

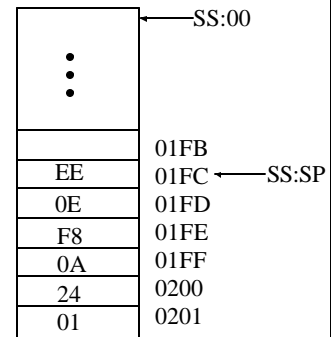
Lecture 16: Passing Parameters on the Stack

- Quick Stack Review
- Passing Parameters on the Stack
- Binary/ASCII conversion

Push Examples

;assume SP =
0202

```
mov ax, 124h
push ax
push 0af8h
push 0eeeh
```



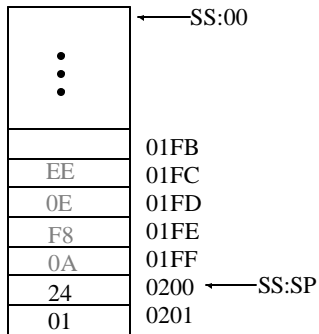
Pop Examples

;initial SP = 01FC

```
pop BX
pop DX
```

BX = 0EEh
DX = 0AF8

If we were to pop again, the next value popped off would be 0124.



CALL and RET

- CALL – pushes IP on the stack (recall, IP holds the address of the next instruction), puts the address of the label (subroutine) into IP.
- RET – pops the stack into IP to return to the point at which the subroutine was called.

Passing Parameters using Registers

- Passing parameter values to a subroutine
- Passing base address of parameter to a subroutine.
- Example – keyboard buffer:


```
maxlen db 20 ; max chars to
input
actualen db ? ;DOS will put the
number read here
inbuf db 20 dup (' '); where
DOS will put the data read in
....
mov dx, offset maxlen
```

Passing Parameters Using Registers, cont.

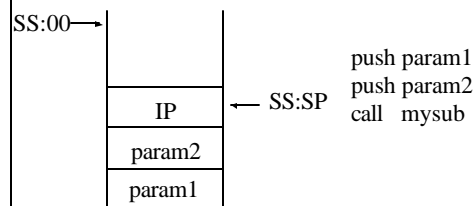
- Advantages:
- Problems:

Passing Parameters using the Stack

- Several parameters? – use the stack
- Push parameters before the CALL
- Why can't you just pop them off again in your subroutine?

Passing Parameters on the Stack, cont.

- When you execute CALL, it pushes IP on the stack.
- If your called routine executes a pop it won't get a parameter, it will get the return IP.



Note: stack drawn as 16-bit WORDS

Getting at Our Parameters

- You can pop the IP off the stack and then get your data.
 - Don't do this!!!!
 - Why?
- A better solution – use the BP register!
 - BP works like BX except it holds an offset from SS (recall BX holds an offset from DS)

Indirect Addressing Using BP

`mov ax, [bx]` ; this moves the
; word at DS:BX
; into AX

`mov dx, [bp]` ; this moves the
; word at SS:BP
; into DX

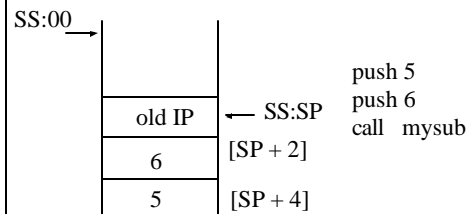
- you can also use indirect addressing with displacement:
 - `mov ax, [bp + 2]` ;take the address
;in BP, add 2
;load AX with the
;contents of the word at
;that address.

Setting up a Stack Frame

- Parameters are passed on the stack by setting up a stack frame.
 - The calling routine pushes the parameters on the stack.
 - The CALL instruction is used to call the subroutine.
 - The subroutine pushes BP on the stack (why?).
 - The subroutine copies the value in SP into BP so it can be used to retrieve the variables.

Stack Frame Example

- After calling our subroutine:

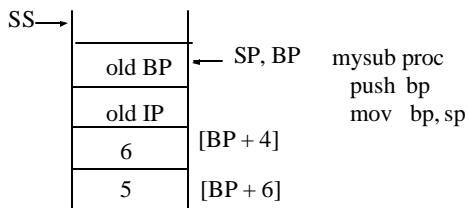


Note: stack drawn as 16-bit WORDS

old IP? The return address – this is the offset of the instruction immediately after the call.

Stack Frame Example, cont.

