

## CS584 HW#6

Due: December 8

1. (2 points) Do **Exercise 26.1-6** of page 650 of CLRS, where the properties are the **Capacity Constraint, Skew Symmetry** and **Flow Conservation**.
2. (6 points) Given a directed graph  $G=(V,E)$ , a positive integer  $k$  and a pair of vertices  $x,y \in V$ , we want to decide if there exist  $k$  edge disjoint paths from  $x$  to  $y$ . Show how to solve this problem in polynomial time.

3. (6 points) Do **Exercise 26.2-9** of page 664 of CLRS. In the statement of the problem, replace “...at most  $|V|$  flow networks...” with “...at most  $|V|^2$  flow networks...”.

4. (5 points) Prove or give a counterexample to the following.

**CONJECTURE:** Assume that for arbitrary network  $G = (V, E)$ ,  $\sigma, \tau \in V$ , with capacities  $c : E \rightarrow \mathfrak{R}^+$ , cut  $(S, T)$  is a minimum cut. In network  $G^* = (V, E)$ ,  $\sigma, \tau \in V$ , with capacities  $c^* : E \rightarrow \mathfrak{R}^+$  with  $c^*(e) = c(e) + 42$  for all  $e \in E$ , cut  $(S, T)$  must be a minimum cut.

5. (8 points) Suppose that we are given

- network  $G = (V, E)$ ,
- $\sigma, \tau \in V$ ,
- $c : E \rightarrow \mathbb{Z}^+$ ,
- maximum flow  $f$  in  $G$ ,
- edge  $e \in E$ .

We want to compute the maximum flow  $f^*$  in  $G^* = (V, E)$ ,  $\sigma, \tau \in V$ ,  $c^* : E \rightarrow \mathfrak{R}^+$  where

$$c^*(e) = \begin{cases} c(e), & \text{if } e \neq e \\ c(e) + 1, & \text{if } e = e \end{cases}.$$

That is,  $G^*$  is identical to  $G$  except that the capacity of  $e$  is increased by 1. Find an algorithm to compute  $f^*$  in time in  $O(|V| + |E|)$ .

6. (6 points) Adapted from *Algorithm Design* by Kleinberg & Tardos.  
As a response to natural disasters, we want an algorithm to assign  $n$  injured people to  $k$  hospitals. There is an unlimited supply of ambulances, but everybody should be transported to a hospital within a 30 minute drive of their current location. Finally, in order to balance the load, every hospital should receive at most  $\lceil n/k \rceil$  people. Describe a polynomial time algorithm to accept as input an  $n \times k$  array of driving times, and which will decide if a balanced assignment of people to hospitals is possible.

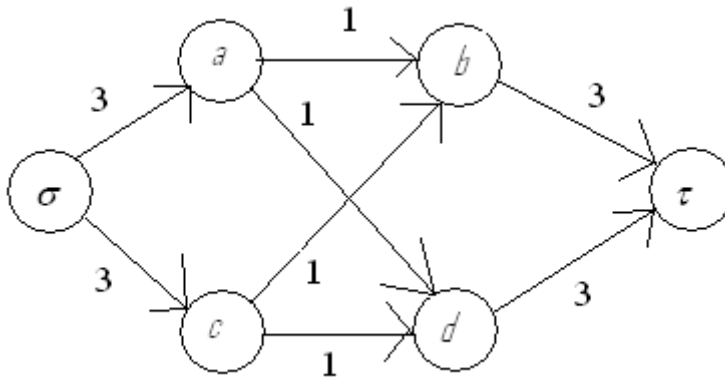
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**HW#6 SOLUTIONS**

1.  $f_1+f_2$  satisfies **Skew Symmetry** and **Flow Conservation**, but might violate the **Capacity Constraint**.

2. Assign a capacity of 1 to each edge  $e \in E$ . Run the Ford-Fulkerson Labelling Algorithm to compute the maximal flow from  $x$  to  $y$ . Each augmenting path from  $x \rightarrow y$  is edge disjoint from all previous augmenting paths. The algorithm should respond  $y \in S$  if there's a flow of at least  $k$  from  $x \rightarrow y$ , and no otherwise.

3. For each pair of vertices  $u, v$ , of  $G$ , define a flow network  $G_{uv}$  with  $s=u, t=v$ , and all edge capacities 1. The maximum flow in  $G_{uv}$  yields a minimum, and hence a minimum set of edges which separate  $u$  and  $v$ . The minimum over all pairs  $u$  and  $v$  is the edge connectivity of  $G$ .

4. The CONJECTURE is false. Network



has minimum cut  $(\{\sigma, a, c\}, \{b, d, \tau\})$ . But the network obtained when 42 is added to every capacity has two min cuts,  $(\{\sigma\}, \{a, b, c, d, \tau\})$  and  $(\{\sigma, a, b, c, d\}, \{\tau\})$ .

5. We form a graph  $\Upsilon(G^*)$  with edges of positive residual capacity relative to  $f$ . That is,  $\Upsilon(G^*) = (V, E')$  where  $E' = \{\varepsilon \in E \mid c^*(\varepsilon) - f(\varepsilon) > 0\}$ . Each  $\varepsilon \in E'$  has capacity  $\kappa(\varepsilon) = c^*(\varepsilon) - f(\varepsilon)$ . That is, the capacity of  $\varepsilon$  is its residual capacity relative to  $f$ . Constructing  $\Upsilon(G^*)$  takes time in  $O(|V| + |E|)$ . We seek a path in  $O(|V| + |E|)$  from  $\sigma$  to  $\tau$ . If such a path exists, it is an augmenting path relative to  $f$ , and it can be found in time in  $O(|V| + |E|)$ .

6. Let the hospitals be  $\{h_1, \dots, h_k\}$  and the people be  $\{p_1, \dots, p_n\}$ . We form a network with vertices  $\{h_1, \dots, h_k, p_1, \dots, p_n, \sigma, \tau\}$  and edges  $\sigma p_i, 1 \leq i \leq n$ , all of capacity 1, and edges

$h_j\tau, 1 \leq j \leq k$ , all of capacity  $\lceil n/k \rceil$ . Finally, there is an edge  $p_i h_j, 1 \leq i \leq n, 1 \leq j \leq k$  of capacity 1 if and only if  $p_i$  is within a 30 minute drive of  $h_j$ . There exists a balanced assignment of people to hospitals if and only if the network admits a flow of volume  $n$ .