

Lightweight Active Router-Queue Management for Multimedia Networking

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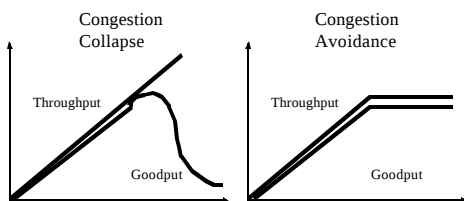


Outline

- Problem
 - Supporting multimedia on the Internet
- Context
 - Drop Tail
 - RED
 - FRED
- Approach
 - CBT
- Evaluation
- Conclusion



Congestion on the Internet

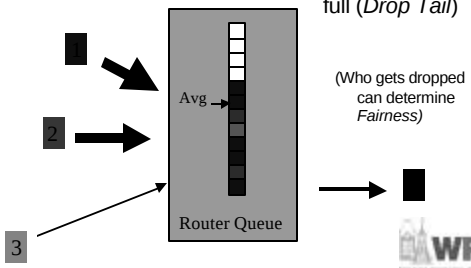


- Drops are the usual way congestion is indicated
- TCP uses congestion avoidance to reduce rate



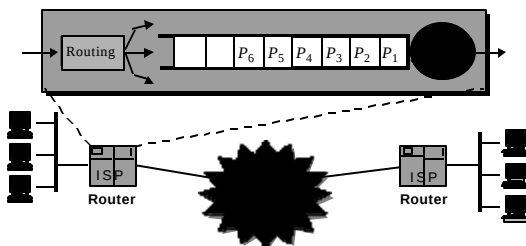
Internet Routers

- Queue to hold incoming packets until can be sent
- Typically, drop when queue is full (*Drop Tail*)



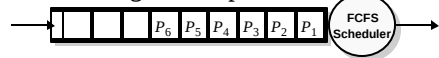
Router-Based Congestion Control Solution 2: Closed-loop congestion control

- Normally, packets are only dropped when the queue overflows
 - "Drop-tail" queueing



Buffer Management & Congestion Avoidance

The case against drop-tail



- Large queues in routers are a bad thing
 - End-to-end latency is dominated by the length of queues at switches in the network
- Allowing queues to overflow is a bad thing
 - Connections that transmit at high rates can starve connections that transmit at low rates
 - Causes connections to synchronize their response to congestion and become unnecessarily bursty



Buffer Management & Congestion Avoidance

Random early detection (RED) packet drop

- Use an exponential average of the queue length to determine when to drop
 - Accommodates short-term bursts
- Tie the drop probability to the weighted average queue length
 - Avoids over-reaction to mild overload conditions

Buffer Management & Congestion Avoidance

Random early detection (RED) packet drop

- Amount of packet loss is roughly proportional to a connection's bandwidth utilization
 - But there is no a priori bias against bursty sources
- Average connection latency is lower
- Average throughput ("goodput") is higher

Buffer Management & Congestion Avoidance

Random early detection (RED) packet drop

- RED is controlled by 5 parameters
 - $qlen$ — The maximum length of the queue
 - w_q — Weighting factor for average queue length computation
 - min_{th} — Minimum queue length for triggering probabilistic drops
 - max_{th} — Queue length threshold for triggering forced drops
 - max_p — The maximum drop probability

Random Early Detection Algorithm

```

for each packet arrival:
  calculate the average queue size ave
  if ave < min_th
    do nothing
  else if min_th < ave < max_th
    calculate drop probability p
    drop arriving packet with probability p
  else if max_th < ave
    drop the arriving packet
  
```

- The average queue length computation needs to be low pass filtered to smooth out transients due to bursts
 - $ave = (1 - w_q)ave + w_q q$

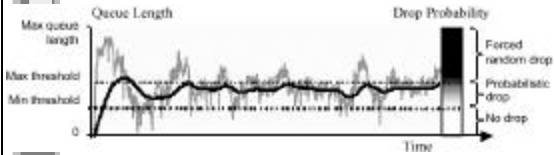
Buffer Management & Congestion Avoidance

Random early detection (RED) packet drop

Random Early Detection Performance

- Floyd/Jacobson simulation of two TCP (*ftp*) flows

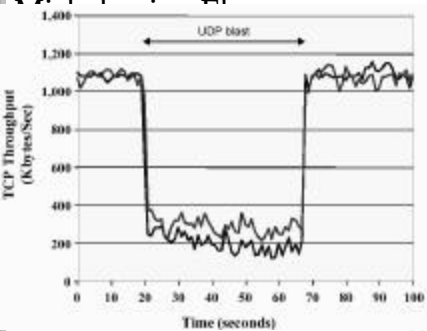
Random Early Detection (RED) Summary



- Controls average queue size
- Drop early to signal impending congestion
- Drops proportional to bandwidth, but drop rate equal for all flows
- Unresponsive traffic will still not slow down!



