# Token Ring and Fiber Distributed Data Interface (FDDI)



# IEEE 802.5 Token Ring

• Proposed in 1969 and initially referred to as a *Newhall ring*.

**Token ring:** a number of stations connected by transmission links in a ring topology. Information flows *in one direction along the ring* from source to destination and back to source.

Medium access control is provided by a small frame, the token, that circulates around the ring when all stations are idle. *Only the station possessing the token is allowed to transmit at any given time*.

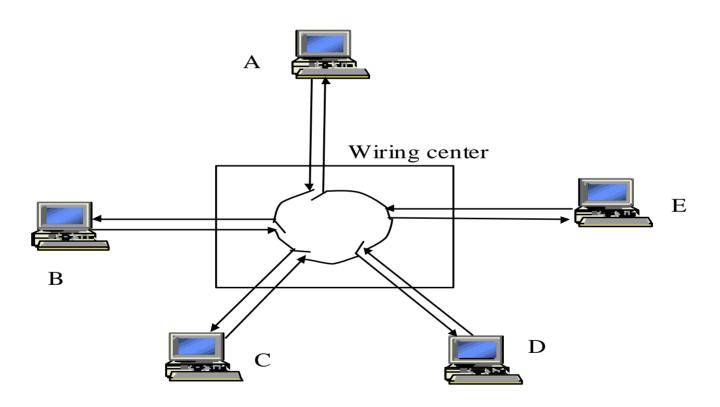


# Token Ring Operation

- When a station wishes to transmit, it must wait for the token to pass by and *seize the token*.
  - One approach: change one bit in token which transforms it into a "start-of-frame sequence" and appends frame for transmission.
  - Second approach: station claims token by removing it from the ring.
- The data frame circles the ring and is removed by the transmitting station.
- Each station interrogates passing frame. If destined for station, it copies the frame into local buffer. {Normally, there is a one bit delay as the frame passes through a station.}



# Token Ring Network with Star Topology



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Figure 6.58



### Token Insertion Choices

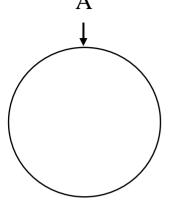
- 1. multi-token: insert token after station has completed transmission of the last bit of the frame.
- 2. single-token: insert token after last bit of busy token is received and the last bit of the frame is transmitted.
- 3. single-frame: insert token after the last bit of the frame has returned to the sending station.

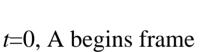
**Performance** is determined by whether more than one frame is allowed on the ring at the same time and the relative propagation time.

