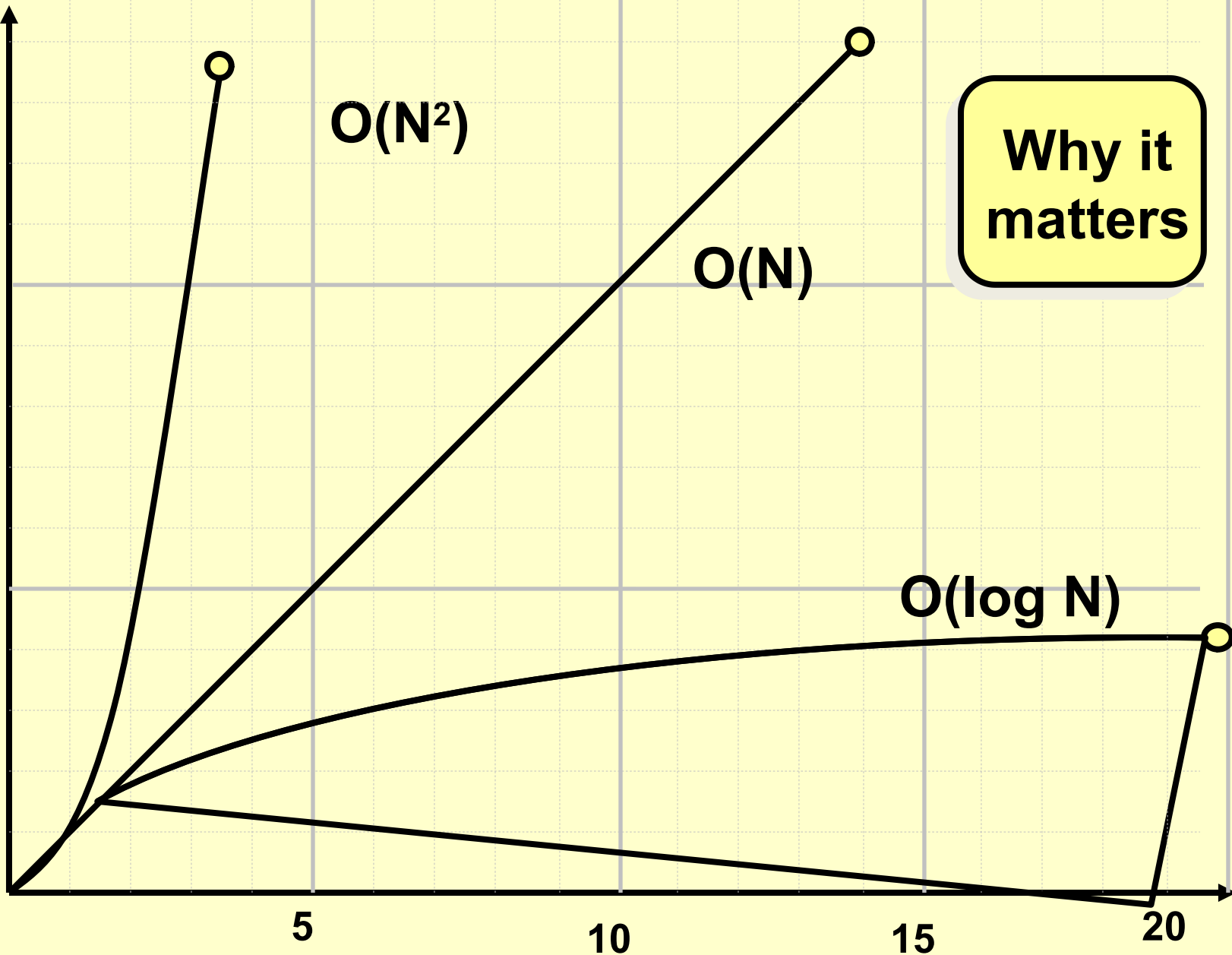
Items

$O(N^2)$

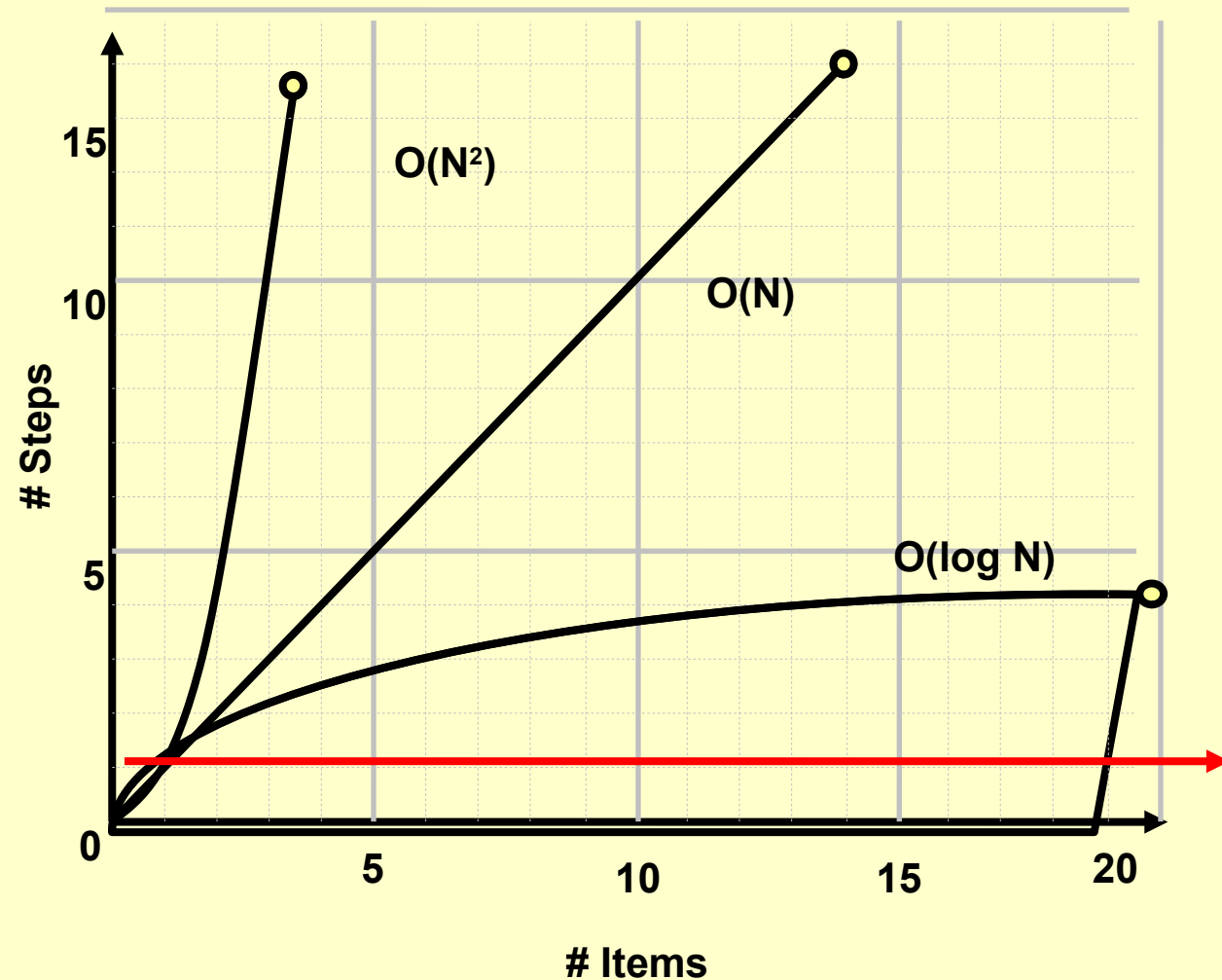
$O(N)$

$O(\log N)$

**Why it
matters**

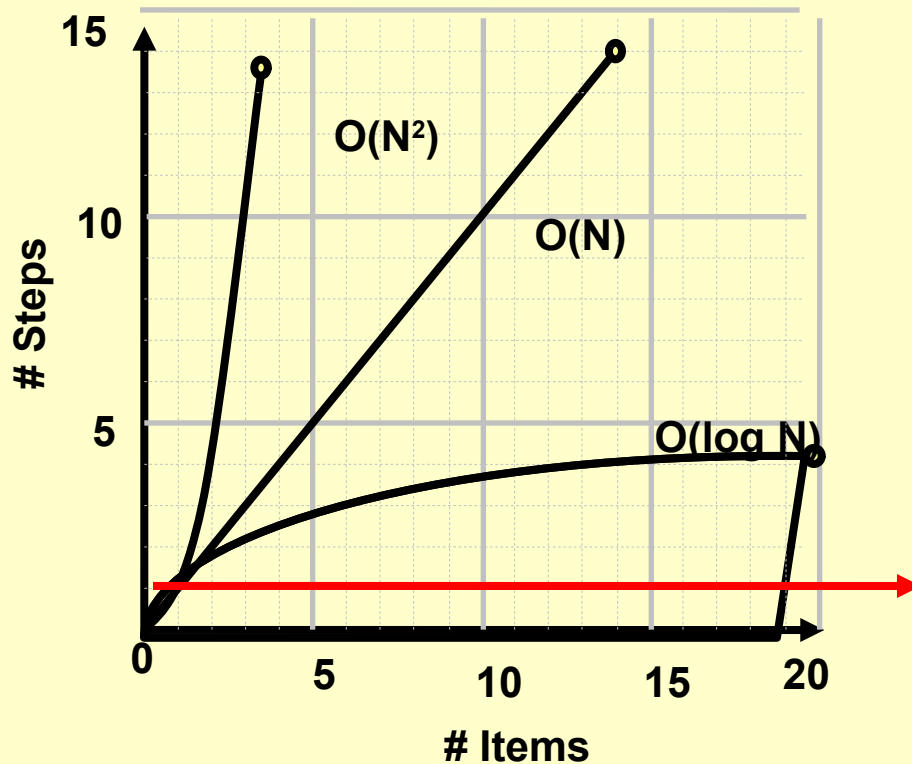


Big Uh Oh



There's a way
to reduce
access time
down to $O(1)$,
or nearly so

Sanity Check



*A search time
of $O(1)$?
How is this
possible?*