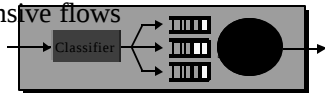
RED Vulnerability to



- TCP performance on a 10 Mbps link under RED in the face of a "UDP" blast



Router-Based Congestion Control Dealing with heterogeneous/non-responsive flows



- TCP requires protection/isolation from non-responsive flows
- Solutions?
 - Employ fair-queuing/link scheduling mechanisms
 - Identify and police non-responsive flows (not here)
 - Employ fair buffer allocation within a RED mechanism



Flows

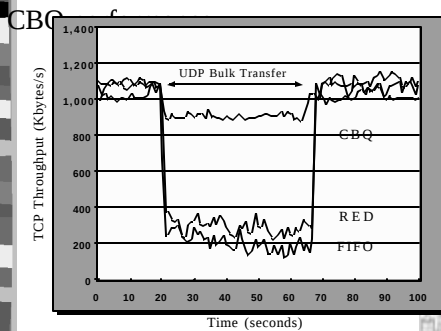
Isolating responsive and non-responsive flows



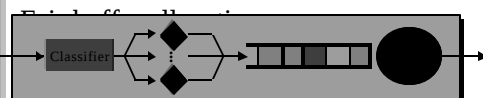
- Class-based Queuing (CBQ) (Floyd/Jacobson) provides fair allocation of bandwidth to traffic classes
 - Separate queues are provided for each traffic class and serviced in round robin order (or weighted round robin)
 - n classes each receive exactly $1/n$ of the capacity of the link
- Separate queues ensure perfect isolation between classes
- Drawback: 'reservation' of bandwidth and state information required



Dealing with Non-Responsive Flows



Dealing with Non-Responsive Flows

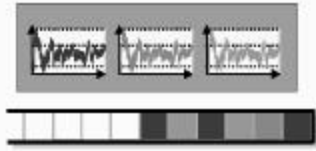


- Isolation can be achieved by reserving capacity for flows within a single FIFO queue
 - Rather than maintain separate queues, keep counts of packets in a single queue
- Lin/Morris: Modify RED to perform fair buffer allocation between active flows
 - Independent of protection issues, fair buffer allocation between TCP connections is also desirable



Flow Random Early Detect (FRED)

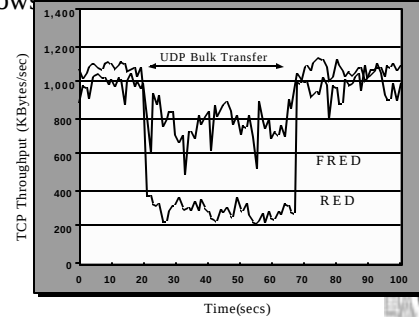
- In RED, 10 Mbps $\rightarrow$ 9 Mbps and 1 Mbps $\rightarrow$.9 Mbps
 - Unfair
- In FRED, leave 1 Mbps untouched until 10 Mbps is down



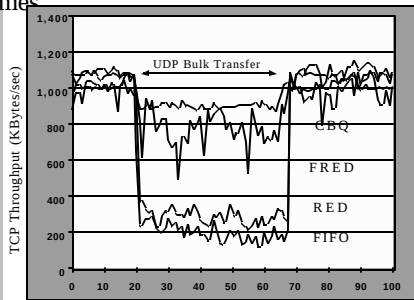
- Separate drop probabilities per flow
- “Light” flows have no drops, heavy flows have high drops



Flow Random Early Detection Performance in the face of non-responsive flows



Congestion Avoidance v. Fair-Sharing TCP throughput under different queue management schemes



- TCP performance as a function of the state required to ensure/approximate fairness



Queue Management Recommendations

- Recommend (Braden 1998, Floyd 1998)
 - Deploy RED
 - + Avoid full queues, reduce latency, reduce packet drops, avoid lock out
 - Continue research into ways to *punish* aggressive or misbehaving flows
- Multimedia
 - Does not use TCP
 - + Can tolerate some loss
 - + Price for latency is too high
 - Often low-bandwidth
 - Delay sensitive



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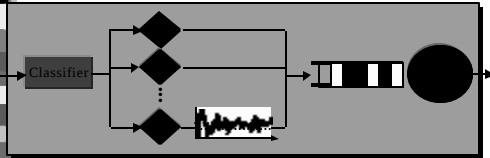


Goals

- Isolation
 - Responsive (TCP) from unresponsive
 - Unresponsive: multimedia from aggressive
- Flexible fairness
 - Something more than equal shares for all
- Lightweight
 - Minimal state per flow
- Maintain benefits of RED
 - Feedback
 - Distribution of drops



Class-Based Threshold (CBT)



