

# Measurement Study of Low-bitrate Internet Video Streaming

Dmitri Loguinov and Hayder Radha  
CS Dept at CUNY NY and EE/ECE at MSU.

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## Introduction

- Many studies of Internet performance
  - Paxson, Mogul, Caceres...
  - Across countries, many sites
  - Well-connected (often schools on backbone)
- But few look at it from the point of dialup user
  - About 50% of home users dialup
    - Peak, but will remain majority for 3-5 years
  - ISP cannot always do 56 kb/s



## Introduction

- Most studies involve TCP
  - 90-99% of traffic on Internet is TCP
- But MM prefers UDP
  - (Why?)
- Also, TCP uses ACK-based scheme
  - MM protocols prefer NACK to scale (why?)
- Video studies have done few paths



## Introduction

- Video streaming experiment
  - Seven month long
  - MPEG-4 (low-bandwidth) over UDP
  - Over dialup
  - 600 major cities
  - 50 States



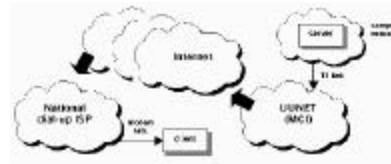
## Outline

- Introduction (done)
- Methodology ←
  - Setup
  - Streaming
  - Client-Server architecture
- Results
- Analysis
- Summary



## Setup

- Clients connected to each long-distance
- Server was in NY




- 3 ISPs in all 50 states
- 1813 different access points
- 1188 major U.S. cities




### Setup

- Dialer
  - Connect to ISP with PPP connection
  - *traceroute* from sender → receiver and receiver → sender
    - Parallel paths
- Detect when modem connection was bad
  - If  $r$  is target bitrate,  $p$  is packet loss
  - If  $B_p < 0.9r$  then bad (toss)
  - If  $B_p$  is  $> 15\%$  then bad (toss)
- Good data was time-stamped
  - Day of week plus 3 eight hour slots each data
  - At least one from each day for each state for each slot




### Streaming

- MPEG-4 stream
  - 2 ten-minute QCIF (176x144) streams
  - $S_1$  14 kbps (Nov-Dec 1999)
  - $S_2$  25 kbps (Jan-May 2000)
- Server split into 576 byte packets
  - With overhead  $S_1$  16 kbps and  $S_2$  27.4 kbps
  - About 6 packets/sec (for  $S_2$ )
- To remove jitter, had delay buffer
  - (*What is this?*)
- Chose 2.7 seconds (1.3 ideal in pilots stud)
  - (*Why might this be a bad idea?*)




### Client-Server Architecture

- Server
  - Multithreaded
  - Bursts of packets (340-500 ms)
- Client
  - Recover lost packets through NACK
  - Collect RTT delay
    - Based on NACK
      - (*When might this not work well?*)
    - Probes every 30 seconds if loss  $< 1\%$
- Evaluated for 9 months
  - Whew!



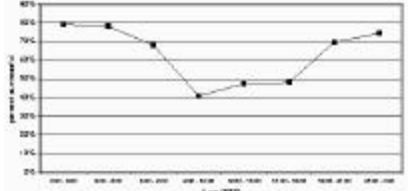
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### Results


- Two datasets
  - $D_1$  16,783 connections, 8429 successful
  - $D_2$  17,465 connections, 8423 successful
  - To get MPEG-4, need 2 attempts on avg