Introducing....

Hashing!

Naive Solution:

```
Data[ ] myRecord = new Data[ 9999999999 ];
```

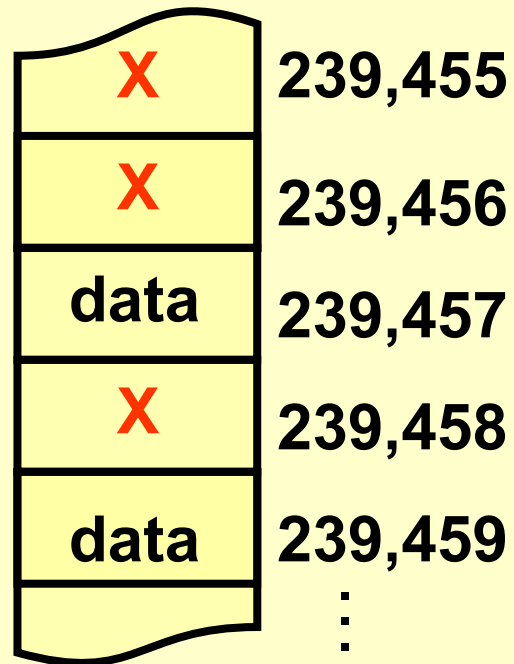
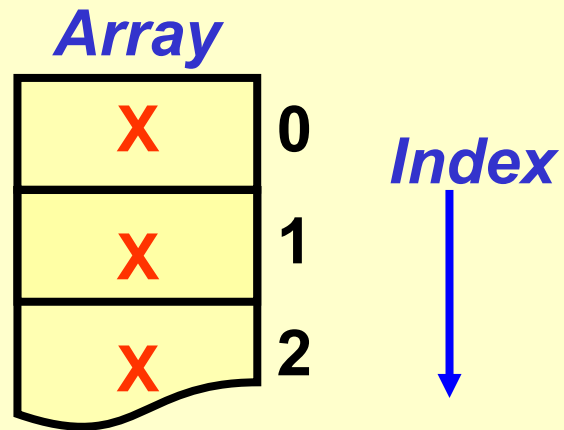
```
/*
```

```
* NOTE: Here, we assume there are  
* approximately a billion social security  
* numbers
```

```
*/
```

Perhaps not the best?

