

CS3133
HW#7 SOLUTIONS

1. Removing unit productions yields

$S \rightarrow \text{if } C \text{ then } A \mid \text{while } C \text{ do } A \mid \text{if } C \text{ then } A \text{ else } S$

$B \rightarrow \text{if } C \text{ then } A \text{ else } S$

$C \rightarrow c_1 \mid c_2 \mid c_3$

$A \rightarrow a_1 \mid a_2 \mid a_3$

Since B is now unreachable, removing useless productions now yields

$S \rightarrow \text{if } C \text{ then } A \mid \text{while } C \text{ do } A \mid \text{if } C \text{ then } A \text{ else } S$

$C \rightarrow c_1 \mid c_2 \mid c_3$

$A \rightarrow a_1 \mid a_2 \mid a_3$

Finally, we can convert to Chomsky Normal Form

$S \rightarrow X_1Y_1 \mid X_3Y_3 \mid X_1Y_5$

$X_1 \rightarrow \text{if}$

$Y_1 \rightarrow CY_2$

$Y_2 \rightarrow X_2A$

$X_2 \rightarrow \text{then}$

$X_3 \rightarrow \text{while}$

$Y_3 \rightarrow CY_4$

$Y_4 \rightarrow X_4A$

$X_4 \rightarrow \text{do}$

$Y_5 \rightarrow CY_6$

$Y_6 \rightarrow X_2Y_7$

$Y_7 \rightarrow AY_8$

$Y_8 \rightarrow X_5S$

$X_5 \rightarrow \text{else}$

$C \rightarrow c_1 \mid c_2 \mid c_3$

$A \rightarrow a_1 \mid a_2 \mid a_3$

2. Let n be given by the PUMPING LEMMA, and consider $z = 0^n 1^n 1^n 0^n 0^n 1^n = 0^n 1^{2n} 0^{2n} 1^n$.

When z is rewritten as $uvwxy$ with $|vwx| \leq n$ and $|vx| \geq 1$, then vx contains at least a 0 or at least a 1, and it can not contain 0's from both clumps of 0's nor can it contain 1's from both clumps of 1's. Then uv^2wx^2y no longer has twice as many 0's in the second clump as the first or no longer has twice as many 1's in the first clump as the second.

3. Let $n=42$. For every $z \in L, |z| \geq 42$, we must be able to express z as $uvwxy$ which satisfies certain conditions. Since there is no such z , these conditions are vacuously satisfied, and we can **not** create a contradiction.

4. **a)** $L(G) = \emptyset$ if and only if S is generating. In fact, since useless symbols have been removed from G , we can conclude that $L(G) = \emptyset$ if and only if $P = \emptyset$.

b) Yes, $L(G)$ must be infinite. Since G has no useless symbols, A is reachable,

$S \Rightarrow^* \mathbf{aAb}$ $\mathbf{a}, \mathbf{b} \in (V \cup T)^*$, $w \in T^*$, and generating, $A \Rightarrow^* w$, $w \in T^*$. Also, there are

derivations $\mathbf{a} \Rightarrow^* u$, $u \in T^*$ and $\mathbf{b} \Rightarrow^* y$, $y \in T^*$, so $S \Rightarrow^* \mathbf{aAb} \Rightarrow^* uwy$. But also

$A \Rightarrow^+ \mathbf{gAV}$, $\mathbf{g}, \mathbf{V} \in (V \cup T)^*$, and $\mathbf{g} \Rightarrow^* v$, $v \in T^*$ and $\mathbf{V} \Rightarrow^* x$, $x \in T^*$ where $|vx| \geq 1$. It

follows that $S \Rightarrow^* uwy$, $S \Rightarrow^* uvwxy$, $S \Rightarrow^* uvvwxy$, ..., and $\{uv^k wx^k y \mid k \geq 0\} \subseteq L(G)$.