- Time of day matters

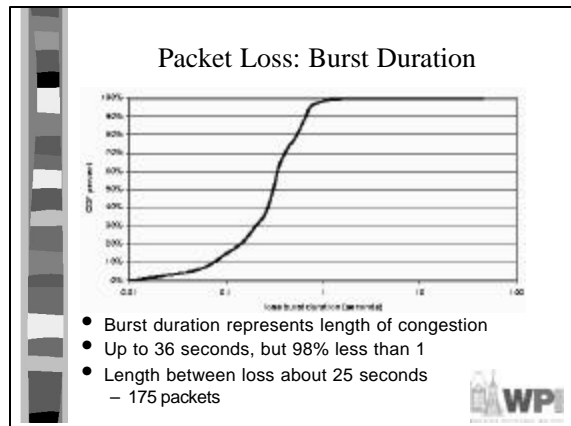
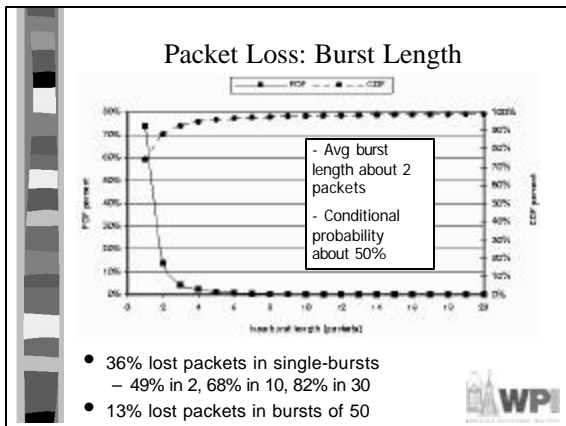
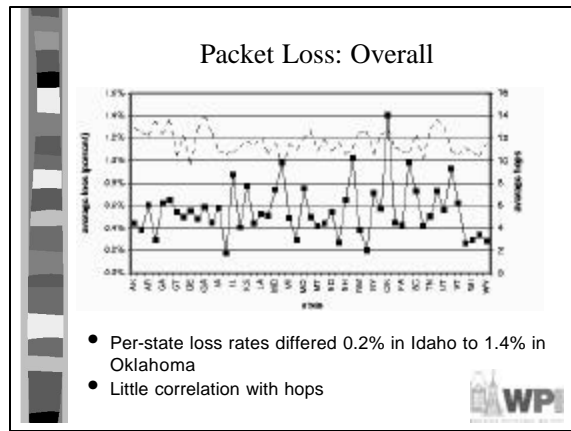
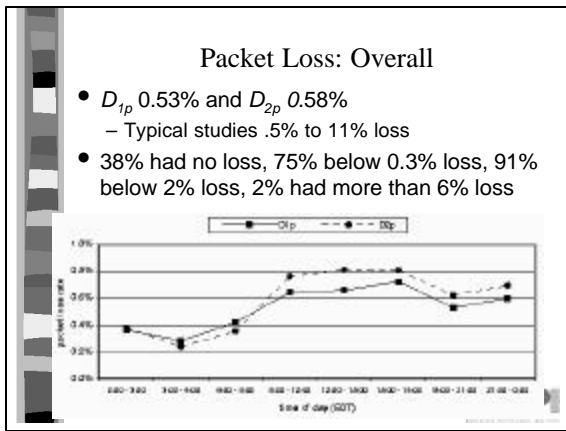
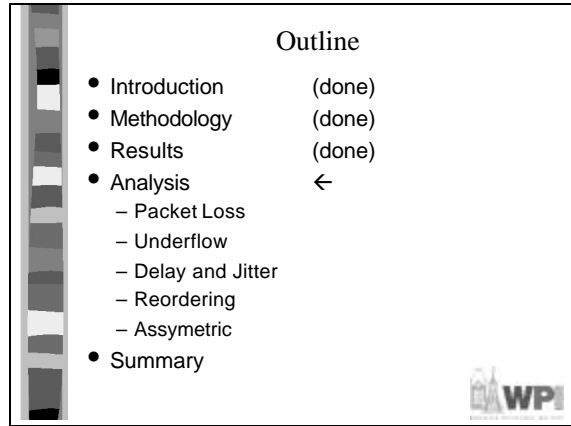
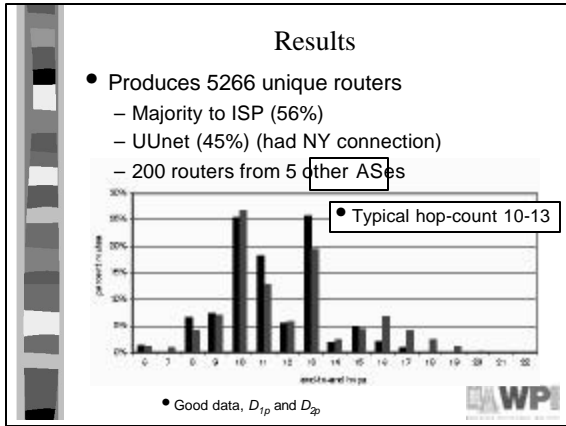


### Results




- $D_1$  had 962 dialup points, 637 cities
- $D_2$  had 880 dialup points, 575 cities






### Outline

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  - Packet Loss (done)
  - Underflow ←
  - Delay and Jitter
  - Reordering
  - Assymmetric
- Summary