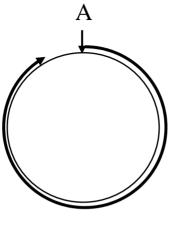


#### Low Latency Ring (a)

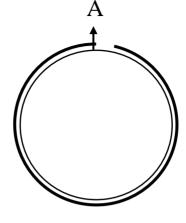
#### Single frame operation



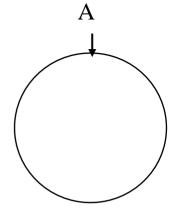




*t*=90, return of first bit

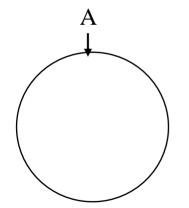


*t*=400, transmit last bit

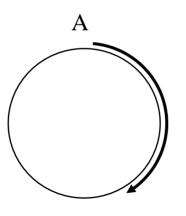


*t*=490, reinsert token

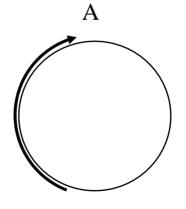
#### (b) High Latency Ring



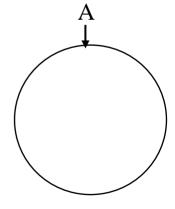
*t*=0, A begins frame



*t*=400, last bit of frame enters ring



*t*=840, return of first bit

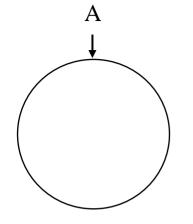


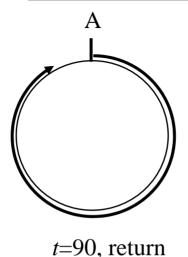
*t*=1240, reinsert token

Networks: Token Ring and FDDI

#### Low Latency Ring (a)

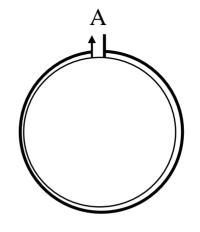
#### Single token operation

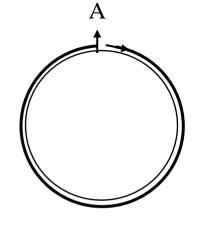




of first bit

A



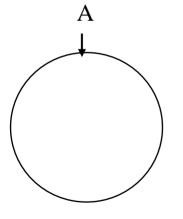


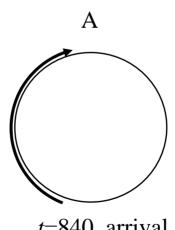
*t*=0, A begins frame

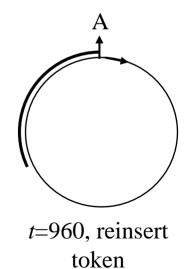
*t*=210, return of header

t=400, last bit enters ring, reinsert token

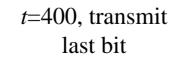
High Latency Ring (b)







*t*=0, A begins frame



*t*=840, arrival first frame bit

# IEEE 802.5 Token Ring

- 4 and 16 Mbps using twisted-pair cabling with differential Manchester line encoding.
- Maximum number of stations is 250.
- 4Mbps 802.5 token ring uses single frame operation.
- 4 Mbps IBM token ring uses single token operation.
- Both 802.5 and IBM 16Mbps token rings use *multi-token* operation.
- 802.5 has 8 priority levels provided via two 3-bit fields (priority and reservation) in data and token frames.
- Permits 16-bit and 48-bit addresses (same as 802.3).



# Token Ring

- Under light load delay is added due to waiting for the token (on average delay is one half ring propagation time).
- Under heavy load ring is "round-robin".
  - Performance is fairer and better than Ethernet!!
- The ring must be long enough to hold the complete token.
- Advantages fair access, no collisions.
- Disadvantages ring is single point of failure, ring maintenance is complex due to token malfunctions.



### Token Maintenance Issues

#### What can go wrong?

- Loss of token (no token circulating)
- Duplication of token (forgeries or mistakes)
- The need to designate one station as the *active ring monitor*.
- Persistently circulating frame
- Deal with active monitor going down.



#### IEEE 802.5 Token and Data Frame Structure

**Token Frame Format** 

SD AC ED

Data Frame Format

1	1	1	2 or 6	2 or 6		4	1	1
SD	AC	FC	Destination	Source	Information	FCS	ED	FS
			Address	Address				

Starting delimiter

J K 0 J K 0 0

J, K non-data symbols (line code)

Access control

PPP T M RRR

PPP Priority; T Token bit M Monitor bit; RRR Reservation

Frame control

FF Z Z Z Z Z Z

FF frame type ZZZZZZ control bit

Ending delimiter

J K 1 J K 1 I E

I intermediate-frame bitE error-detection bit

Frame status

A C xx A C x x

A address-recognized bitxx undefinedC frame-copied bit

Networks: Token Ring and FDDI

# Fiber Distributed Data Interface (FDDI)

- FDDI uses a ring topology of multimode or single mode optical fiber transmission links operating at 100 Mbps to span up to 200 kms and permits up to 500 stations.
- Employs dual counter-rotating rings.
- 16 and 48-bit addresses are allowed.
- In FDDI, token is absorbed by station and released as soon as it completes the frame transmission *[multi-token operation]*.



