

**CS2223**  
**FINAL EXAM**

Name \_\_\_\_\_

**Date:** May 2, 2006  
All documentation permitted

1 (25) \_\_\_\_\_

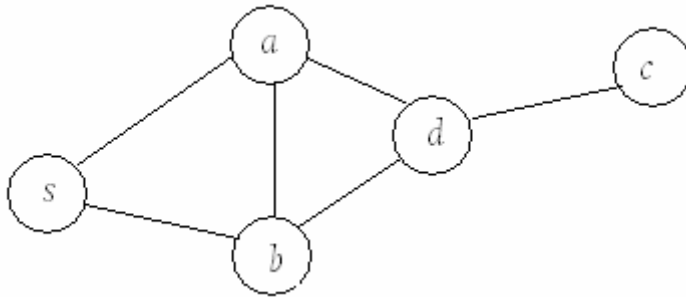
2 (25) \_\_\_\_\_

3 (25) \_\_\_\_\_

4 (25) \_\_\_\_\_

TOTAL \_\_\_\_\_

1. (25 points) We define the *length* of a path in a graph  $G = (V, E)$  to be the number of edges on the path, and the *distance* between two vertices is the length of a shortest path. Describe an algorithm which accepts as input a connected graph  $G = (V, E)$  and a source  $s \in V$  and which computes the distance between  $s$  and  $v$  for every  $v \in V$ . For graph



the output would be

$v$	$s$	$a$	$b$	$c$	$d$
<i>distance</i>	0	1	1	3	2

The execution time of your algorithm should be in  $O(|V| + |E|)$ .

**Note:** The graph is not a weighted graph. Every edge contributes 1 to the length of a path.

2. (25 points) Let  $G = (V, E)$  be a connected weighted graph with weight function  $w: E \rightarrow \mathfrak{R}^+$ , and assume that for fixed  $v \in V$ , the edge  $(v, y) \in E$  is the unique edge of minimum weight incident with  $v$ . That is, for all edges  $(v, z) \in E$  with  $y \neq z$ ,  $w((v, y)) < w((v, z))$ . Either prove that  $(v, y)$  must belong to all minimum spanning trees of  $G$  or provide a counter-example in which there is a minimum spanning tree to which  $(v, y)$  doesn't belong.

3. (25 points) Given an array  $A[1..n]$  of non-negative reals, we want to find

$\max_{1 \leq i \leq j \leq n} \left( \prod_{i \leq k \leq j} A[k] \right)$ . That is, we want to find the maximum product realizable as the

product of a contiguous subsequence of numbers of  $A$ . For example, if

$A=(0.6, 0, 23, 0.5, 28, 0.9)$ , then the answer is  $322=23*0.5*28$ . If  $A=(0.8, 0.9, 0.6)$ , then

the answer is  $0.9$ . If  $A=(1, 0, 12, 3, 1, 2, 0.001, 39)$ , then the answer is  $72=12*3*1*2$ .

Describe an algorithm with time complexity in  $O(n)$  to solve this problem.

4. (25 points) A graph  $G = (V, E)$  has a *bridge* if  $G$  is connected and has an edge  $e \in E$  whose removal disconnects  $G$ . Give an algorithm to test if  $G$  has a bridge. The execution time of your algorithm should be in  $O(|E| * (|V| + |E|))$ .

**CS2223**  
Solutions to FINAL EXAM

1. Do a breadth-first search from  $s$ , and the number of edges on the path in the bfs-tree from  $s$  to any  $v \in V$  is the distance from  $s$  to  $v$ .
2. Given the conditions of the problem,  $(v, y)$  must belong to all minimum spanning trees of  $G$ . If not, there must be a minimum spanning tree  $T$  of  $G$  which doesn't contain  $(v, y)$ . Adding  $(v, y)$  to  $T$  yields a connected graph with exactly one cycle, and  $v$  is on the cycle. The other edge of the cycle incident with  $v$  has a higher weight than  $(v, y)$ . So removing that edge from  $T \cup \{(v, y)\}$  weighs less than  $T$  and is a spanning tree, contradicting the fact that  $T$  is a minimum spanning tree.
3.  $best\_to[1] \leftarrow A[1]$   
**for**  $k \leftarrow 2$  **to**  $n$  **do**  
     $best\_to[k] \leftarrow \max(A[k], best\_to[k-1]*A[k])$
4. Do a dfs of  $G$  to test if it is connected.  
    **if**  $G$  is not connected **then return**  $G$  does not have a bridge  
    **for each**  $e \in E$  test if  $G-e$  is connected  
        **if** it is not **then return**  $G$  has a bridge  
    **return**  $G$  does not have a bridge