### Video Quality

- No user studies, no PSNR
  - Do not provide insight into network
- Instead, consider *underflow event*
  - When there is no frame to play
- Consider repair?
  - No standardized techniques to conceal loss
  - Techniques range from simple to complex
  - Performance depends upon:
    - Motion in video
    - Type of frame from packet (I, P, B)
  - Don't want this to be a study evaluating repair
    - Every packet loss may cause an *underflow event*




### Video Quality

- Too much delay can cause underflow
  - Retransmitted packet will still be late
- Too much jitter can cause underflow
  - Retransmitted or original packet late
- Two types of late
  - Completely late (of no use)
  - Partially late (can help decode other frames in GOP)
- GOP: IPPPPPPPPP

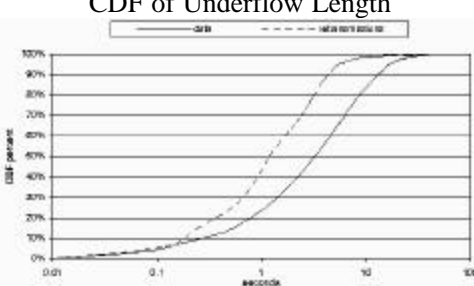


### Underflow Results from Delay and Jitter

- For  $D_{1p}$  and  $D_{2p}$ , 431,000 lost packets
  - 160,000 found after deadline (37%) so no NACK
  - 257,000 (94%) sent NACK and recovered
    - 9,000 recovered late
      - 4000 (about 50%), "rescue" about 5 frames
    - 5,000 never recovered
- Jitter caused 1,100,000 underflow events
  - 98% of underflow events
  - 73% if don't use retransmission
- (*How to improve these numbers?*)




### CDF of Underflow Length



Retransmit: 25% late by 2+, 10% by 5+, 1% by 10+


Jitter: 25% by 7+, 10% by 13+, 1% by 27

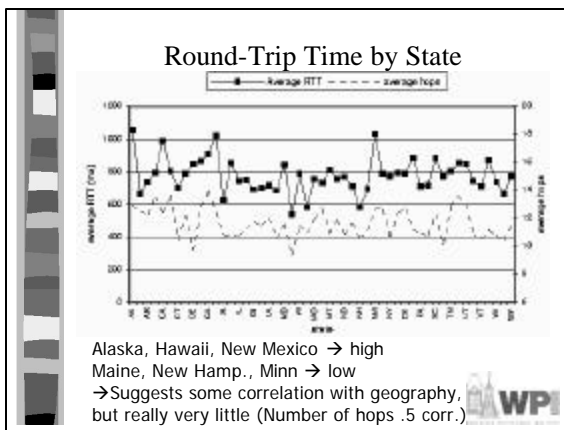
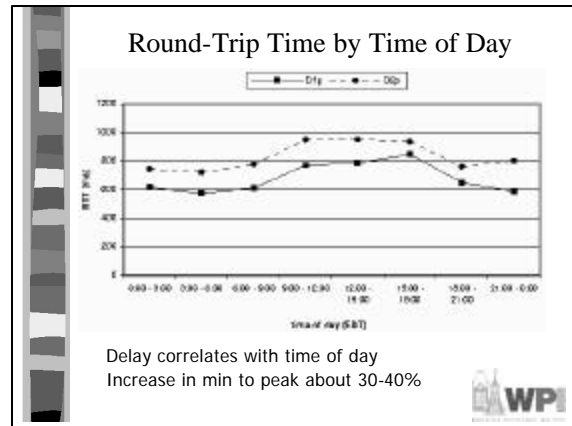
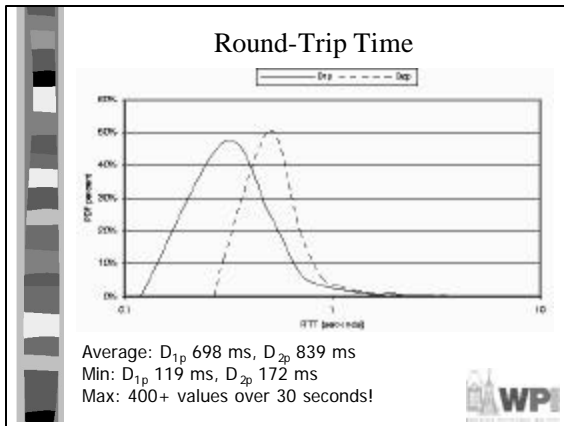
Buffer of 13 seconds would recover 99% of retransmissions and 84% of jitter