- Designate a set of traffic classes and allocate a fraction of a router's buffer capacity to each class
- Once a class is occupying its limit of queue elements, discard *all* arriving packets
- Within a traffic class, further active queue management may be performed

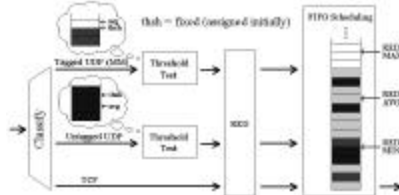


Class-Based Threshold (CBT)

- Isolation
 - Packets are classified into 1 of 3 classes
 - Statistics are kept for each class
- Flexible fairness
 - Configurable thresholds determine the ratios between classes during periods of congestion
- Lightweight
 - State per class and not per flow
 - Still one outbound queue
- Maintain benefits of RED
 - Continue with RED policies for TCP



CBT Implementation



- Implemented in Alt-Q on FreeBSD
- Three traffic classes:
 - TCP
 - marked non-TCP ("well behaved UDP")
 - non-marked non-TCP (all others)
- Subject TCP flows get RED and non-TCP flows to a weighted average queue occupancy threshold test



Class-Based Thresholds

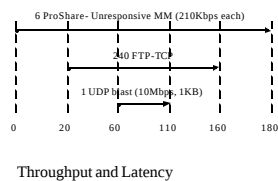
Evaluation



- Compare:
 - FIFO queuing (Lower bound baseline)
 - RED (The Internet of tomorrow?)
 - FRED (RED + Fair allocation of buffers)
 - CBT
 - CBQ (Upper bound baseline)



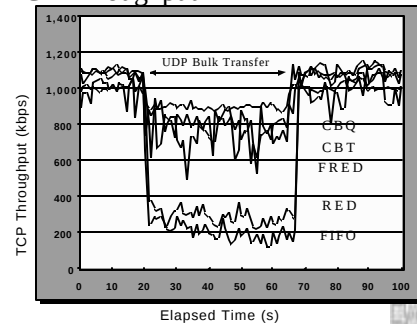
CBT Evaluation Experimental design

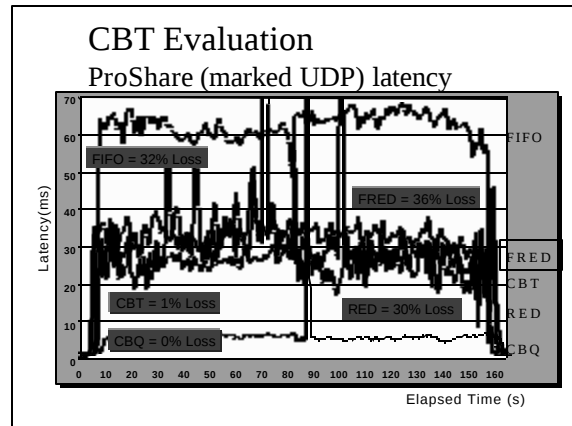
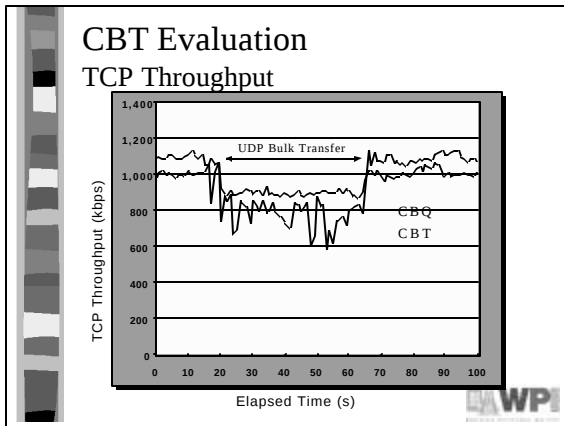


- RED Settings:
 - qsize = 60 pkts
 - max-th = 30 pkts
 - min-th = 15 pkts
 - wq = 0.002
 - max-p = 0.1
- CBT Settings:
 - mm-th = 10 pkts
 - udp-th = 2 pkts

CBT Evaluation

TCP Throughput





- ### Conclusion
- RED/FIFO scheduling not sufficient
 - Aggressive unresponsive flows cause trouble
 - Low bandwidth unresponsive (VoIP) punished
 - CBT provides
 - Benefits of RED for TCP only traffic
 - Isolation of TCP vs. Unresponsive
 - Isolation of Aggressive vs. Low Bandwidth
 - Lightweight overhead
- WPI

Future Work?

WPI

- ### Future Work
- How to pick thresholds?
 - Implies reservation
 - Dynamic adjustments of thresholds (D-CBT)
 - Additional queue management for classes
 - Classes use "Drop Tail" now
 - Extension to other classes
 - Voice
 - Video
- WPI