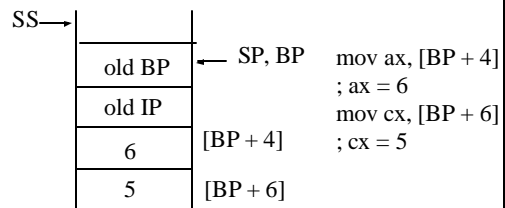
- Save and set up BP (in procedure mysub):



Note: stack drawn as 16-bit WORDS

Retrieving Parameters

- Use indirect addressing with displacement:



Note: stack drawn as 16-bit WORDS

```

TITLE Demonstrates parameter passing on the stack
;
; This program adds a constant value to each element of an array.
; It uses a procedure, and parameters are passed on the stack in
; the following order: the start address of the array, the constant
; value to be added, and the number of elements in the array.

.model small
.stack 100H

AddNumber EQU 12H ;number to be added
ArraySize EQU 9H

.data
TestArray DW 3AH, 4AH, 5AH, 6AH, 7AH, 1234H, 5678H
           DW 9ABCH, 0DEFH

.code
.startup

mov ax, offset TestArray
push ax
mov ax, AddNumber
push ax
mov ax, ArraySize
push ax
call ArrayInc
nop
.exit
    
```

```

ArrayInc PROC NEAR

push bp
mov bp, sp ;set up frame for this call

push ax ;save registers destroyed by this
push bx ;procedure
push cx
push dx

mov cx, [bp+4] ;size of array
mov ax, [bp+6] ;the constant to be added
mov bx, [bp+8] ;start address of array

next: mov dx, [bx] ;get array element
      add dx, ax ;add constant to it
      mov [bx], dx ;store it back in array
      add bx, 2 ;point to next array element
      loop next

pop dx ;restore registers
pop cx
pop bx
pop ax

pop bp
ret ;returns

ArrayInc ENDP
end
    
```

What's Wrong with this Example?

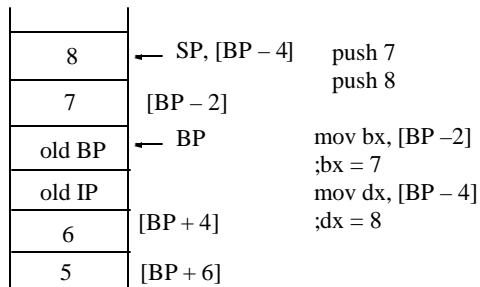
Cleaning up the Stack

- Clean up the stack using RET x, where x is the number of bytes to add to SP after the return:

```
RET 6 ;clean up the 6 bytes of
      ;parameters by adding
      ;6 to SP after popping IP
```

Using the Stack for Local Variables

- The stack is useful for local variables



Note: stack drawn as 16-bit WORDS

Pass By Reference vs. Pass By Value

- Defn:
- Registers or the stack can be used to pass by value or by reference.
- Why does it matter?

ASCII-Binary Conversions

- In a high-level language, you just read the number:
 - read (num) or
 - scanf(“%d”, &num),
 - cin >> num, or...
- What’s going on behind the scenes?
- Say the user enters 361
 - They enter 3 separate keys: “3”, “6”, “1”
 - These come in as ASCII values
 - They must be converted into the integer 361 and stored.

Algorithm (in Pseudo-Code)

Result <- 0
Multiplier <- 10

Convert:

Get a character

If not a digit char, go to Finish

Else

Strip the ASCII bias off of the digit character (subtract 30h)

Result <- Result * multiplier + digit

go to Convert

Finish:

```

mov     bx, 0           ;bx contains running total
mov     di, 10          ;di is the multiplier
sub     si, s           ;s=0 means positive,
                       ;s=1 means negative

mov     ah, 01h        ;read one character from keyboard
int     21h
cmp     al, '-'         ;was character a minus sign?
jne     Convert        ;no, it was a digit
inc     si             ;yes, remember to negate at end
jmp     GetNext        ;go get first digit

Convert:
cmp     al, '0'        ;when you read a non-digit, you're
jb     Finish         ;done
cmp     al, '9'
ja     Finish

mov     cx, ax         ;al will be used by MUL
mov     ax, bx         ;put running total in ax
mul     di             ;multiply AX by 10

sub     cx, 30h        ;convert single digit to binary
mov     ch, 0
add     ax, cx         ;add new digit to running total

mov     bx, ax         ;put running total back in BX,
                       ;because AX will be needed by next
                       ;input function

GetNext:
mov     ah, 1          ;prepare to read in next character
int     21h
jmp     Convert

Finish:
cmp     si, 1          ;see if you need to negate number
jne     NotNeg
neg     bx             ;yes you do

NotNeg:
nop                   ;check number with debugger

```

Binary-ASCII Conversions

- To print a number, it needs to be converted from binary to ASCII.
- For example:
 - 361 decimal -> “3”, “6”, “1”

Algorithm (in Pseudo-Code)

```
put number in AX
repeat
  divide AX by 10
  convert remainder (DL) to ASCII
  save remainder in a buffer
until AX = 0

-> numbers are generated in *reverse order*
```