### Outline

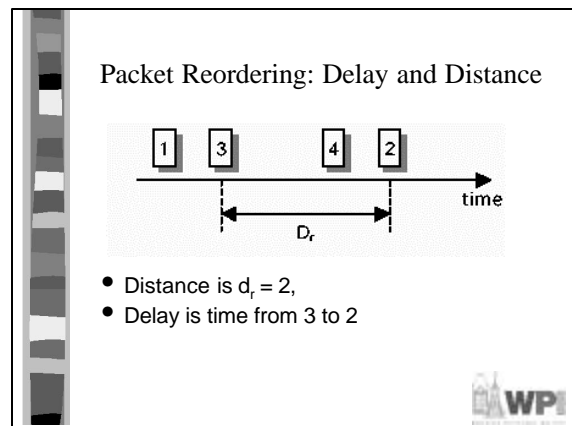
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  - Assymmetric
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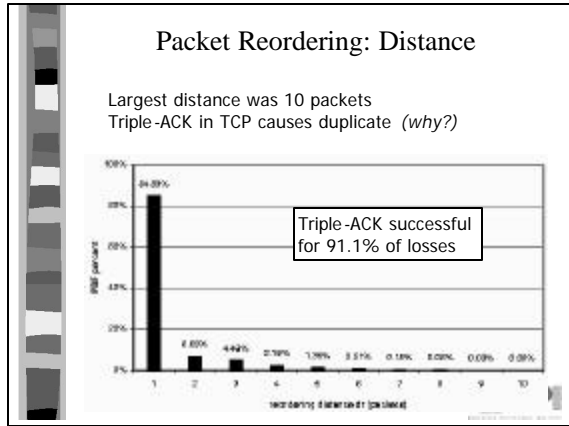
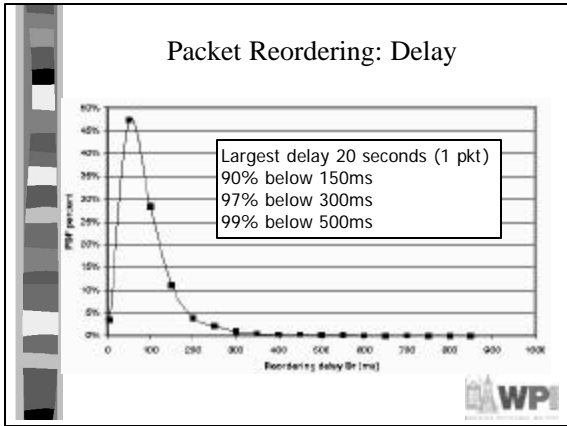




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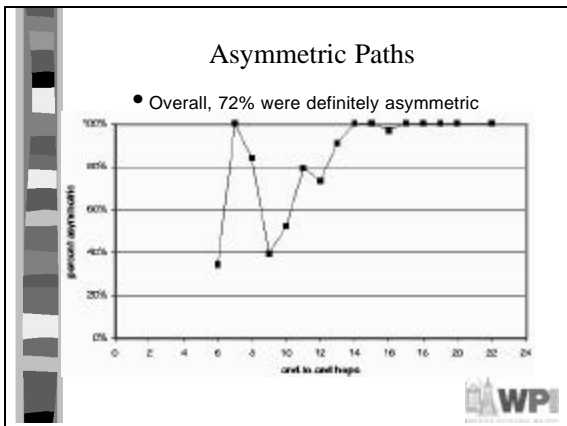
- ### Packet Reordering: Overview
- Gap in sequence numbers indicates loss
    - (When might this fail?)
  - For  $D_{1p}^a$ , 1 in 3 missing packets arrived out of order
    - Simple streaming protocol with NACK could waste bandwidth
  - Average
    - was 6.5% of missing
    - 0.04% of sent packets
  - Of 16,952 sessions, 9.5% have at least 1
    - ½ of sessions from ISP a
  - No correlation with time of day





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- ### Asymmetric Paths
- If number of hops from sender → receiver different than receiver → sender
    - then *asymmetric*
  - If number of hops from sender → receiver same as receiver → sender
    - then *probably symmetric*



- ### Conclusion
- Internet packet loss is bursty
  - Jitter worse than packet loss or RTT
  - RTTs on the order of seconds are possible
  - RTT correlated with number of hops
  - PacktlLoss not correlated with number of hops or RTT
  - Most paths asymmetric