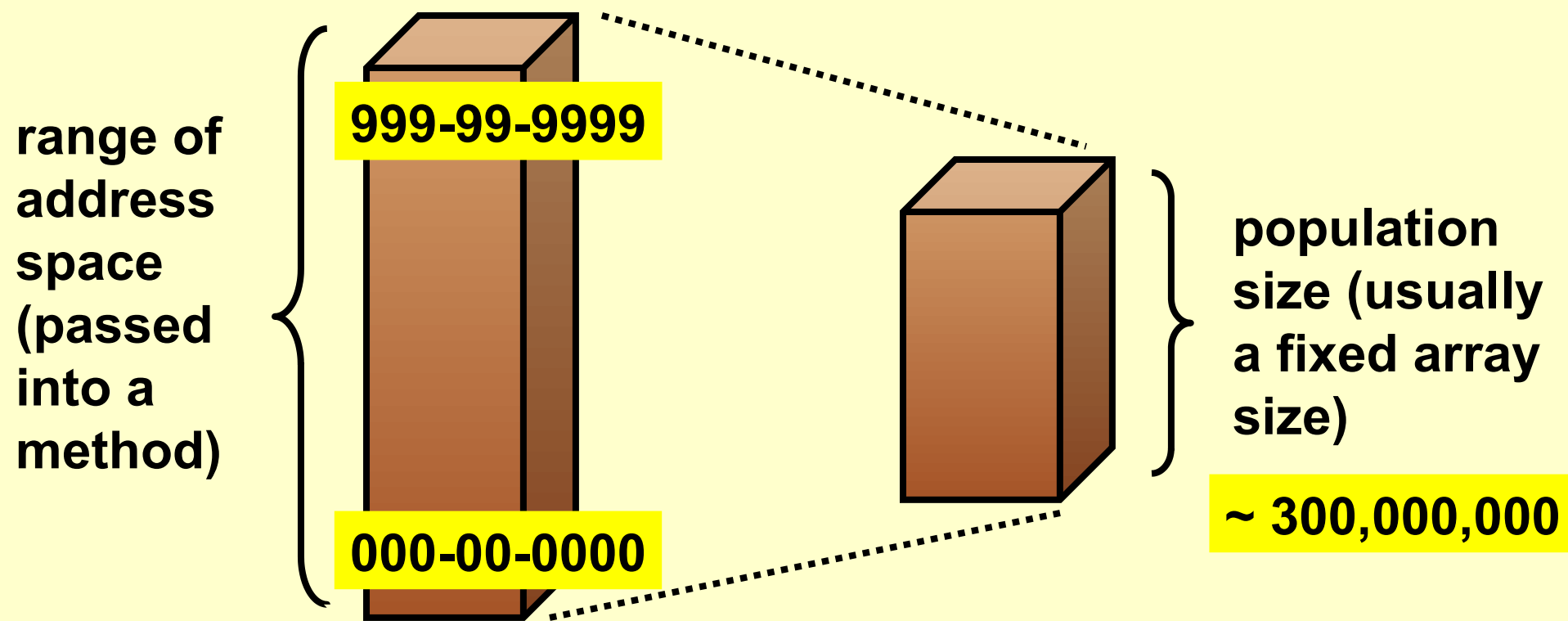
Example



Hashing

Idea:

Shrink the *address space* to fit the *population size*.



```
class AllContacts {  
    // storing contacts in a hashmap from names to contacts  
    HashMap<String, Contact> contactMap  
        = new HashMap<String, Contact>();  
  
    AllContacts() {}  
  
    // return contact that goes with a given name  
    // will return null if the name isn't in the map  
    Contact findContactByName(String name) {  
        return contactMap.get(name);  
    }  
  
    // takes the whole contact as input  
    AllContacts addContact(Contact aContact) {  
        Contact oldContact  
            = contactMap.put(aContact.name, aContact);  
        return this;  
    }  
}
```

```
class AllContacts {  
    // storing contacts in a hashmap from names to contacts  
    HashMap<String, Contact> contactMap  
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    // will return null if the name isn't in the map  
    Contact findContactByName(String name) {  
        return contactMap.get(name);  
    }  
  
    // takes the whole contact as input  
    AllContacts addContact(Contact aContact) {  
        Contact oldContact  
            = contactMap.put(aContact.name, aContact);  
        return this;  
    }  
}
```

Reality Check

Everyone getting the idea?

Hash Function Example

Instead of

```
StudentFile temp = studentRecords[iSocSecNum] ;
```

reduce the range:

```
StudentFile temp =  
    record[iSocSecNum % record.length] ;
```



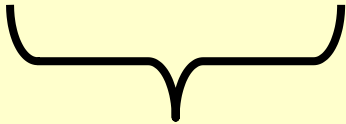
returns an index within the
appropriate range

Recall

- Our friend the Mod Function %

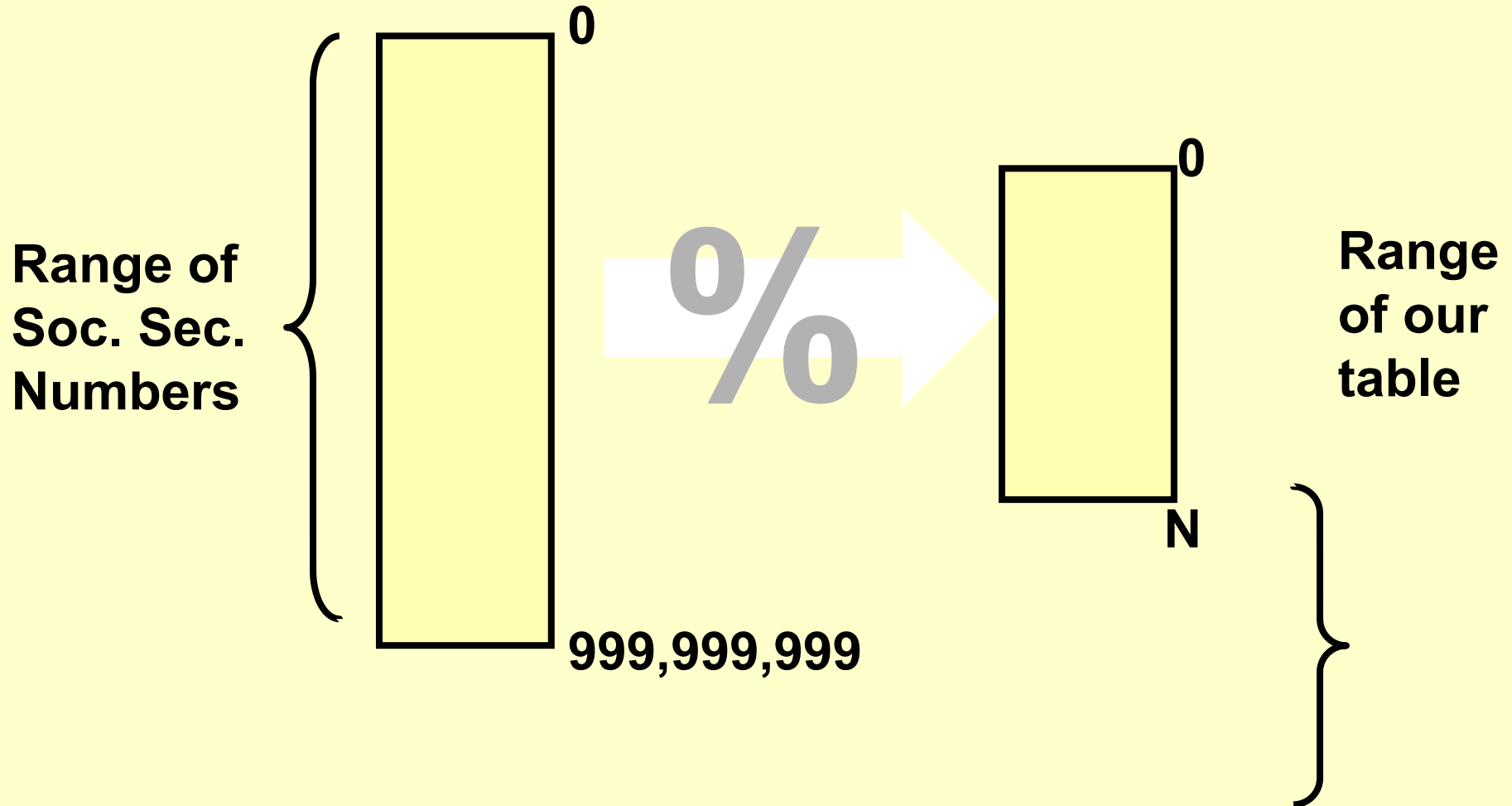
$$x \% y$$

- will yield values between 0 and $y - 1$



***Note remarkable similarity
to range of values for the
index of an array of size y***

The Art of Hashing



A problem...