# FDDI: Dual Token Ring

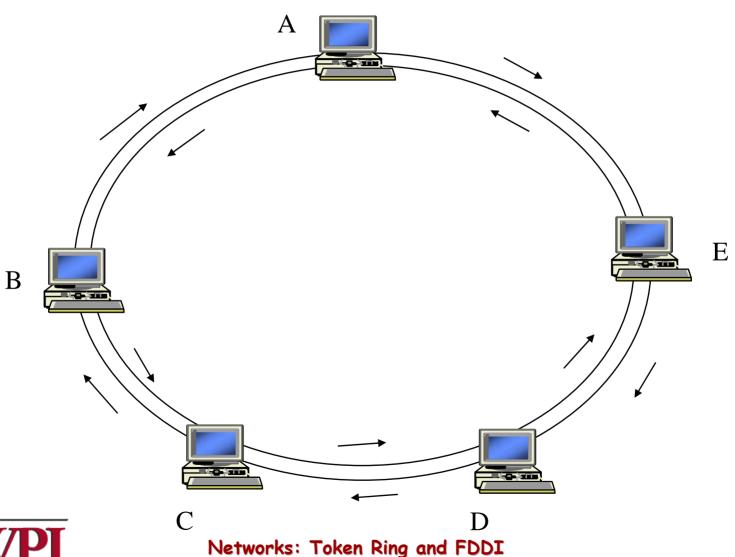
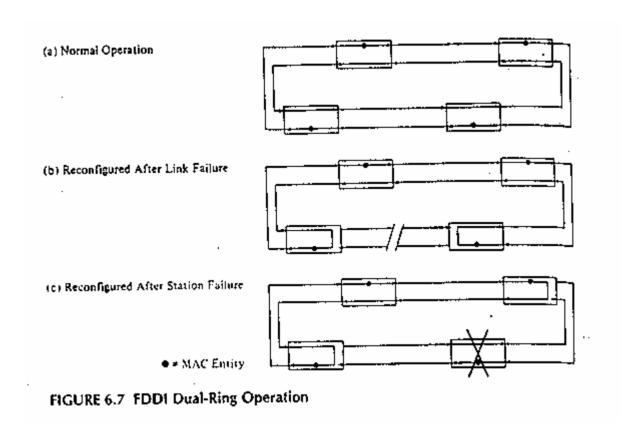


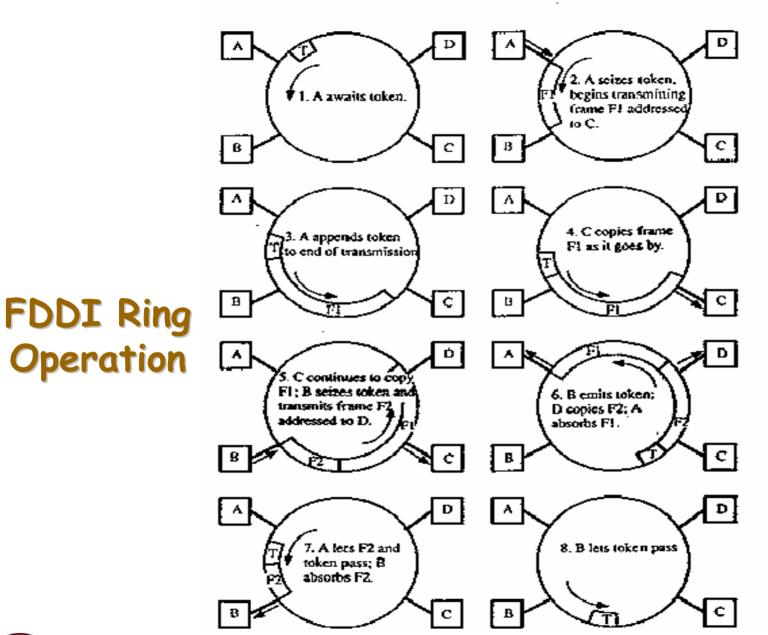
Figure 6.62

# FDDI Repair





14





15

# **FDDI**

- To accommodate a mixture of stream and bursty traffic, FDDI is designed to handle two types of traffic:
  - Synchronous frames that typically have tighter delay requirements (e.g., voice and video)
  - Asynchronous frames have greater delay tolerances (e.g., data traffic)
- FDDI uses TTRT (Target Token Rotation Time) to ensure that token rotation time is less than some value.



