

### Homework #3

**1. True or False:**

- a)  $abab$  matches  $a^* + b^*$     T    F
- b)  $babab$  matches  $b(ab)^*$     T    F
- c) If  $A = \Phi$  then  $AB = \Phi$  for all languages B    T    F
- d) If  $A = \{\epsilon\}$  then  $AB = \Phi$  for all languages B    T    F
- e) If  $A = a^*$  and  $\Sigma = \{a,b\}$ , then  $\Sigma^* - A = b^*$     T    F

#2. Write regular expressions for the set of strings of 0's and 1's with at most one pair of consecutive 1's

**Exercise says "at most one" occurrence of "11", so start with**

$$(\epsilon + 1 + 11)$$

**Now build around it the combination of 1's and 0's that will prevent another occurrence of "11":**

**In front:**

$$(0 + 10)^*$$

**After:**

$$(0 + 01)^*$$

**So altogether:  $(0 + 10)^* (\epsilon + 1 + 11) (0 + 01)^*$**

#3. Draw the graph for the following DFA and then convert to regular expression.

	<b>0</b>	<b>1</b>
* $\rightarrow$ p	s	p
q	p	s
r	r	q
s	q	r

This gives you an idea how messy these are, and how you'd really like to plug them into a computer program!

Book's method, removing q

$$\alpha_{pp}^{pqrs} = \alpha_{pp}^{prs} + \alpha_{pq}^{prs} (\alpha_{qq}^{prs})^* \alpha_{qp}^{prs}$$

$$\alpha_{pp}^{prs} = 1^*$$

$$\alpha_{pq}^{prs} = 1 * 0 (0 + 10 * 1)$$

$$\alpha_{qq}^{prs} = 01 * 0 (0 + 10 * 1) + 1 (0 + 10 * 1)$$

$$\alpha_{qp}^{prs} = 01 *$$

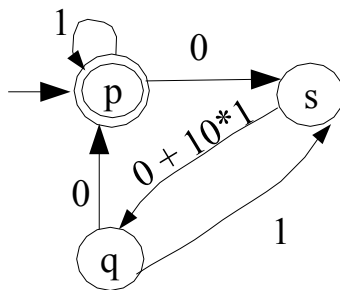
$$\text{So } \alpha_{pp}^{pqrs} = 1 * + 1 * 0 (0 + 10 * 1) (01 * 0 (0 + 10 * 1) + 1(0 + 10 * 1)) * 01 *$$

**I'm sure this can be simplified. If anyone wants to send me their simplification, I'll post it with full credit to the sender!**

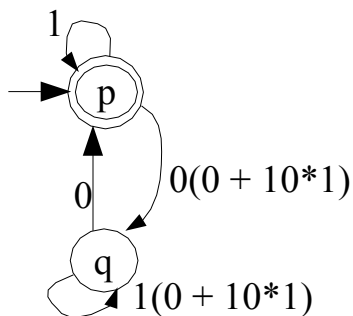
Other techniques Other texts have us eliminate states until we just have an initial and a final state (same here) left:

Step 1:

Eliminating "r":

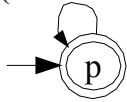


Step 2 Then eliminating "s":



Step 3: Eliminating "q"

$1 + (0(0 + 10^*1))(1(0 + 10^*1))^*0$



Now we \* the expression on the loop to get the regular expression:

$(1 + (0(0 + 10^*1))(1(0 + 10^*1))^*0)^*$

Are these two the same? I hope so, and again, using the equivalencies and many hours of our time, we could probably show it.

#4. Does  $(R+S)^*S = (R^*S)^*$  Justify your answer

Counter-example: if  $\epsilon$  is not in  $S$ , then they are not equal.

#5 - #6. Given  $R$  is a regular language and  $N$  is a non-regular language

#5. (5 Points) Suppose  $X$  is a language such that  $N = \sim X$  ( $\sim$  means complement). Does it follow that  $X$  must be regular? If so, state why. If not, does it follow that  $X$  must be non-regular? If so, state why. If neither of these is true name 1) a specific non-regular  $N$  such that  $N = \sim X$  with  $X$  non-regular and 2) a specific non-regular  $N$  satisfying  $N = \sim X$  with  $X$  regular.

$X$  cannot be regular because then its complement would be regular and it is given that  $N$  is not regular.

#6. (5 Points) Suppose  $X$  is a language such that  $X = R \cap N$ . Does it follow that  $X$  must be regular? If so, state why. If not, does it follow that  $X$  must be non-regular? If so, state why. If neither of these is true name 1) a specific non-regular  $N$  and a regular  $R$  such that  $X = R \cap N$  with  $X$  non-regular and 2) a specific non-regular  $N$  and a regular  $R$  satisfying  $X = R \cap N$  with  $X$  regular.

Neither is true:

1) Let  $R = 0^*1^*$  and  $N = 0^n1^n$ . Then the intersection is  $0^n1^n$  which is non-regular.

2) Let  $R = \emptyset$  and  $N = 0^n1^n$ . Then the intersection is  $\emptyset$  which is regular.

#7. Create a dfa to accept  $(0+1)^*1(0+1)^*$