- We have an array of length 100
- We have about 50 students
- We hash using: `ssn % 100`

- George P. Burdell

- 123-45-6789

Collision!

- George W. Bush

- 321-54-7689

Designing Hash Functions

The Perfect Hash Function:

- Fast
- No collisions

Common Hash Functions:

- ***Digit selection***: e.g., last 4 digits of phone number
- ***Division***: modulo
- ***Character keys***: use ASCII num values for chars (e.g., 'R' is 82)

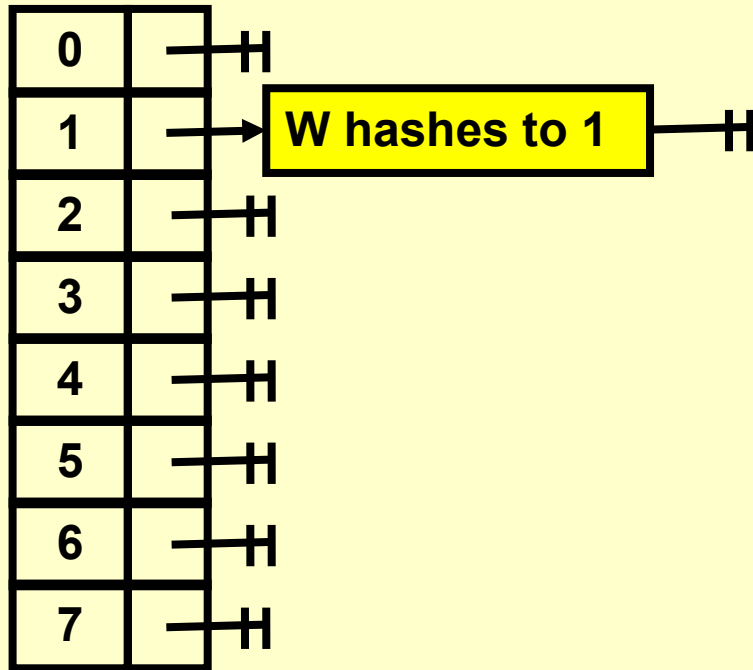
Collision Resolution

Technique: External chaining:

0	—	H
1	—	H
2	—	H
3	—	H
4	—	H
5	—	H
6	—	H
7	—	H

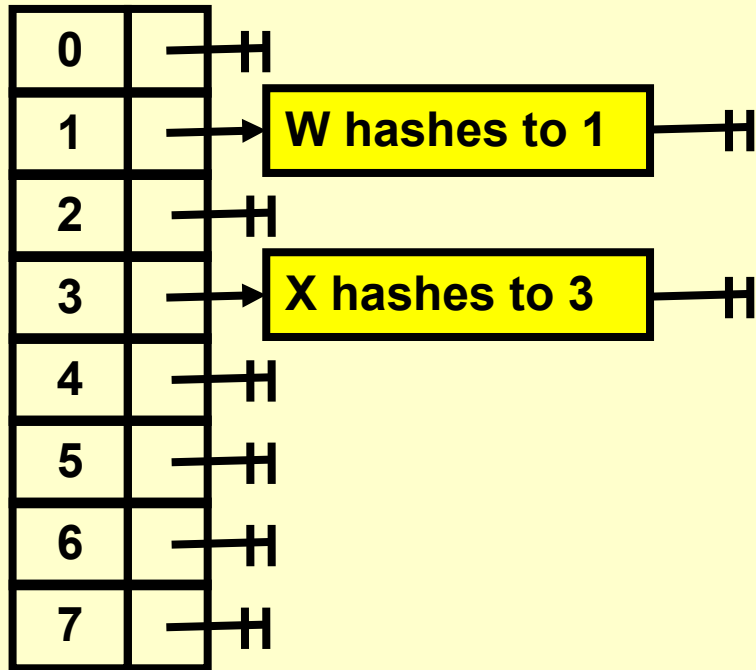
Collision Resolution

Technique: External chaining:



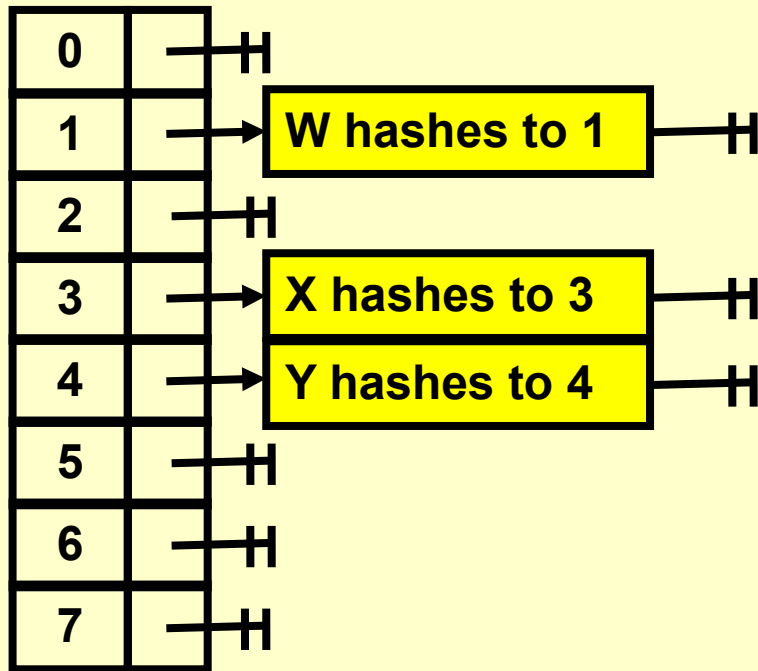
Collision Resolution

Technique: External chaining:



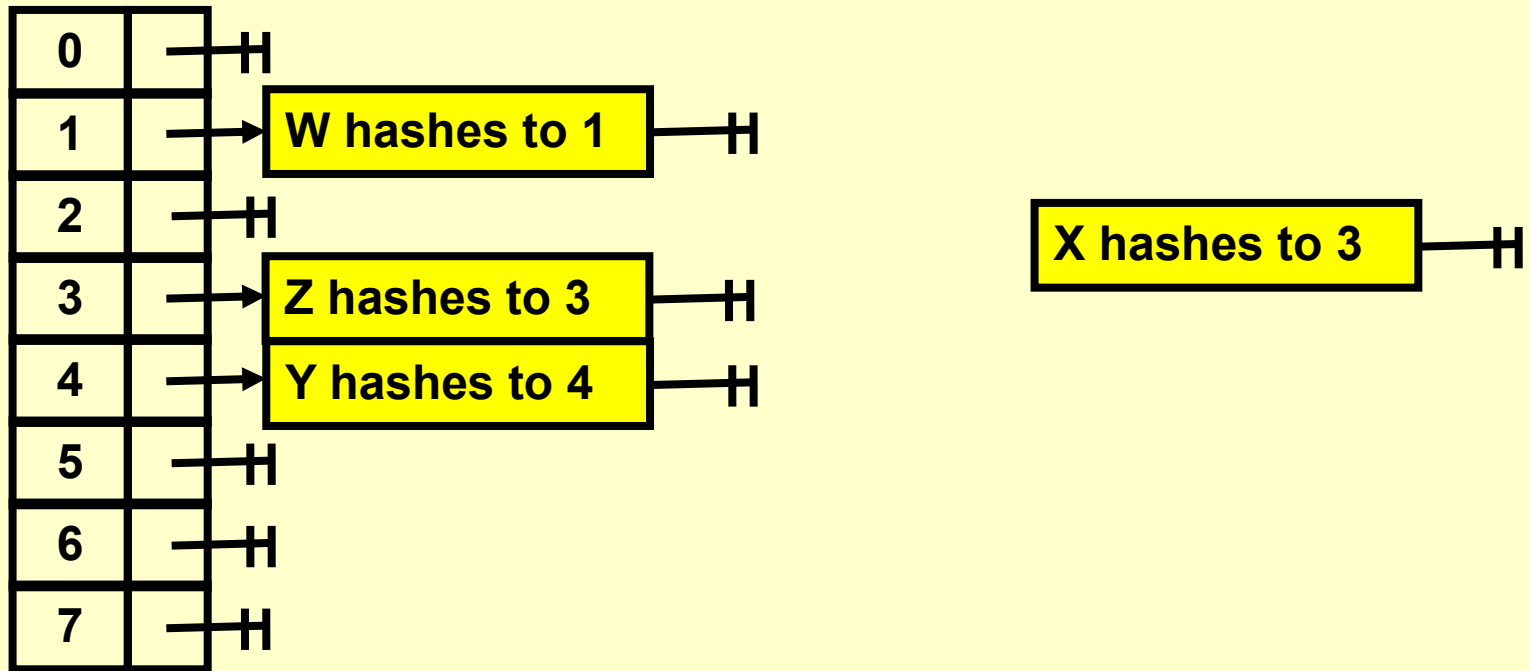
Collision Resolution

Technique: External chaining:



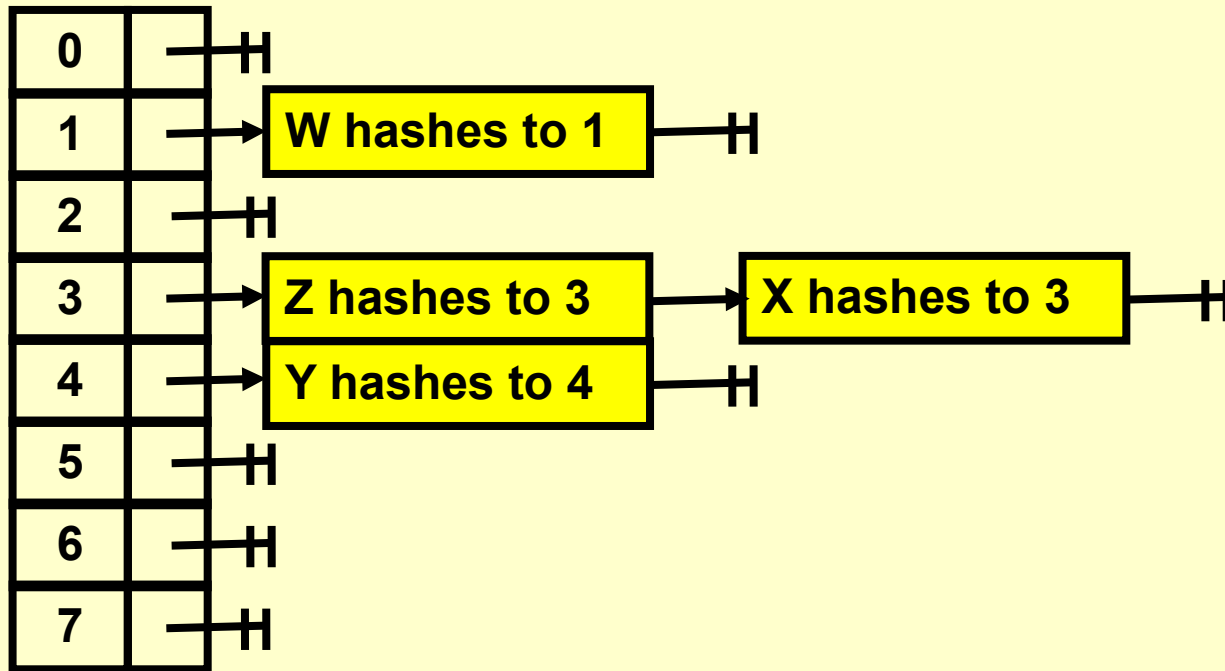
Collision Resolution

Technique: External chaining:



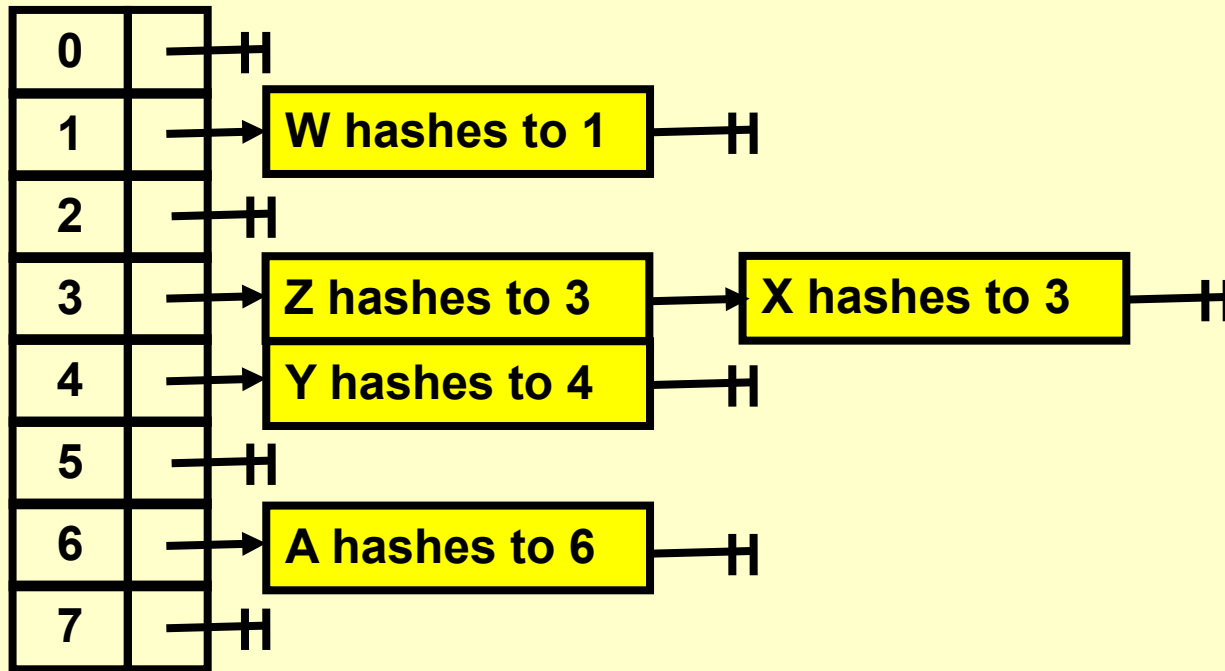
Collision Resolution

Technique: External chaining:



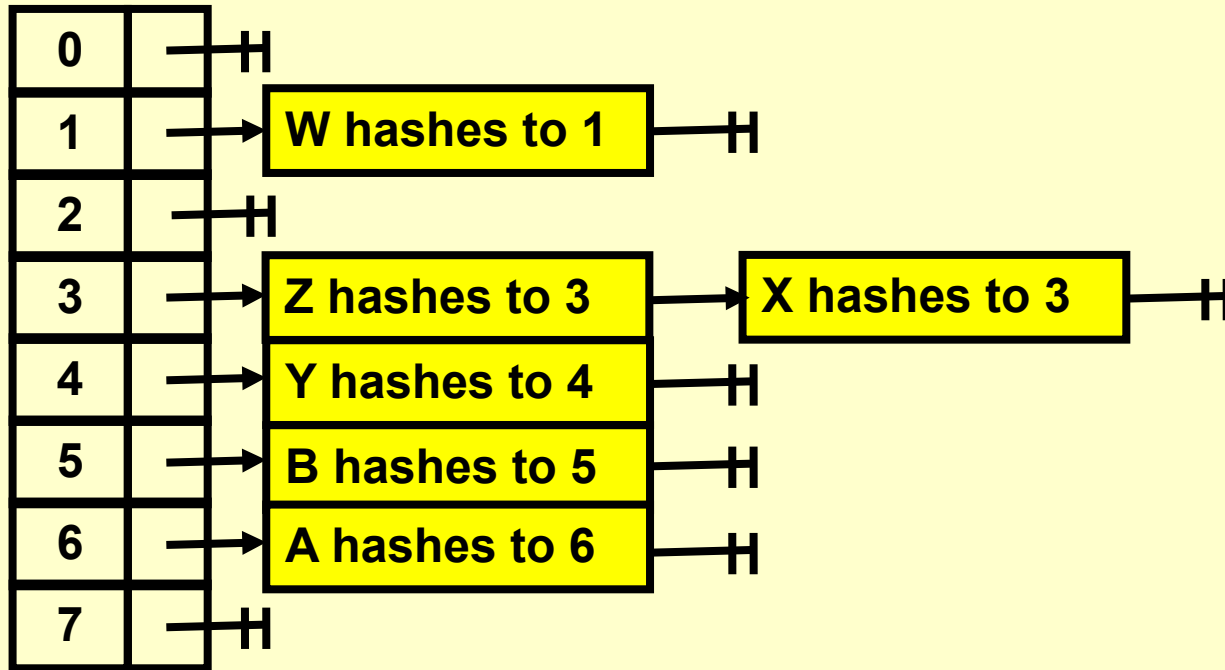
Collision Resolution

Technique: External chaining:



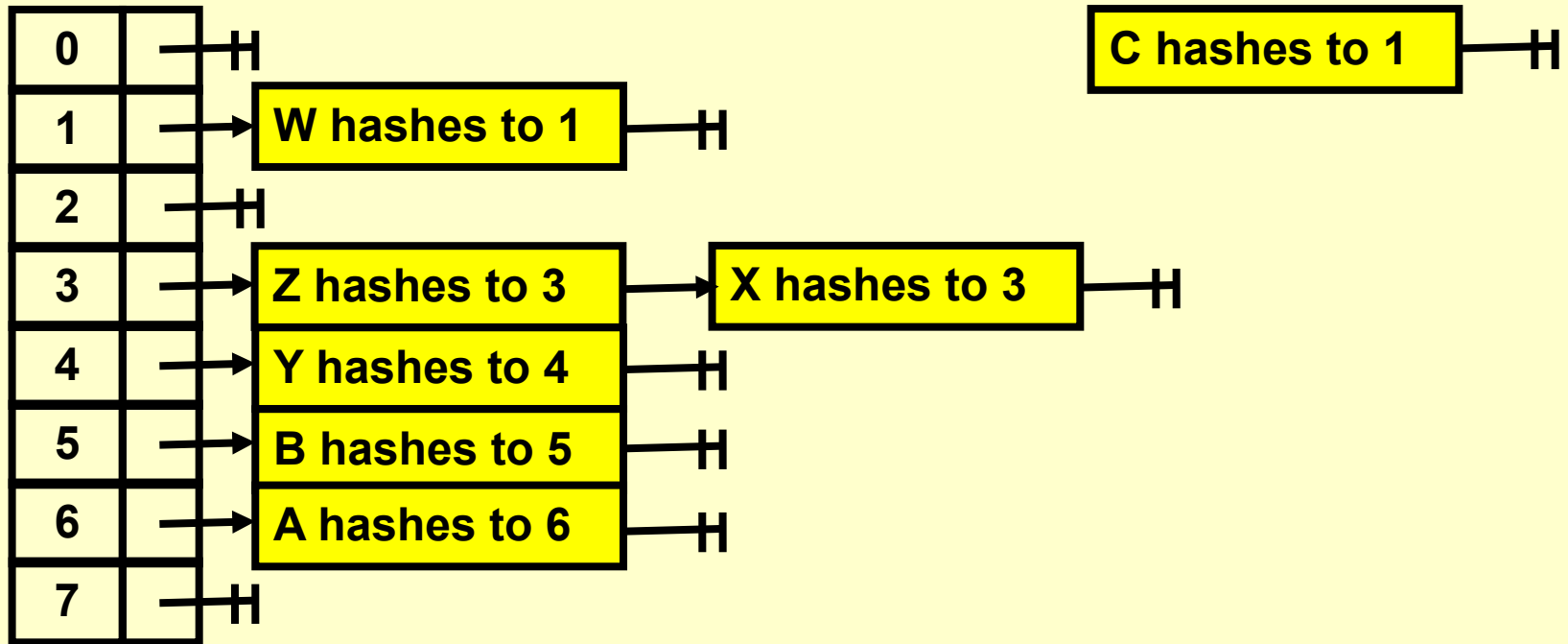
Collision Resolution

Technique: External chaining:



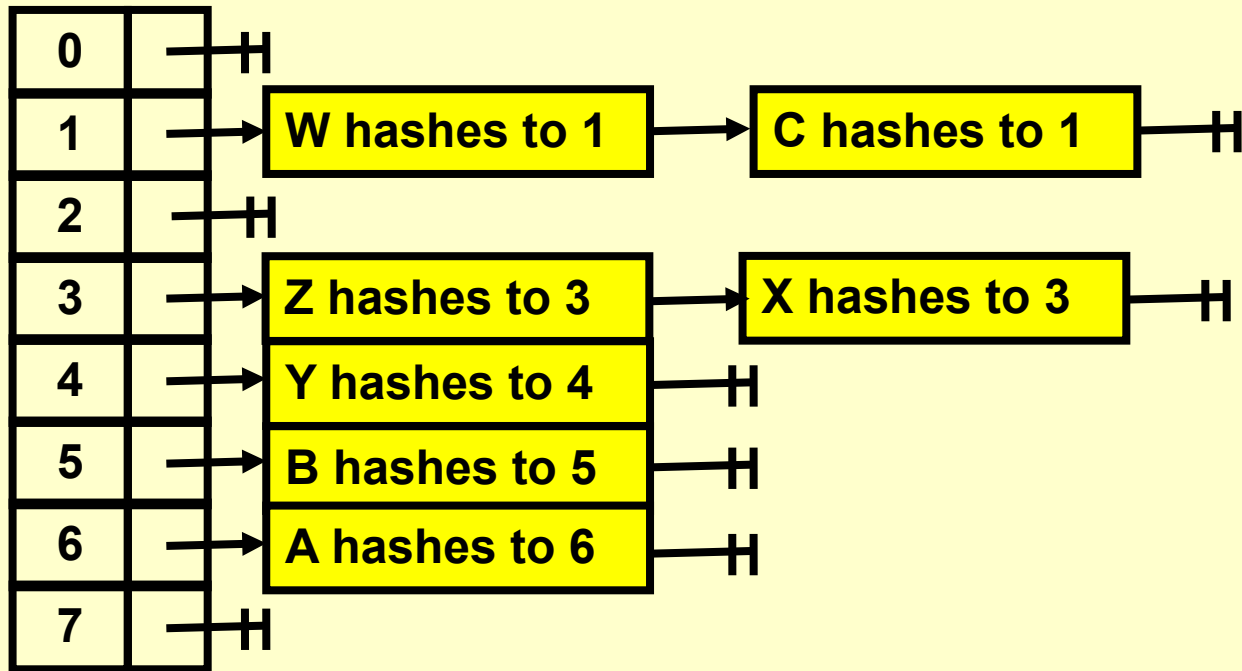
Collision Resolution

Technique: External chaining:



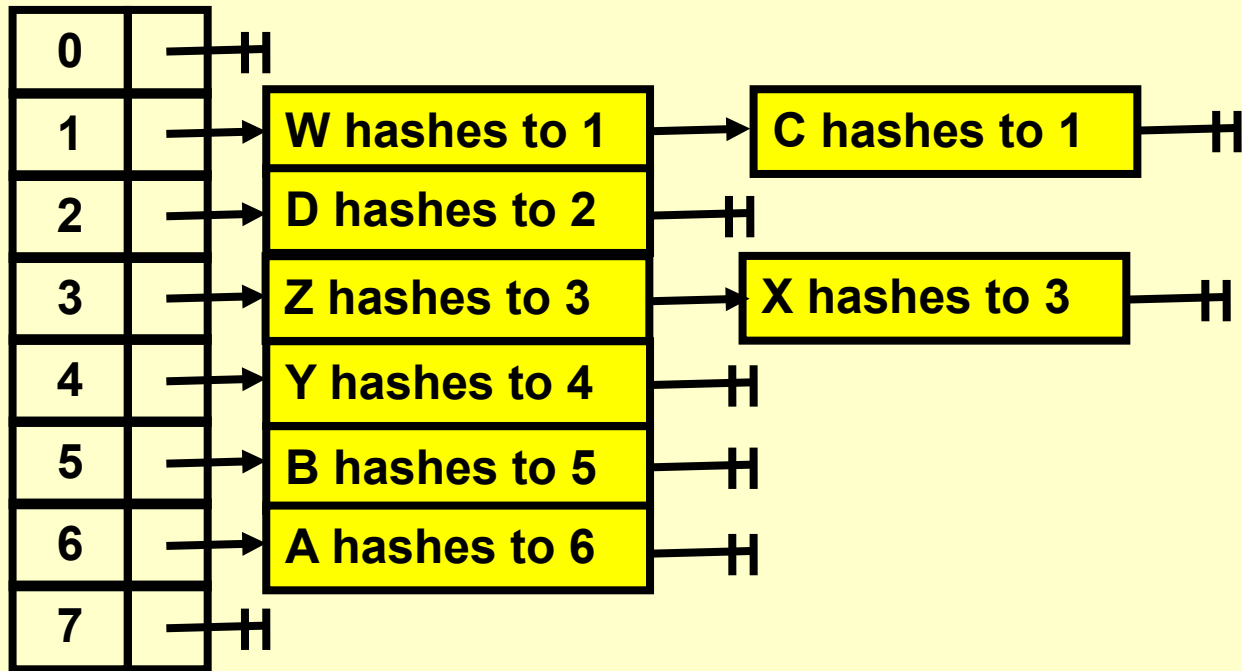
Collision Resolution

Technique: External chaining:



Collision Resolution

Technique: External chaining:



```
class AllContacts {  
    // storing contacts in a hashmap from names to contacts  
    HashMap<String, Contact> contactMap  
        = new HashMap<String, Contact>();  
  
    AllContacts() {}  
  
    // return contact that goes with a given name  
    // will return null if the name isn't in the map  
    Contact findContactByName(String name) {  
        return contactMap.get(name);  
    }  
  
    // takes the whole contact as input  
    AllContacts addContact(Contact aContact) {  
        Contact oldContact  
            = contactMap.put(aContact.name, aContact);  
        return this;  
    }  
}
```

```
class AllContacts {  
    // storing contacts in a hashmap from names to contacts  
    HashMap<String, LinkedList<Contact>> contactMap  
    = new HashMap<String, LinkedList<Contact>>();  
  
    AllContacts() {}  
  
    // return contact that goes with a given name  
    // will return null if the name isn't in the map  
    LinkedList<Contact> findContactByName(String name) {  
        return contactMap.get(name);  
    }  
  
    // takes the whole contact as input  
    AllContacts addContact(LinkedList<Contact> aContact) {  
        Contact oldContact  
        = contactMap.put(aContact.name, aContact);  
        return this;  
    }  
}
```

```
AllContacts addContact(LinkedList<Contact> aContact) {  
    Contact oldContact  
        = contactMap.put(aContact.name, aContact);  
    return this;  
}
```

```
AllContacts addContact(Contact aContact) {  
    LinkedList<Contact> contacts  
        = this.findContactByName(aContact.name);  
    contacts.add(aContact);  
    Contact oldContact  
        = contactMap.put(aContact.name, contacts);  
    return this;  
}
```

Better Hashing

The key to efficient hashing is the hash function.

How can we *efficiently* convert a name into a key number?

```
public int getHash(String strName) ;
```

Hashing Names

**Version
1**

```
public int getHash (String strName) {  
  
    int hash = 0;  
  
    for (int i =0; i < strName.length(); i++) {  
        hash += (int) strName.charAt(i);  
    }  
  
    hash %= tableSize;  
  
    return hash;  
}
```

**This works,
but what are
its limitations?**

Hashing Names

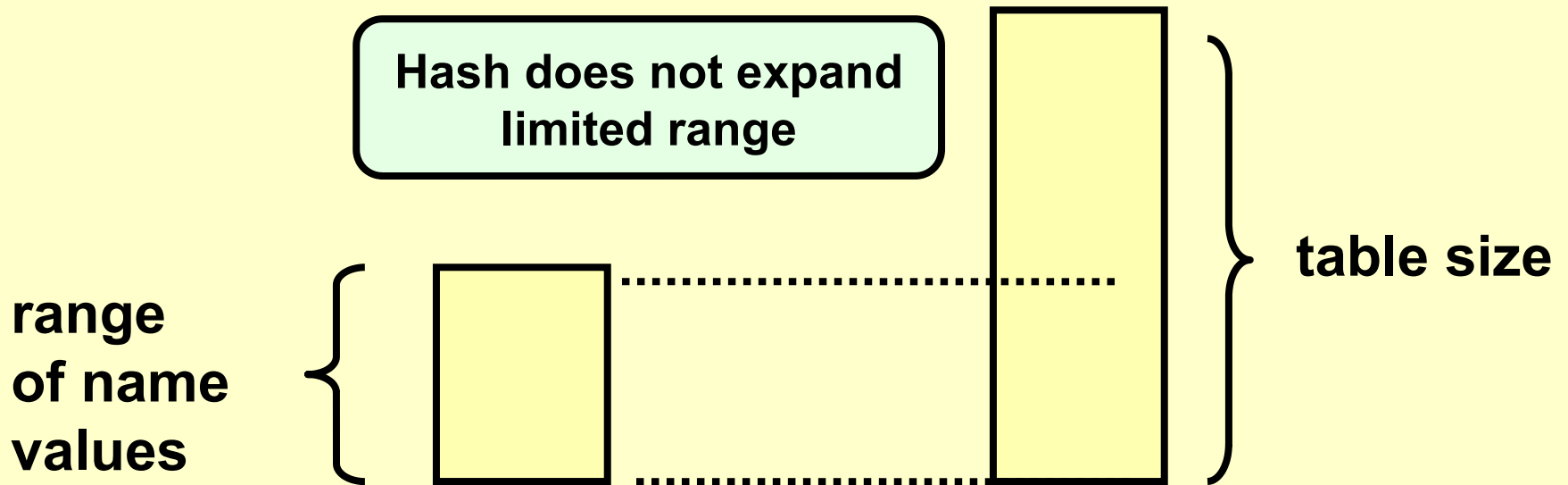
```
public int getHash (String strName) {  
  
    int hash = 0;  
  
    hash = (int) strName.charAt(0) +  
           128 * (byte) strName.charAt(1) +  
           16384 * (byte) strName.charAt(2);  
  
    hash %= tableSize;  
  
    return hash;  
}
```

**Version
2**

**Strategy:
only examine
first three
characters**

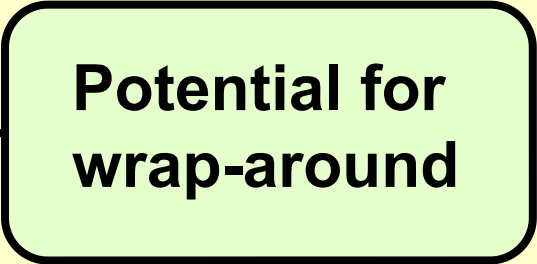
**This works,
but what are
its limitations?**

Inductive Analysis



Improved Hash Function

```
public int getHash (String strName) {  
  
    int hash = 0;  
  
    for (int i=0; i< strName.length(); i++)  
        hash = 128 * hash + (int) strName.charAt(i);  
  
    hash %= tableSize;  
  
    if (hash < 0 )  
        hash += tableSize;  
  
    return hash;  
}
```



Potential for
wrap-around

```
import java.util.HashMap;
```

```
class Addr {
```

```
    int num;
```

```
    String street;
```

```
    //... code snipped
```

```
    // a simple hashCode function multiplies the hashCode of each
```

```
    // field that matters by a unique prime,
```

```
    // and sums the results
```

```
    public int hashCode() {
```

```
        return (this.num * 3) + (this.street.hashCode() * 11);
```

```
    }
```

```
}
```

```
import java.util.HashMap;

class Addr {
    int num;
    String street;

    //... code snipped

    // a simple hashCode function multiplies the hashCode of each
    // field that matters by a unique prime,
    // and sums the results
    public int hashCode() {
        return (this.num * 3) + (this.street.hashCode() * 11);
    }
}
```

Hard Lessons about Hashing

Your hash function must be *carefully* selected.

It varies with your data. You have to study your input, and base your hash on the properties of the input data.

Your range of input should be larger than your table size (else your hashing will under utilize the table).

Experience (and theory beyond this course) show that prime numbers not close to a power of two will lead to the best distribution of keys

Summary of Hash Tables

- Purpose: Allows extremely fast lookup given a key value. Reduce the address space of the information to the table space of the hash table.
- Hash function: the reduction function
- Collision: $\text{hash}(a) == \text{hash}(b)$, but $a \neq b$
- Chaining is most general solution to collision