# FDDI Data Encoding

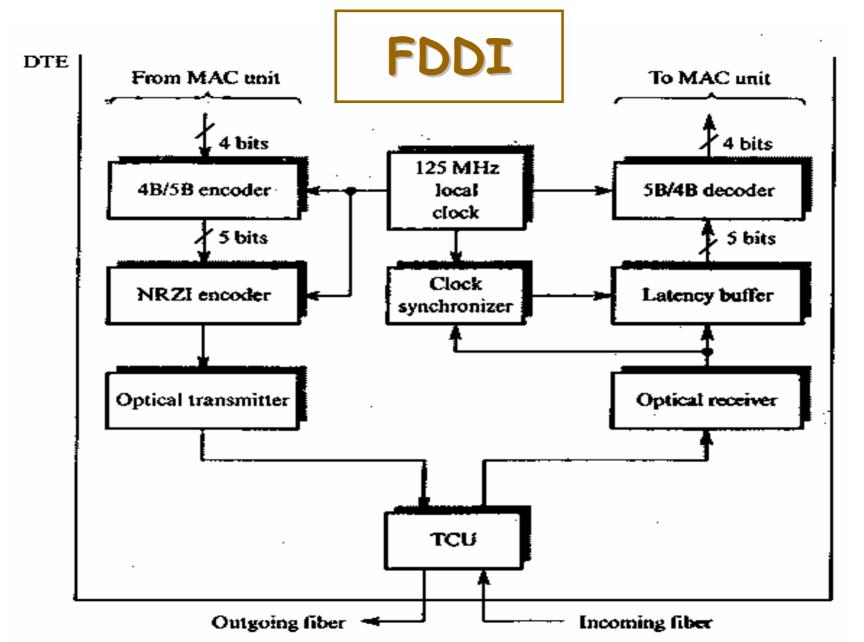
- Cannot use differential Manchester because 100 Mbps FDDI would require 200 Mbaud!
- Instead each ring interface has its own <u>local</u> <u>clock</u>.
  - Outgoing data is transmitted using this clock.
  - Incoming data is received using a clock that is frequency and phase locked to the transitions in the incoming bit stream.



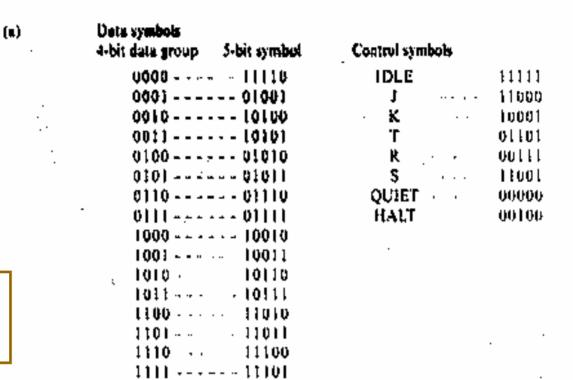
# FDDI Data Encoding

- Data is encoded using a 4B/5B encoder.
  - For each four bits of data transmitted, a corresponding 5bit codeword is generated by the encoder.
  - There is a maximum of two consecutive zero bits in each symbol.
- The symbols are then shifted out through a NRZI encoder which produces a signal transition whenever a 1 bit is being transmitted and no transition when a 0 bit is transmitted.
- Local clock is 125MHz. This yields 100 Mbps (80% due to 4B/5B).









FDDI

(b) PA SD FC DA SA INFORMATION FCS ED FS Information

	والمنافذ الأراب			in .
PA		FC	ΕD	Token
				ı

Figure 7.15
FDDI line coding and framing detail:

- (a) 4858 codes;
- (b) frame formats.

- PA = Preamble (16 or more symbols)
- SD = Start delimiter (2 symbols)
- FC = Frame control (2 symbols)
- DA = Destination address (4 or 12 symbols)

SA = Source address (4 or 12 symbols)

FCS = Frame check sequence (8 symbols)

ED = End delimiter (1 or 2 symbols)

FS = Frame status (3 symbols)

## FDDI Frame Structure

**Token Frame Format** 

PRE SI	) FC	ED
--------	------	----

Data Frame Format

8	1	. 1	2 or 6	2 or 6		4	1	1
PRE	SD	FC	Destination	Source	Information	FCS	ED	FS
			Address	Address				

Preamble

Frame CLFFZZZZ

Control

C = Synch/Asynch

L = Address length (16 or 48 bits)

FF = LLC/MAC control/reserved frame type

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Figure 6.63



### More FDDI Details

- FDDI Transmission on optical fiber requires ASK.
- The simplest case: coding is done via the absence or presence of a carrier signal {Intensity Modulation}.
- Specific 5-bit codeword patterns chosen to guarantee no more than **three zeroes in a row** to provide for adequate synchronization.
- 1300 nm wavelength specified.
- Dual rings (primary and secondary) transmit in opposite directions.
- Normally, second ring is **idle** and used for redundancy for automatic repair (self-healing).



# IEEE 802.5 and FDDI Differences

#### Token Ring

- Shielded twisted pair
- 4, 16 Mbps
- No reliability specified
- Differential Manchester
- Centralized clock
- Priority and Reservation bits
- All three token operations possible

#### **FDDI**

- Optical Fiber
- 100 Mbps
- Reliability specified (dual ring)
- 4B/5B encoding
- Distributed clocking
- Timed Token Rotation Time
- Multi-token operation

