

- Quick Stack Review
- Passing Parameters on the Stack
- Binary/ASCII conversion





CALL and RET

- CALL pushes IP on the stack (recall, IP holds the address of the next instruction), puts the address of the label (subroutine) into IP.
- RET pops the stack into IP to return to the point at which the subroutine was called.

Passing Parameters using Registers

- We've been using registers to pass parameters (when? for the INT 21h calls!)
- You can also use registers to pass a base address of a group of memory locations.
- Example our keyboard buffer:
 - maxlen db
 20 ; max chars to input

 actuallen db
 ? ;DOS will put the

 number read here

 inbuf
 db
 20 dup (' ') ; where DOS

 will put the data read in
 - If you load the address of maxlen into a register (using MOV BX offset maxlen) before calling the subroutine then the subroutine will be able to access the other data members by using indirect addressing.

Passing Parameters Using Registers, cont.

- Advantages:
 - Easy for passing a small number of parameters
 - Useful for passing the start address of a block of parameters
- Problems:
 - There are only four general purpose registers (AX, BX, CX, DX) and they are needed for many things.

Passing Parameters using the Stack

• For several parameters, it's better to use the stack.

(this is what high level languages do)

- Before the CALL, push the parameters on the stack.
- But... you can't just pop them off again in your program... why?

Passing Parameters on the Stack, cont.

- When you execute CALL, it pushes IP on the stack.
- If your called routine executes a pop it won't get a parameter, it will get the return IP.



Getting at Our Parameters

- You can pop the IP off the stack and then get your data.
 - Don't do this!!!!
 - If you forget to put IP back on top, you will return off into nevernever land!
- A better solution use the BP register!
 - BP works like BX except it holds an offset from SS (recall BX holds an offset from DS)

Indirect Addressing Using BP

mov ax, [bx] ; this moves the ; word at DS:BX ; into AX

mov dx, [bp] ; this moves the ; word at SS:BP ; into DX

• you can also use indirect addressing with displacement:

mov ax, [bp + 2] ;take the address ;in BP, add 2 ;load AX with the ;contents of the word at ;that address.



Stack Frame Example

• After calling our subroutine:



Note: stack drawn as 16-bit WORDS

old IP? The return address – this is the offset of the instruction immediately after the call.





TITLE Demonstrates parameter passing on the stack ; ; This program adds a constant value to each element of an array. ; It uses a procedure, and parameters are passed on the stack in ; the following order: the start address of the array, the constant ; value to be added, and the number of elements in the array.				
.model small .stack 100H				
AddNumber EQU 12H ;number to be added ArraySize EQU 9H				
.data TestArray DW 3AH, 4AH, 5AH, 6AH, 7AH, 1234H, 5678H DW 9ABCH, 0DEFH				
.code .startup				
mov	ax. offset TestArray			
push	ax			
mov	ax, AddNumber			
push	ax			
mov	ax, ArraySize			
push	ax			
call	ArrayInc			
nop				
.exit				

ArrayInc PROC		NEAR		
	push	bp		
	mov	bp, sp	;set up frame for this call ;save registers destroyed by this	
	push	ax		
	push bx		;procedure	
	push	cx		
	push	dx		
	mov	cx, [bp+₄	4]	;size of array
	mov	ax, [bp+0	5]	;the constant to be added
	mov	bx, [bp+	8]	;start address of array
next:	mov	dx [bx]	.get array	v element
nonti	add	dx ax	add con	stant to it
	mov	[bx] dx	store it l	pack in array
	add	bx 2	point to	next array element
	loop	next	,point to	nent unuj element
	non	dv	restore r	egisters
	non	CX	,10310101	egisters
	non	by		
	pop	ax		
	Pob	u.		
	pop	bp		
	ret		;returns	
ArrayIr	nc ENDP end			

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What's Wrong with this Example?

Cleaning up the Stack

- What you should do is clean up the stack using RET x, where x is the number of bytes to add to SP after the return:
 - RET 6 ;clean up the 6 bytes of ;parameters by adding ;6 to SP after popping IP

Using the Stack for Local Variables • Once you've set up your stack frame, you can use it for local variables as well:



Pass By Reference vs. Pass By Value

- The book talks about pass by value and pass by reference.
 - Pass by value the value of the parameter is passed into the subroutine
 - Pass by reference the address of the parameter is passed into the subroutine (we just saw this in the array example!)
- Either the registers or the stack can be used to pass by value or by reference.
- Why does it matter? Pass by reference is used if you want the procedure to modify a variable.

ASCII-Binary Conversions

- In a high-level language, you just read the number:
 - read (num) or
 - scanf("%d", &num),
 - cin >> num, or...
- What's going on behind the scenes?
- Say the user enters 361
 - They enter 3 separate keys: "3", "6", "1"
 - These come in as ASCII values
 - They must be converted into the integer 361 and stored.

Algorithm (in Pseudo-Code)

Result <- 0 Multiplier <- 10

Convert: Get a character If not a digit char, go to Finish Else Strip the ASCII bias off of the digit character (subtract 30h) Result <- Result * multiplier + digit go to Convert Finish:

	mov	bx, 0	;bx contains running total		
	mov	di, 10	di is the multiplier		
	sub	si, si	;si=0 means positive,		
			;si=1 means negative		
	mov	ah, 01h ;read one character from keyboard			
	int	21h	1h		
	cmp	al, '-' ;was cl	haracter a minus sign?		
	jne	Convert	;no, it was a digit		
	inc	si	;yes, remember to negate at end		
	jmp	GetNext	:go get first digit		
Convert:					
	cmp	al, '0' ;when	you read a non-digit, you're		
	jb	Finish	;done		
	cmp	al, '9'			
	ја	Finish			
	mov	cl, al	;al will be used by MUL		
	mov	ax, bx	put running total in ax		
	mul	di	;multiply AX by 10		
	sub	cl, 30h ;con	vert single digit to binary		
	mov	ch, 0			
	add	ax, cx	;add new digit to running total		
	mov	bx, ax	:put running total back in BX, :because AX will be needed by next ;input function		
GetNext:					
	mov	ah, 1	prepare to read in next character		
	int	21h			
	jmp	Convert			
Finish:					
	cmp	si, 1	;see if you need to negate number		
	jne	NotNeg			
	neg	bx ;yes y	rou do		
NotNeg:					
	nop	;check number with debugger			



- To print a number, it needs to be converted from binary to ASCII.
- For example:
 361 decimal -> "3", "6", "1"

Algorithm (in Pseudo-Code)

put number in AX repeat divide AX by 10 convert remainder (DL) to ASCII save remainder in a buffer until AX = 0

-> numbers are generated in *reverse order*